Improving Cloud Security with Attacker Profiling

Bryan D. Payne Engineering Manager, Platform Security



NETFLIX Who is out to get me? What do they want? Why are we losing?



Platform Security at Netflix

Platform Security Overview

- Content Protection - Device Security

> Device or Browser



Platform Security

- Foundational Security Services
- Security in Common Platform
- Security by Default in base AMI



Classic Security VIa AWS

CloudHSM

Key Management Supply Chain

Ubiquitous Security

- Partner with other teams
- Make security transparent (or easy)
- Focus on common components
- Also focus on strategic risks







Who is out to get me?









BBC Newsnight, 11 February 2010 <u>https://www.youtube.com/watch?v=1pMuV2o4Lrw</u>



Murdoch et al, Chip and PIN is Broken, IEEE Symposium on Security and Privacy, 2010





Greenberg, X-Ray Scans Expose and Ingenious Chip-and-PIN Card Hack, Wired, 19 October 2015

Attacker Motivations

- financial / business
- political / idealogical
- revenge
- demonstration
- fun

Political & Industrial Espionage

Financial

Financial & Idealogical

Financial, Revenge, Fun

Fun, Demonstration

OpenStack Security Guide (CC BY 3.0) http://docs.openstack.org/sec/





- Little trust in authorities
- Desire control
- Hacker life kept secret
- "Don't foul your own nest"

PROFILING HACKERS

The Science of Criminal Profiling as Applied to the World of Hacking

CRC Press

Attacker Characteristics

- creative and brilliant
- curious
- motivated
- shy in real life
- comfortable with computers

"Yes, I am a criminal. My crime is that of curiosity." The Hacker Manifesto

NETFLIX Attack Characteristics

- access (nmap, exploit, configuration error, etc)
- file cleaners
- backdoor •
- password cracking •
- monitor system admin •
- proceed with goals (files, network sniffing, etc) •



Photo Credit: Google http://www.google.com/about/datacenters/gallery/





What do they want?



"Diamonds" by Swami http://flickr.com/photos/swamibu/1182138940/. Licensed under CC BY 2.0 via Commons

NE C



1



19 February 2003 BBC News

http://news.bbc.co.uk/2/hi/europe/2782305.stm

Diamond heist baffles police

Belgian police are trying to unravel events behind a daring robbery in the diamond-cutting capital of the world.

Thieves cleared out 123 of the 160 vaults in the maximum security cellars at Antwerp's Diamond Centre at the weekend, but the raid was only discovered the next day.

The precise value of the stolen diamonds is not known, but Belgian media have speculated it could run to millions of dollars.

Diamond traders in the city have been shocked by the audacity of the robbery and fear it could be a blow to their industry.

Inside job?

The Diamond Centre building, located in the heart of Antwerp's historic diamond district, is closely guarded.

There are surveillance camers, entry codes, 24-hour security guards, and even cameras in the vaults.

But with no signs of a break-in, police suspect the thieves could have had inside help and have been questioning staff and owners of the safes at the centre.

Antwerp's Diamond High Council, which represents the gemstone traders, has admitted the robbery could have serious implications for an industry proud of its discretion and security.



Antwerp: At the centre of the diamond trade for 400 years



Joshua Davis. *The Untold Story of the World's Biggest Diamond Heist.* Wired, http://archive.wired.com/politics/law/magazine/17-04/ff_diamonds

- 1. Combination dial
- 2. Keyed lock
- 3. Seismic sensor
- 4. Locked steel grate
- 5. Magnetic sensor
- 6. External security camera
- 7. Keypad to disarm sensors
- 8. Light sensor
- 9. Internal security camera
- 10. Heat / motion sensor













T··Mobile·

Experian**





- USG employee background checks & fingerprints
- Credit cards
- User data
- PPI: SSN, driver's license, phone, address, DoB, etc
- Passwords





3 1

Photo Credit: Jonathunder (CC BY-SA 3.0) https://en.wikipedia.org/wiki/Bank_vault#/media/File:WinonaSavingsBankVault.JPG

-





 $risk \propto threat \cdot vulnerability \cdot consequence$



$risk \propto threat \cdot vulnerability \cdot consequence$

asset

attack vectors

controls





http://lockheedmartin.com/content/dam/lockheed/data/isgs/documents/Threat-Driven%20Approach%20whitepaper.pdf



http://lockheedmartin.com/content/dam/lockheed/data/isgs/documents/Threat-Driven%20Approach%20whitepaper.pdf

NETFLIX Cloud Attack Graphs

- Cloud account credentials
- Instance account credentials
- Your employees, supply chains, code
- Provider's employees, supply chains, code
- Corporate network
- Build pipeline

ode ains, code

Why are we losing? ... and how can we improve?

Increasing Security Investment Increasing Security Engineering Efficiencies





Simple Libraries

(e.g., python-cryptography)

from cryptography.fernet import Fernet

key = Fernet.generate_key() f = Fernet(key)ciphertext = f.encrypt(b"A message.") plaintext = f.decrypt(ciphertext)

Traditional Libraries

(e.g., openssl)

#include <openssl/conf.h> /* Initialise the encryption operation. IMPORTANT - ensure you use a key #include <openssl/evp.h> * and IV size appropriate for your cipher #include <openssl/err.h> * In this example we are using 256 bit AES (i.e. a 256 bit key). The #include <string.h> * IV size for *most* modes is the same as the block size. For AES this * is 128 bits */ int main(int arc, char *argv[]) if(1 != EVP_EncryptInit_ex(ctx, EVP_aes_256_cbc(), NULL, key, iv)) handleErrors(); /* Set up the key and iv. Do I need to say to not hard code these in a * real application? :-) /* Provide the message to be encrypted, and obtain the encrypted output. */ * EVP_EncryptUpdate can be called multiple times if necessary /* A 256 bit key */ if(1 != EVP_EncryptUpdate(ctx, ciphertext, &len, plaintext, plaintext_len)) unsigned char *key = "01234567890123456789012345678901"; handleErrors(); ciphertext_len = len; /* A 128 bit IV */ unsigned char *iv = "01234567890123456"; /* Finalise the encryption. Further ciphertext bytes may be written at * this stage. /* Message to be encrypted */ */ unsigned char *plaintext = if(1 != EVP_EncryptFinal_ex(ctx, ciphertext + len, &len)) handleErrors(); "The quick brown fox jumps over the lazy dog"; ciphertext_len += len; /* Buffer for ciphertext. Ensure the buffer is long enough for the /* Clean up */ * ciphertext which may be longer than the plaintext, dependant on the EVP_CIPHER_CTX_free(ctx); * algorithm and mode */ return ciphertext_len; unsigned char ciphertext[128]; /* Buffer for the decrypted text */ unsigned char decryptedtext[128]; int decrypt(unsigned char *ciphertext, int ciphertext_len, unsigned char *key, unsigned char *iv, unsigned char *plaintext) int decryptedtext_len, ciphertext_len; EVP_CIPHER_CTX *ctx; /* Initialise the library */ ERR_load_crypto_strings(); int len; OpenSSL_add_all_algorithms(); OPENSSL_config(NULL); int plaintext_len; /* Encrypt the plaintext */ /* Create and initialise the context */ ciphertext_len = encrypt(plaintext, strlen(plaintext), key, iv, if(!(ctx = EVP_CIPHER_CTX_new())) handleErrors(); ciphertext); /* Initialise the decryption operation. IMPORTANT - ensure you use a key /* Do something useful with the ciphertext here */ * and IV size appropriate for your cipher printf("Ciphertext is:\n"); * In this example we are using 256 bit AES (i.e. a 256 bit key). The BIO_dump_fp(stdout, ciphertext, ciphertext_len); * IV size for *most* modes is the same as the block size. For AES this * is 128 bits */ /* Decrypt the ciphertext */ if(1 != EVP_DecryptInit_ex(ctx, EVP_aes_256_cbc(), NULL, key, iv)) decryptedtext_len = decrypt(ciphertext, ciphertext_len, key, iv, handleErrors(); decryptedtext); /* Provide the message to be decrypted, and obtain the plaintext output. /* Add a NULL terminator. We are expecting printable text */ * EVP_DecryptUpdate can be called multiple times if necessary decryptedtext[decryptedtext_len] = '\0'; if(1 != EVP_DecryptUpdate(ctx, plaintext, &len, ciphertext, ciphertext_len)) /* Show the decrypted text */ handleErrors(); printf("Decrypted text is:\n"); plaintext_len = len; printf("%s\n", decryptedtext); /* Finalise the decryption. Further plaintext bytes may be written at /* Clean up */ * this stage. EVP_cleanup(); */ ERR_free_strings(); if(1 != EVP_DecryptFinal_ex(ctx, plaintext + len, &len)) handleErrors(); plaintext_len += len; return 0; /* Clean up */ EVP_CIPHER_CTX_free(ctx); int encrypt(unsigned char *plaintext, int plaintext_len, unsigned char *key, unsigned char *iv, unsigned char *ciphertext) return plaintext_len; EVP_CIPHER_CTX *ctx; [edit] int len; int ciphertext_len;

/* Create and initialise the context */ if(!(ctx = EVP_CIPHER_CTX_new())) handleErrors();



Sidebar: Key Management @Netflix



Simple Framework for Key Handling

	Throughput	Protection	It's Exposed!	It lives
Construction Constitution Const	High	Low	No biggie	In lots of VMs
Medium Sensitivity	Medium	Medium	It'll be a long week.	In very few VMs
High Sensitivity	Low	High	No. Just. No.	In Special Hardware



Use Case of a Key Implies Handling Requirements

TLS Session Key - Fast, Handled in Dynamic Environment But easy to have a reasonable policy if we lose it

•

Certificate Authority Private Key - Maybe not used so much

• Probably way more important that you just don't lose it



Cryptex - Our Framework for Key Handling

Netflix Business Application

Web Server Logic

Cryptex Client Library

Netflix IPC Components (Ribbon/Hystrix/etc)





Cloud HSMs - Dedicated Hardware



"Low" Key Handling



Key Exported Out to Every Client Extremely High Throughput

• Client Library Attempts to be Mindful of Key Handling



"Medium" Key Handling



Every Operation is a REST Call

Luckily we don't have many bulk encrypt use cases for these • Cryptex servers not publicly facing; ostensibly harder to get onto



"High" Key Handling



Every Operation is a call to specialized hardware

HSM API challenging relative to REST calls (only Cryptex does it) Very constrained throughput;VM side channel attacks negated



"Asymmetric" Key Handling

Netflix Business Application

Cryptex Client Library

Verify

We support the basics: AES, HMAC-SHA, RSA Optimize RSA verify/encrypt by pushing public key to edge At scale computational intensity of RSA quite apparent



GetKey(ID=111)

Resp(PubValue=iXKQ...)

Client Auth TLS







LIBERTY





Photo Credit: Kayamon (CC BY-SA 3.0) https://en.wikipedia.org/wiki/File:Penny_Harvest_Field_2007.jpg



Managing Security at Scale

Spinnaker

what you deploy



deployment pipeline

runtime consistency

Attackers Are Creative

- 802.11a/b/g/n/ac
- Bluetooth
- Gigabit Ethernet
- Out-of-band SSH access over 4G/GSM cell networks



https://www.pwnieexpress.com/product/pwn-plug-r3penetration-testing-device/



A team participating in a CTF competition at DEFCON 17

Photo Credit: Nate Grigg (CC BY 2.0) http://www.flickr.com/photos/nateone/3792232737/



Questions?

bryanp@netflix.com http://bryanpayne.org

[PS... I'm hiring!]

