## Stateful Stream Processing with Apache Flink

The Power of Snapshots



QCon San Francisco, 2017



# dataArtisans





PLATFORM

Original creators of Apache Flink® dA Platform 2 Open Source Apache Flink + dA Application Manager



# Stream Processing

## What changes faster? Data or Query?

Data changes slowly compared to fast changing queries

ad-hoc queries, data exploration, ML training and (hyper) parameter tuning

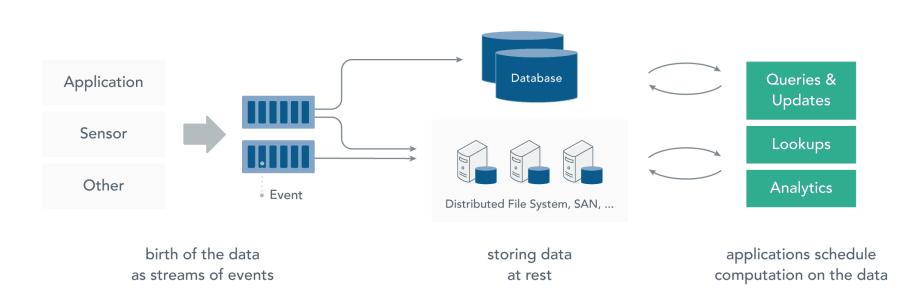
> Batch Processing Use Case

Data changes fast application logic is long-lived

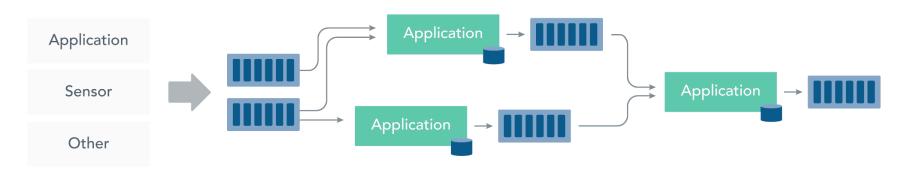
continuous applications, data pipelines, standing queries, anomaly detection, ML evaluation, ...

> Stream Processing Use Case

#### **Batch Processing**



#### Stream Processing



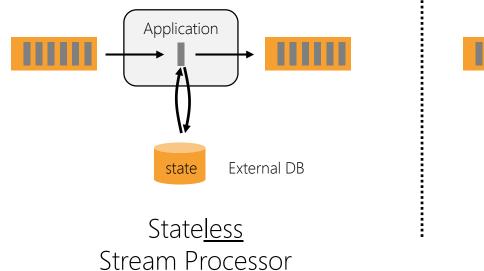
birth of the data as streams of events

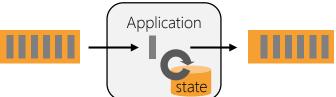
applications computing over event data streams



# Stateful Stream Processing

#### Moving State into the Processors





#### State<u>ful</u> Stream Processor



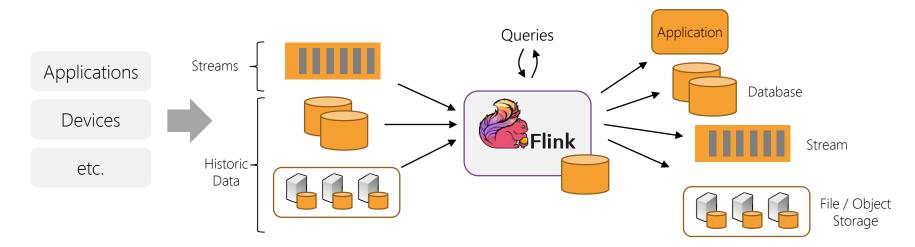


Apache Flink

#### Apache Flink in a Nutshell



Stateful computations over streams real-time and historic fast, scalable, fault tolerant, in-memory, event time, large state, exactly-once





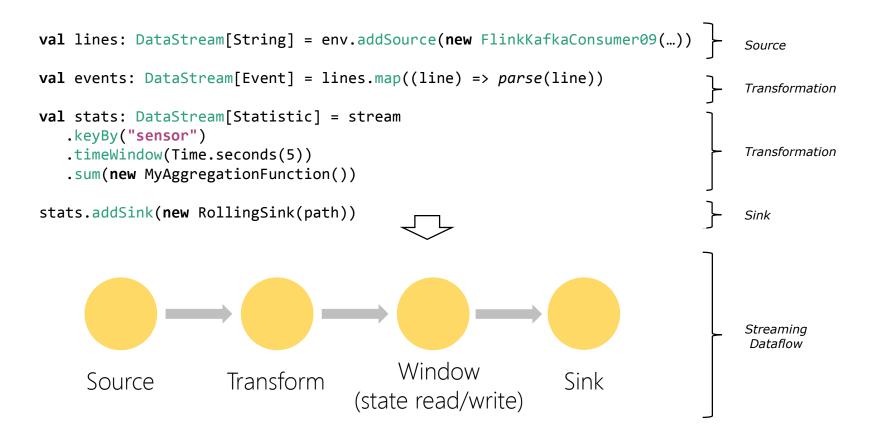


#### The Core Building Blocks

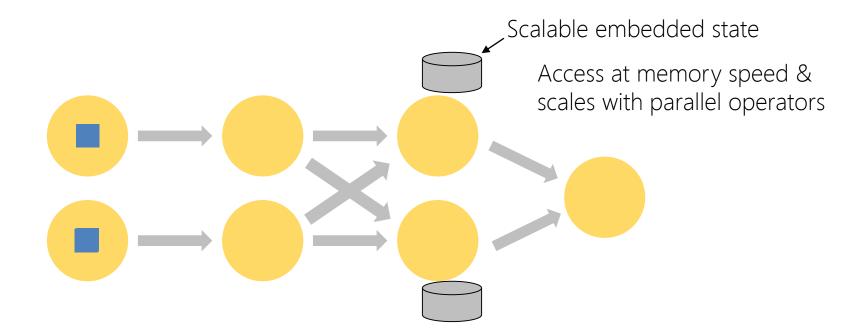
#### Event Streams State (Event) Time Snapshots

real-time and hindsight complex business logic consistency with out-of-order data and late data forking / versioning / time-travel

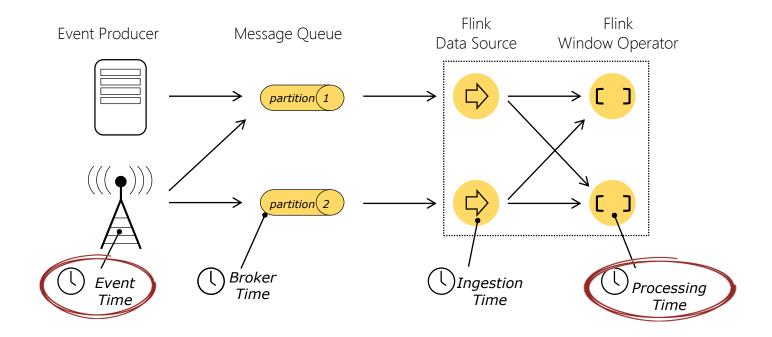
#### Stateful Event & Stream Processing



#### Stateful Event & Stream Processing



#### Event time and Processing Time



Event time, Watermarks, as in the Dataflow model

#### **Powerful Abstractions**

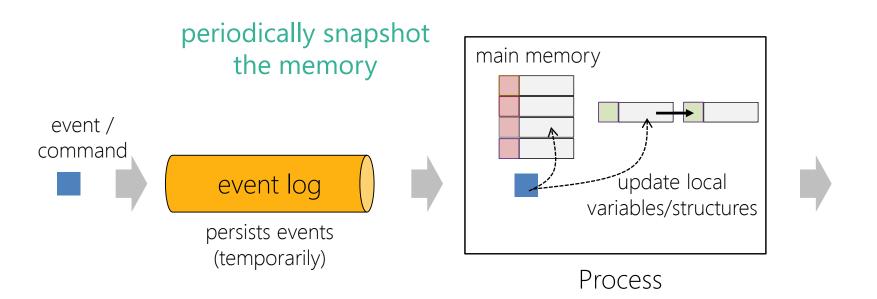


Layered abstractions to navigate simple to complex use cases SELECT room, TUMBLE\_END(rowtime, INTERVAL '1' HOUR), AVG(temp) FROM sensors GROUP BY TUMBLE(rowtime, INTERVAL '1' HOUR), room High-level Stream SQL / Tables (dynamic tables) Analytics API **val** stats = stream Stream- & Batch .keyBy("sensor") DataStream API (streams, windows) .timeWindow(Time.seconds(5)) Data Processing  $.sum((a, b) \rightarrow a.add(b))$ Stateful Event-Process Function *(events, state, time)* Driven Applications def processElement(event: MyEvent, ctx: Context, out: Collector[Result]) = { // work with event and state (event, state.value) match { ... } out.collect(...) // emit events state.update(...) // modify state // schedule a timer callback ctx.timerService.registerEventTimeTimer(event.timestamp + 500) 15



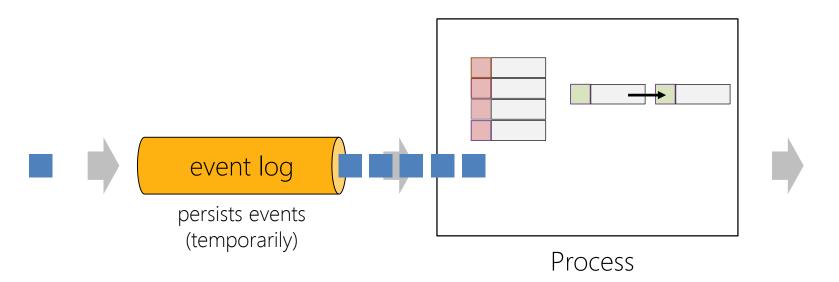
# Distributed Snapshots

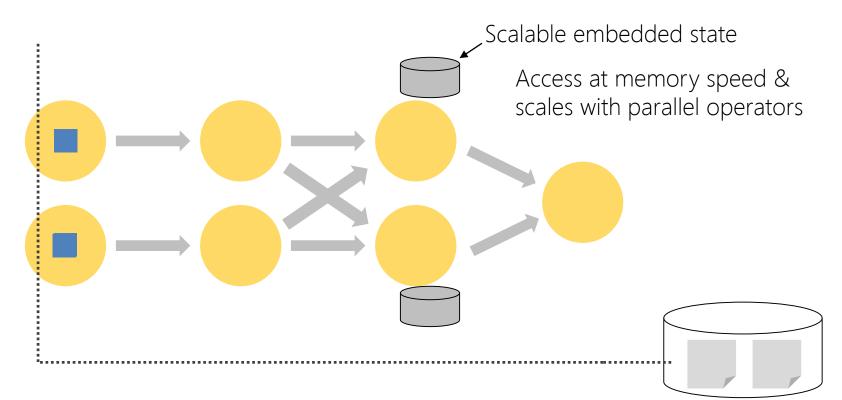
### Event Sourcing + Memory Image



### Event Sourcing + Memory Image

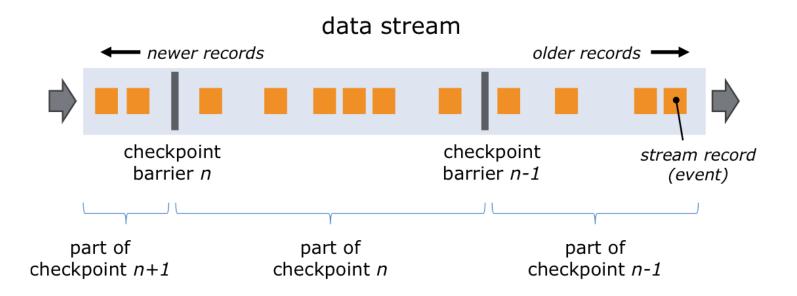
Recovery: Restore snapshot and replay events since snapshot

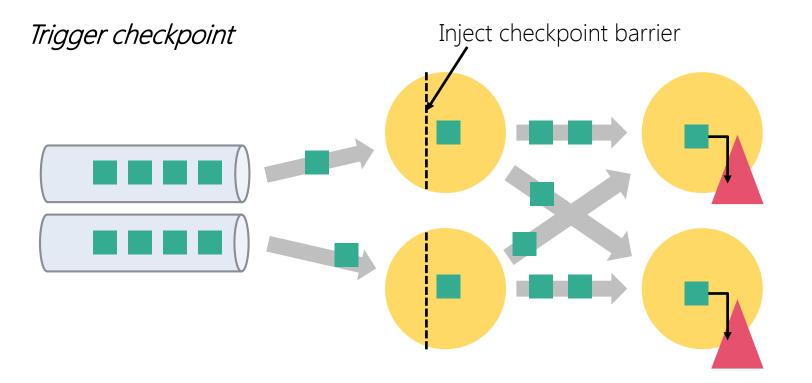


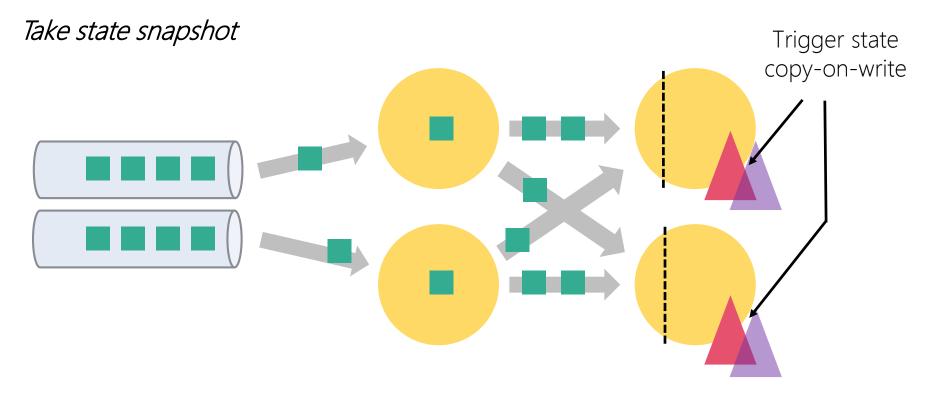


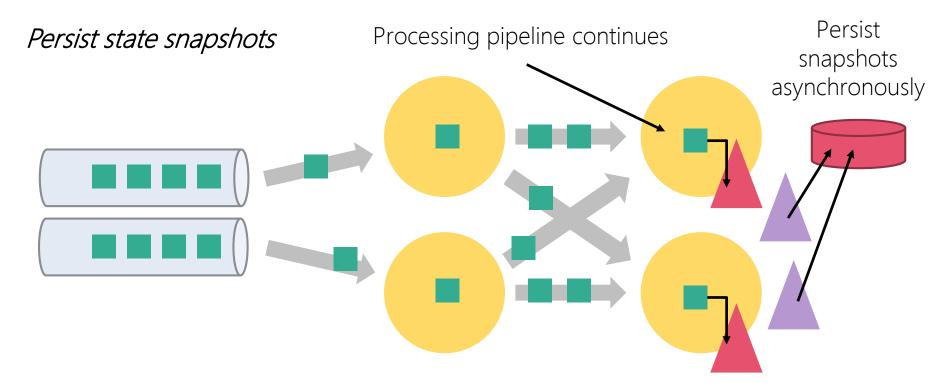
#### **Checkpoint Barriers**

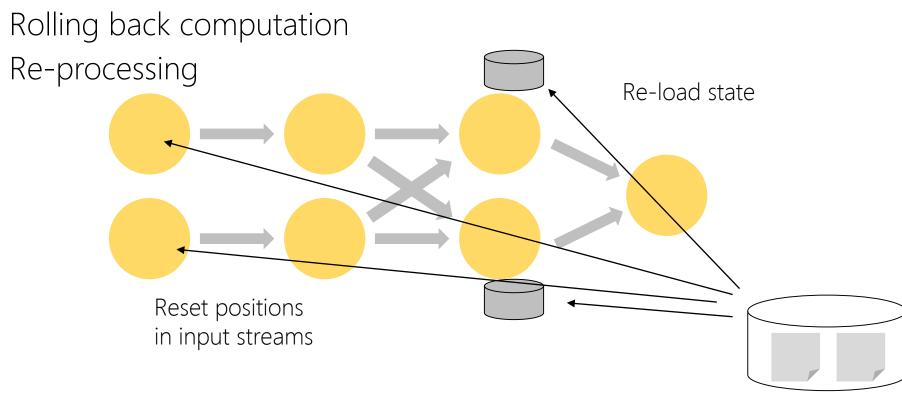


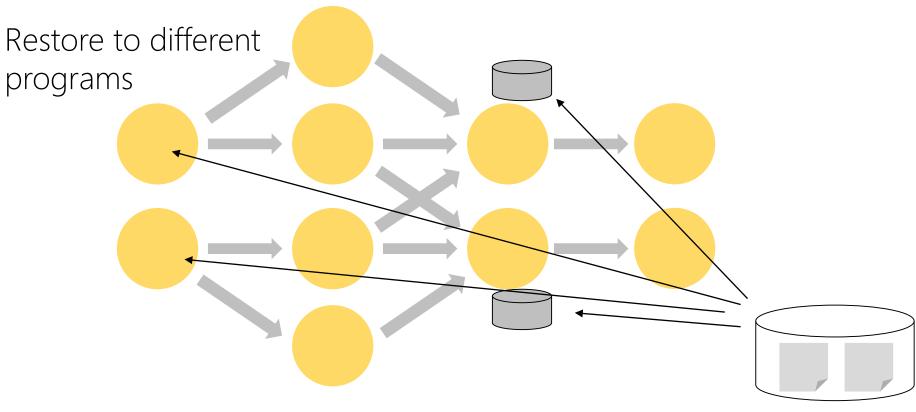














# Checkpoints and Savepoints in Apache Flink





#### What to optimize for?

#### Fast snapshots

Flexible Operations on Snapshots

Checkpoint

Savepoint

## Savepoints: Opt. for Operability

- Self contained: No references to other checkpoints
- Canonical format: Switch between state structures
- Efficiently re-scalable: Indexed by key group
- Future: More self-describing serialization format for to archiving / versioning (like Avro, Thrift, etc.)

## Checkpoints: Opt. for Efficiency

#### Often incremental:

- Snapshot only diff from last snapshot
- Reference older snapshots, compaction over time
- Format specific to state backend:
  - No extra copied or re-encoding
  - Not possible to switch to another state backend between checkpoints
- Compact serialization: Optimized for speed/space, not long term archival and evolution
- Key goups not indexed: Re-distribution may be more expensive



# What else are snapshots / checkpoints good for?

### What users built on checkpoints

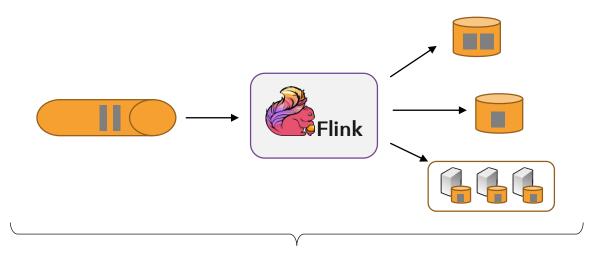
- Upgrades and Rollbacks
- Cross Datacenter Failover
- State Archiving
- State Bootstrapping
- Application Migration
- Spot Instance Region Arbitrage
- A/B testing



## Distributed Snapshots and side effects

#### Transaction coordination for side fx

Snapshots may include side effects



One snapshot can transactionally move data between different systems

#### Transaction coordination for side fx

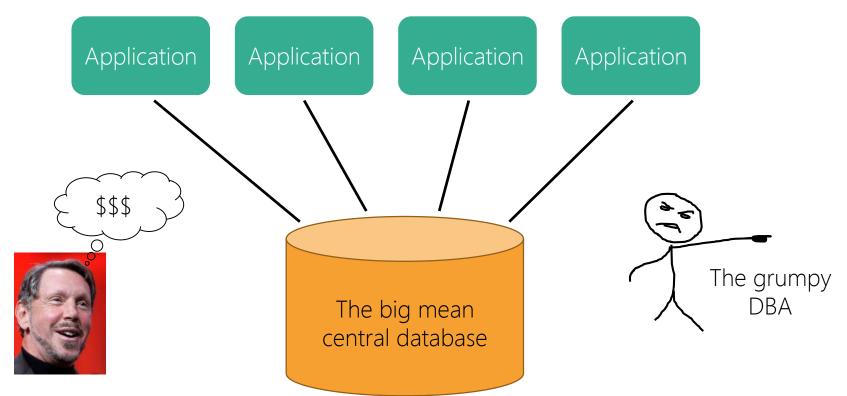
- Similar to a distributed 2-phase commit
- Coordinated by asynchronous checkpoints, no voting delays
- Basic algorithm:
  - Between checkpoints: Produce into transaction or Write Ahead Log
  - On operator snapshot: Flush local transaction *(vote-to-commit)*
  - On checkpoint complete: Commit transactions *(commit)*
  - On recovery: check and commit any pending transactions



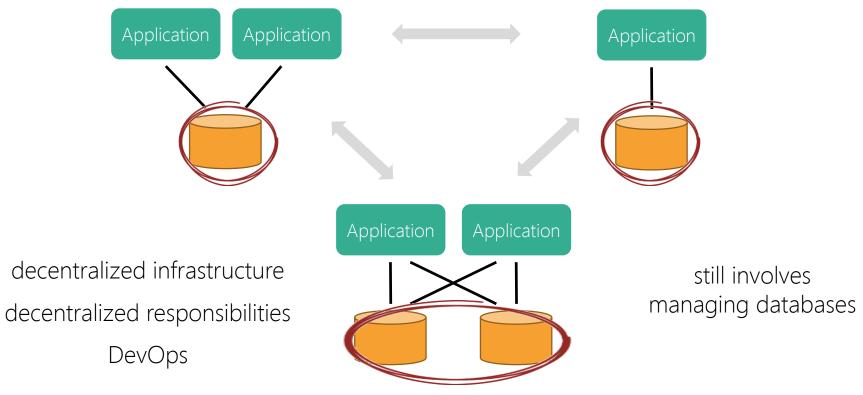
# Distributed Snapshots and Application Architectures

# (A Philosophical Monologue)

#### Good old centralized architecture



### Stateful Stream Proc. & Applications



#### Stateless Application Containers



State management is nasty, let's pretend we don't have to do it

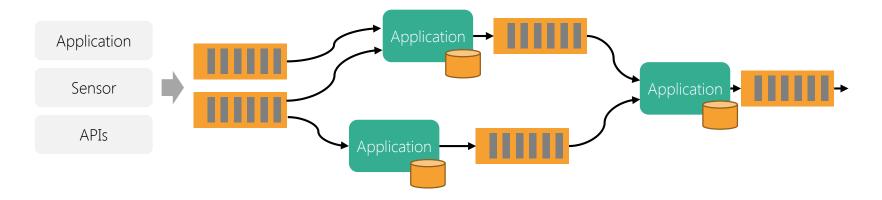
### Stateless Application Containers

Broccoli <del>(state management)</del> is nasty, let's pretend we don't have to eat <del>do</del> it

> Kudos to Kiki Carter for the Broccoli Metaphor

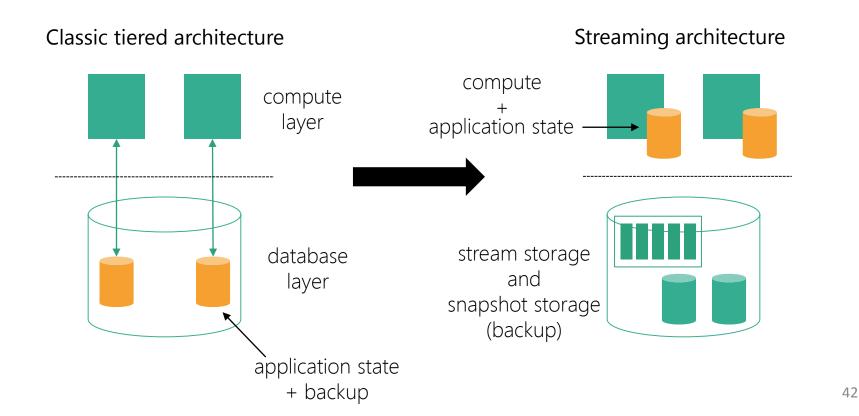
#### Stateful Stream Proc. to the rescue

very simple: state is just part of the application



### Compute, State, and Storage





#### Performance

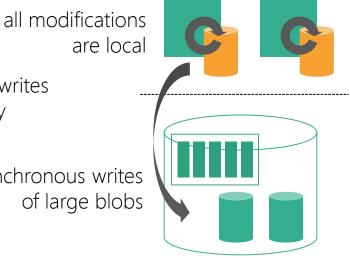
Classic tiered architecture

Streaming architecture

asynchronous writes of large blobs

synchronous reads/writes

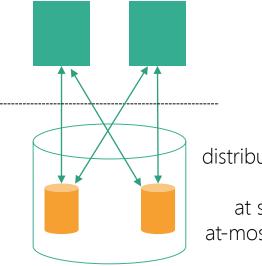
across tier boundary





### Consistency

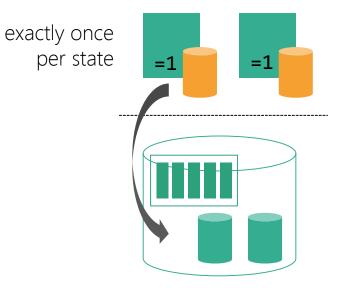
#### Classic tiered architecture



distributed transactions

at scale typically at-most / at-least once

#### Streaming architecture

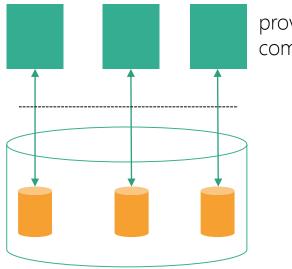




### Scaling a Service



#### Classic tiered architecture

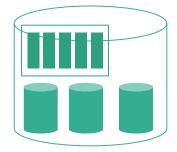


provision compute

> provision compute and state together

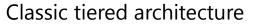
#### Streaming architecture

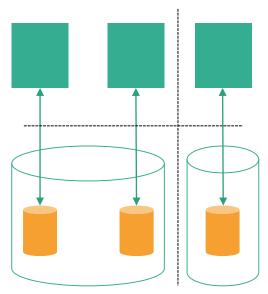




separately provision additional database capacity

### Rolling out a new Service



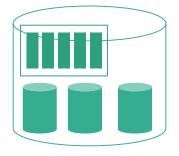


provision a new database (or add capacity to an existing one)

#### Streaming architecture



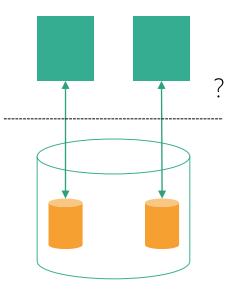
and state together



simply occupies some additional backup space

### Time, Completeness, Out-of-order

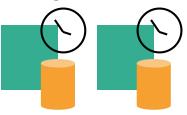
#### Classic tiered architecture

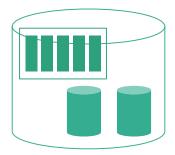


event time clocks define data completeness

> event time timers handle actions for out-of-order data

#### Streaming architecture

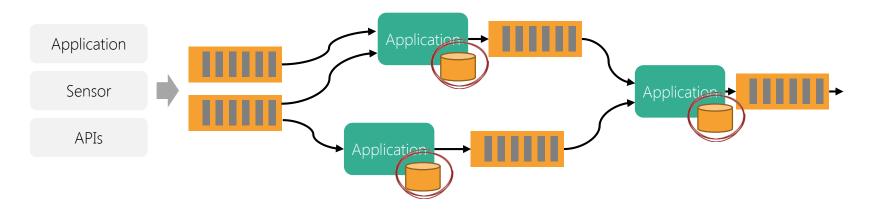




#### Stateful Stream Processing



very simple: state is just part of the application



### The Challenges with that:

- Upgrades are stateful, need consistency
  - application evolution and bug fixes
- Migration of application state
  - cluster migration, A/B testing
- Re-processing and reinstatement
  - fix corrupt results, bootstrap new applications
- State evolution (schema evolution)



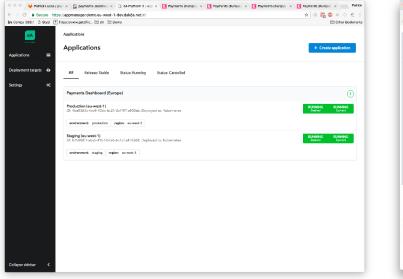
#### The answer

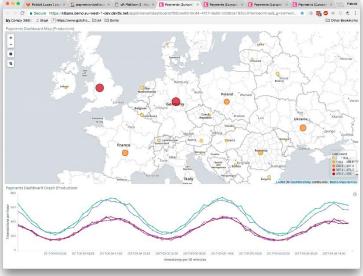
(my personal and obviously biased take)

### Consistent Distributed Snapshots



### Demo Time!





#### Payments Dashboard



# Thank you very much 😊

#### (shameless plug)



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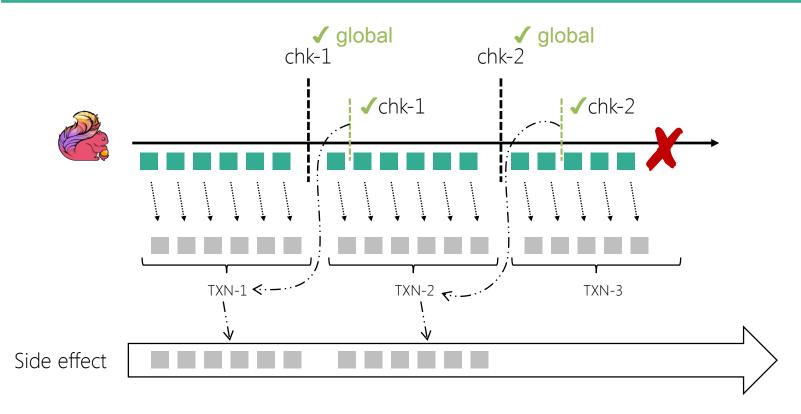
## Appendix



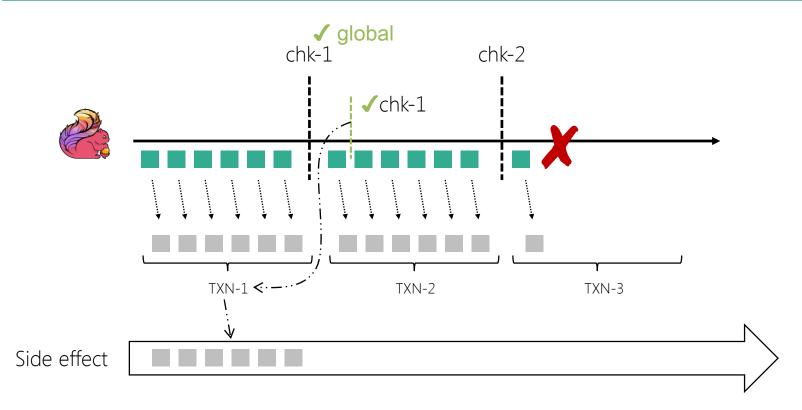
# Details about Snapshots and Transactional Side Effects

#### Exactly-once via Transactions

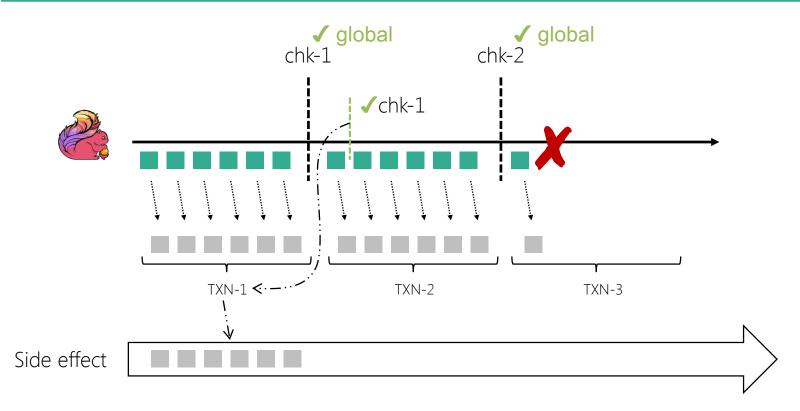




### Transaction fails after local snapshot



### Transaction fails before commit...



#### ... commit on recovery

