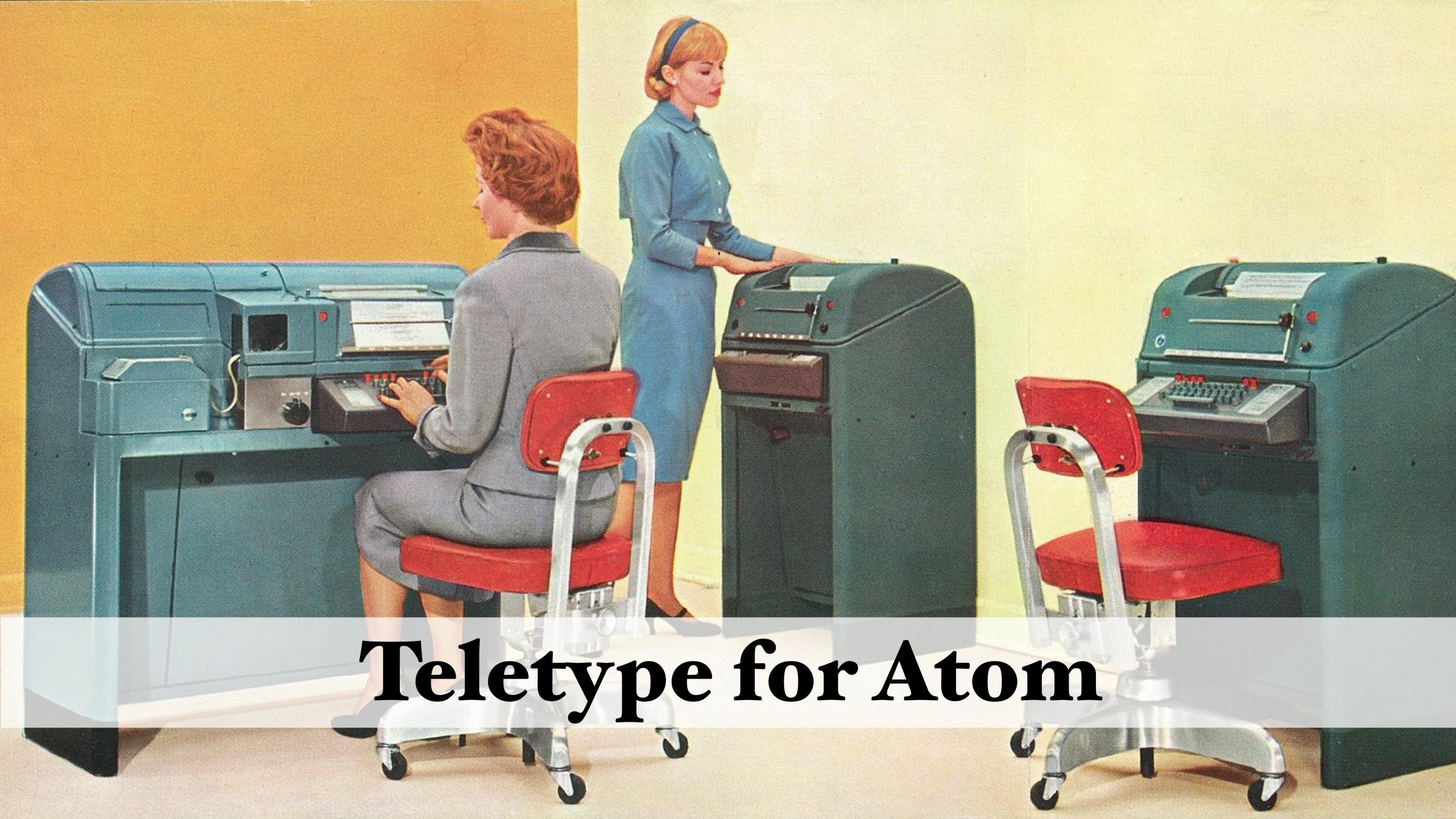
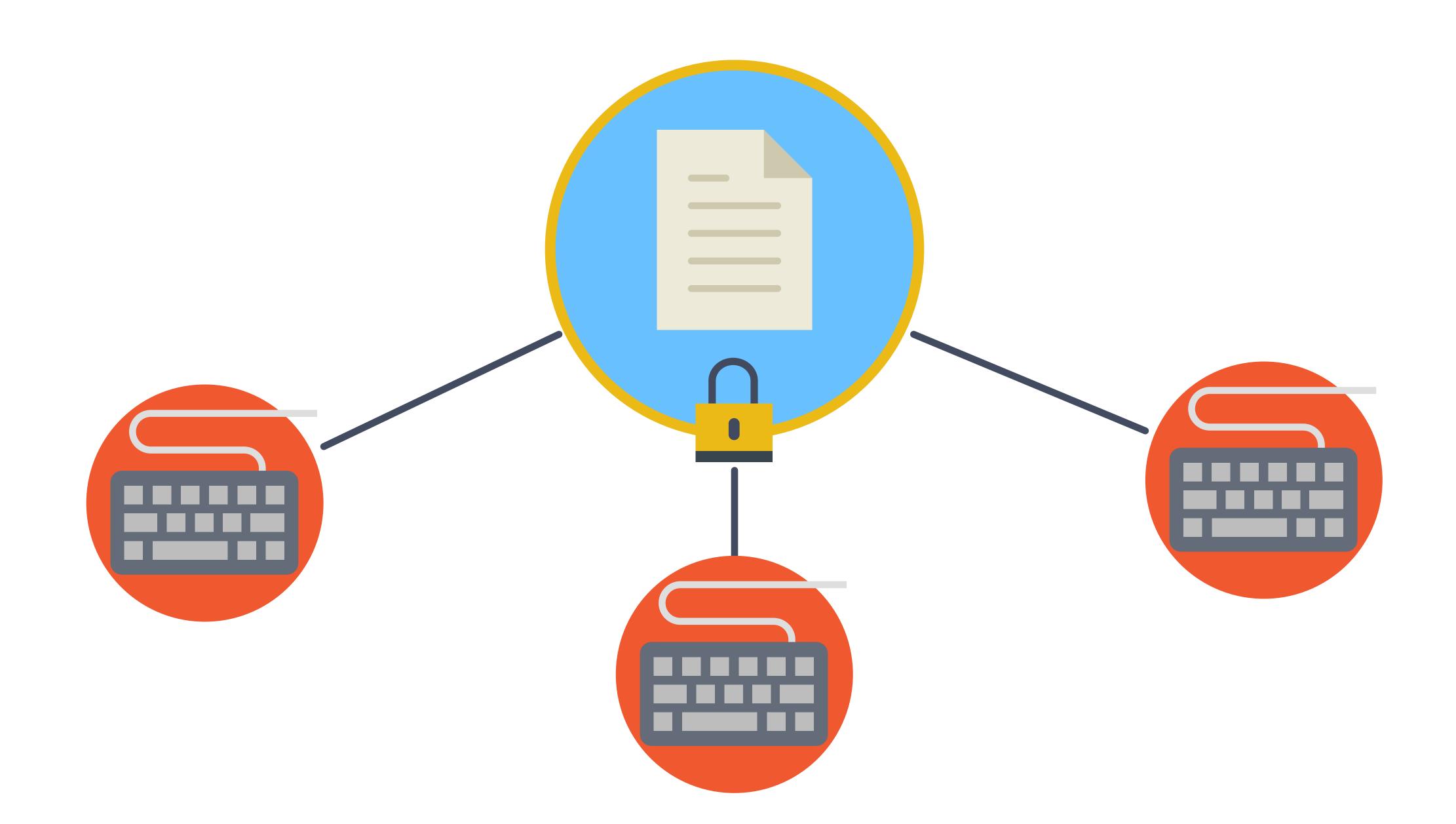
Real-Time Collaborative Editing With CRDTs

Real-Time Collaborative Editing With CRDTs

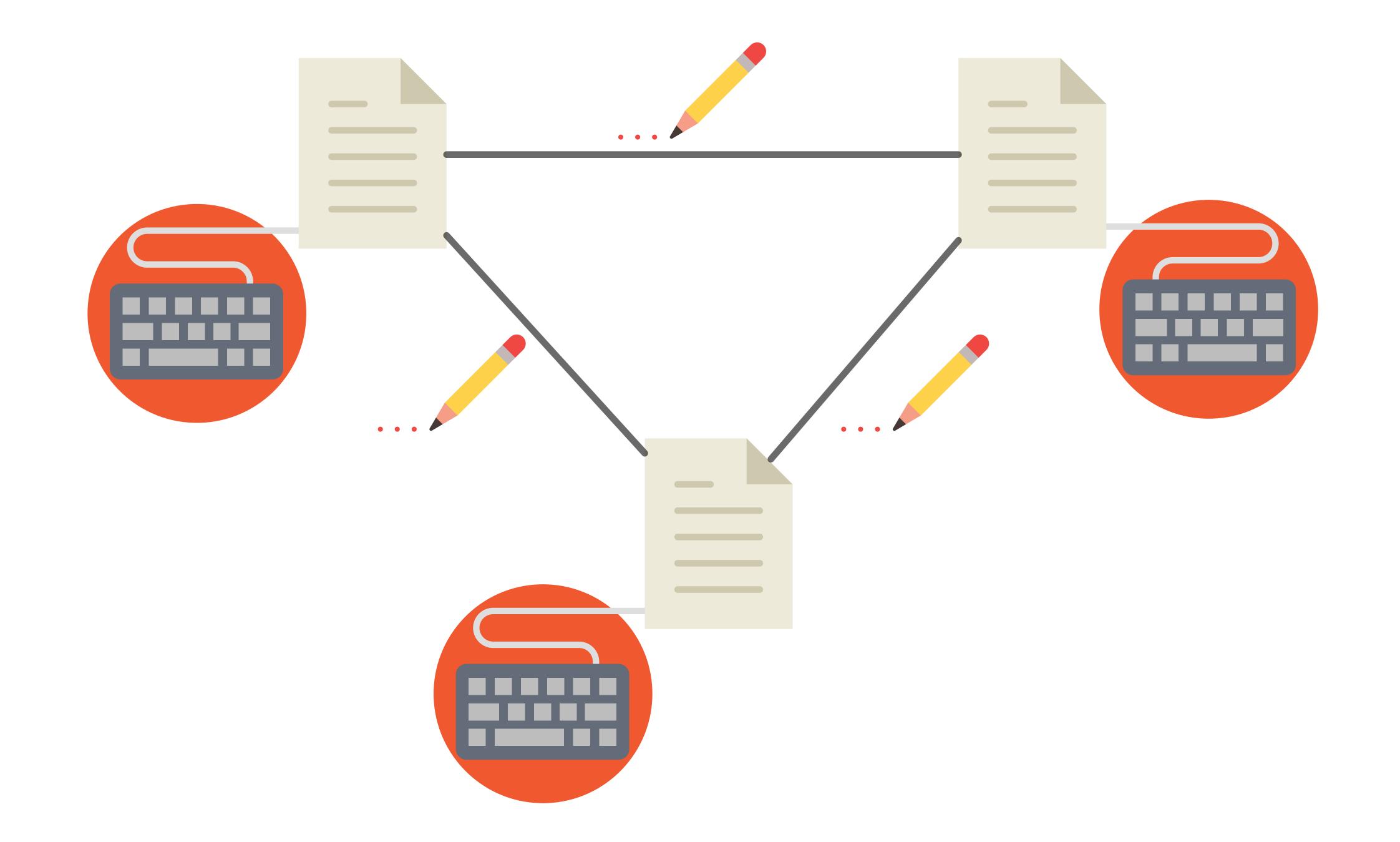


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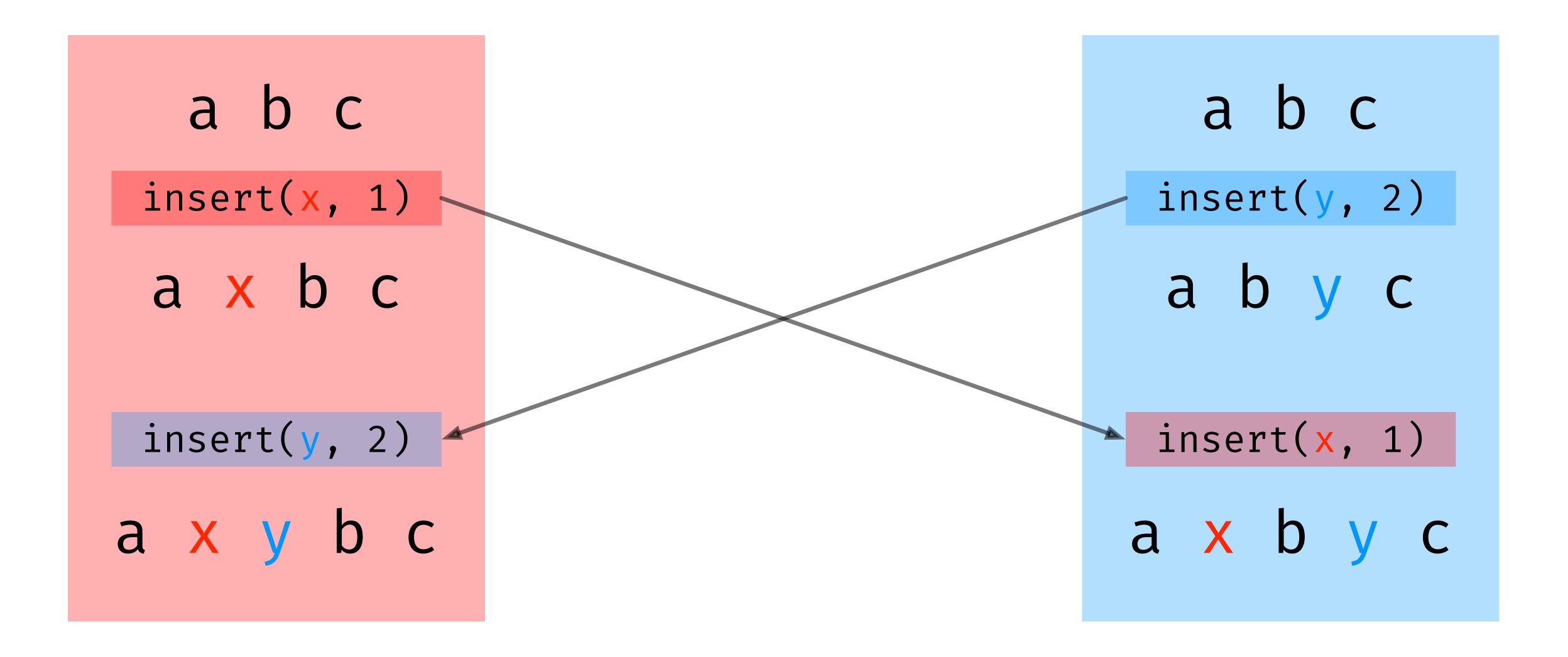
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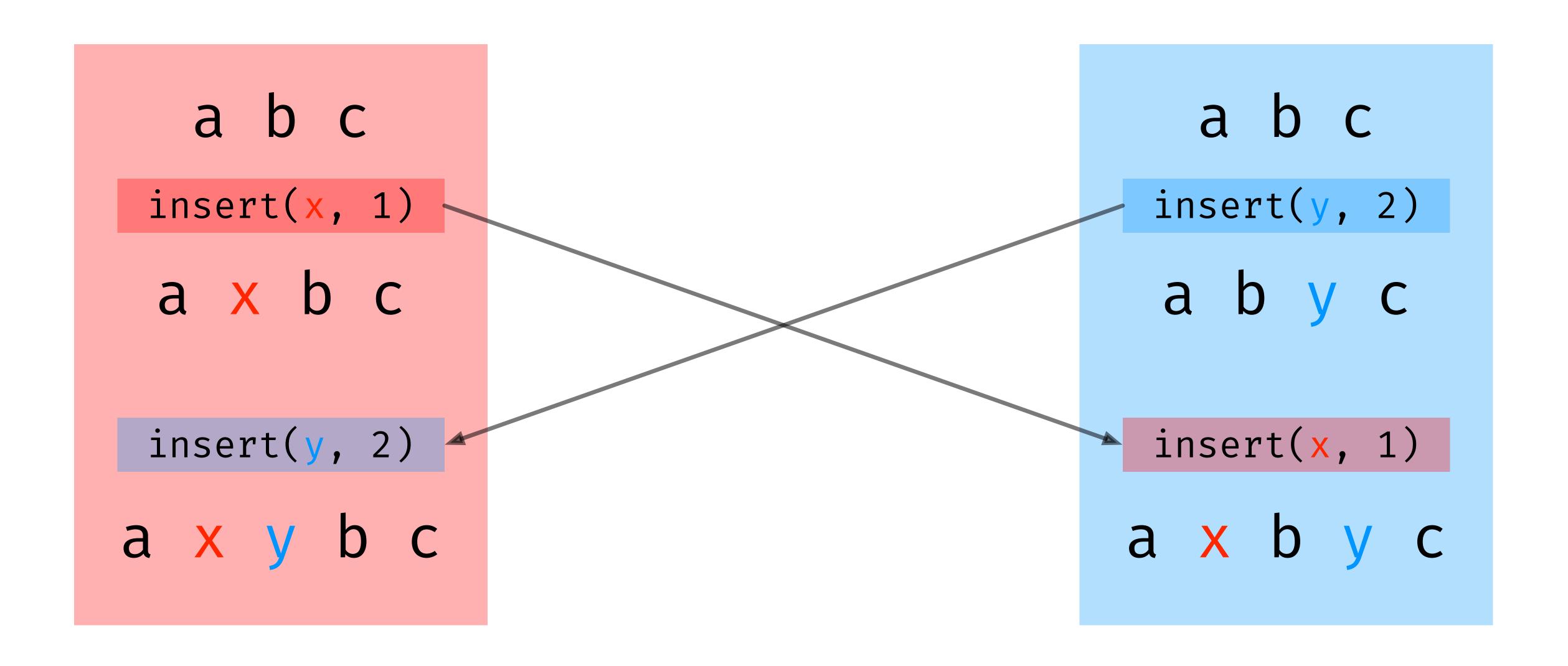
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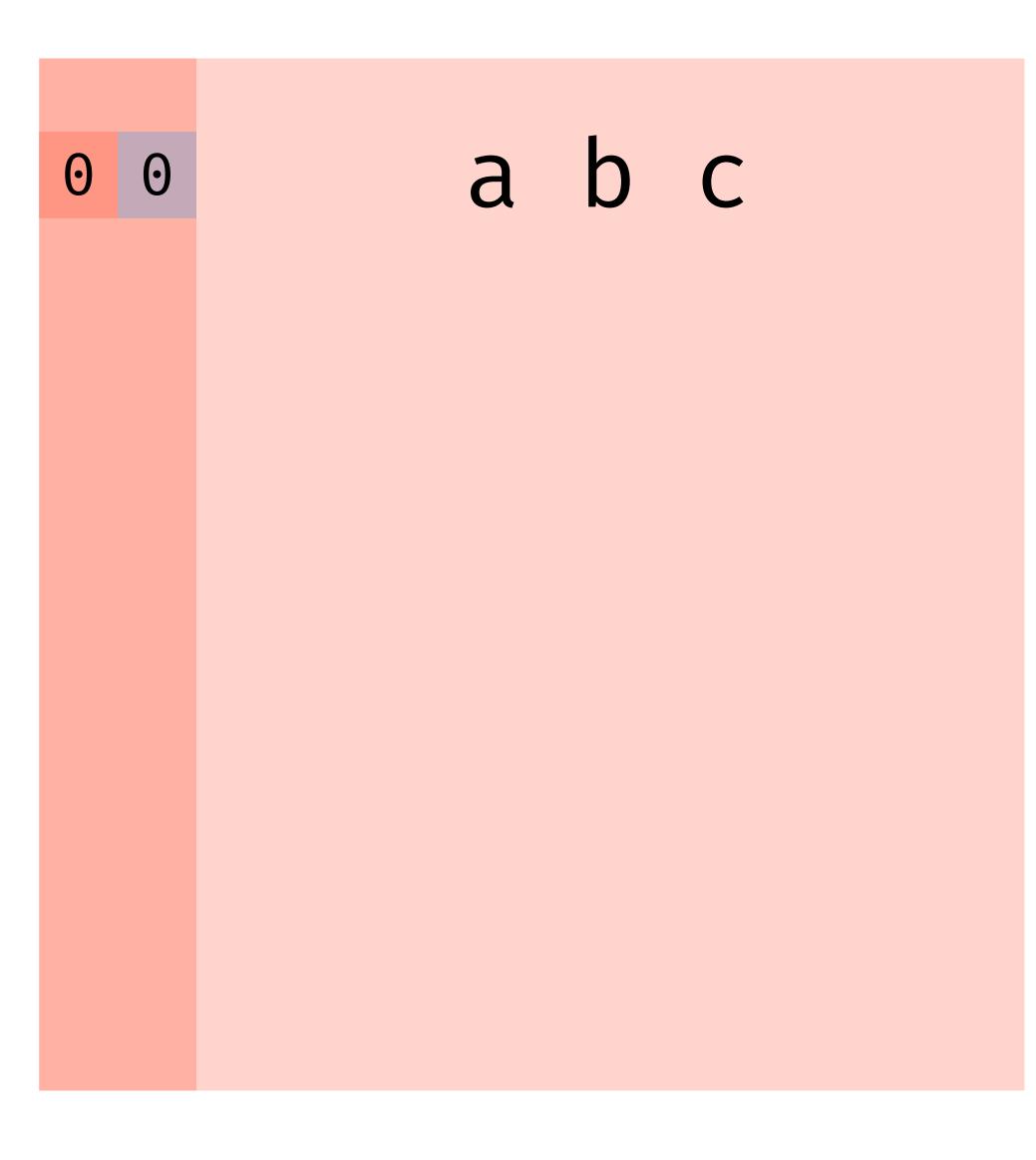
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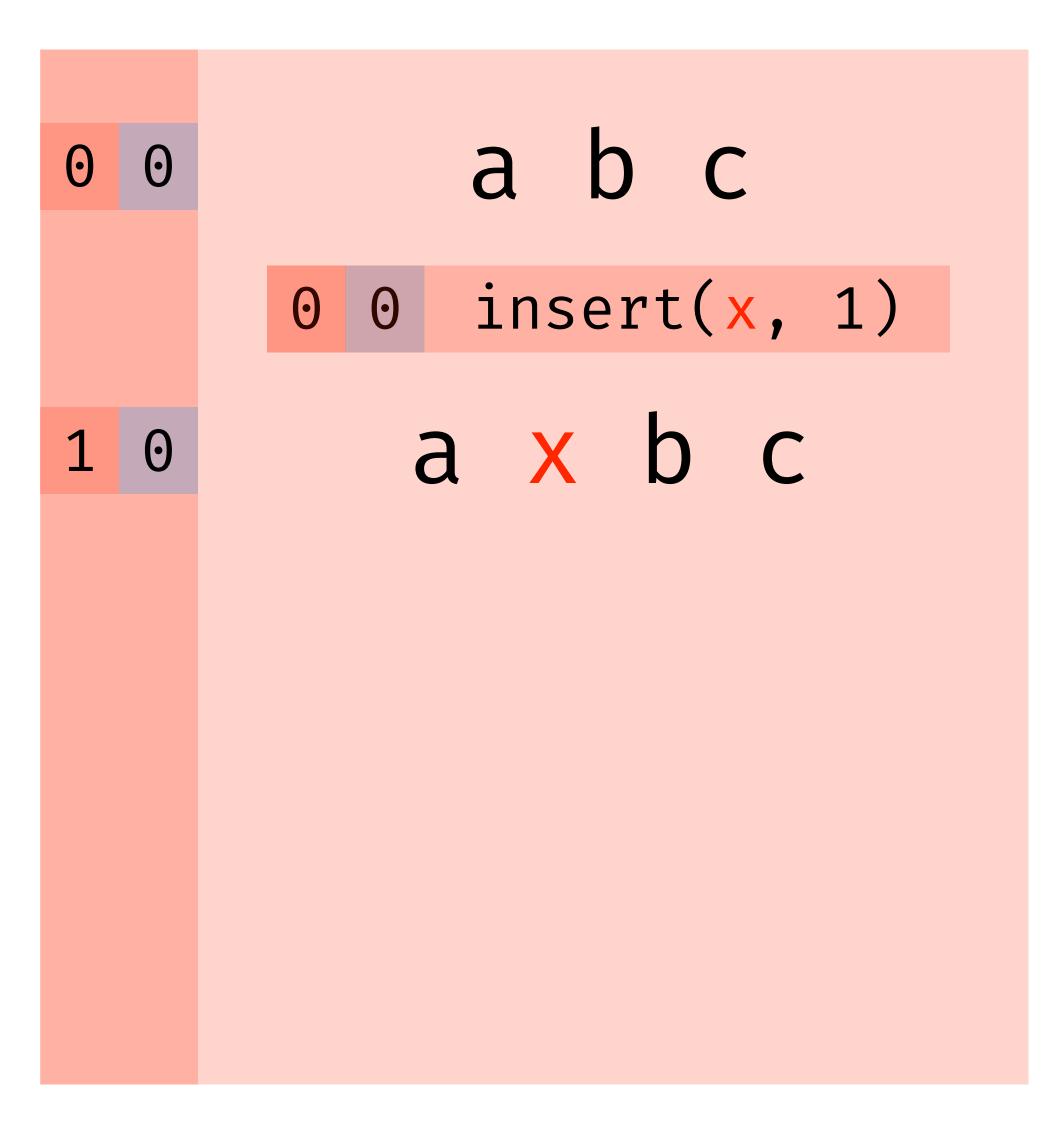
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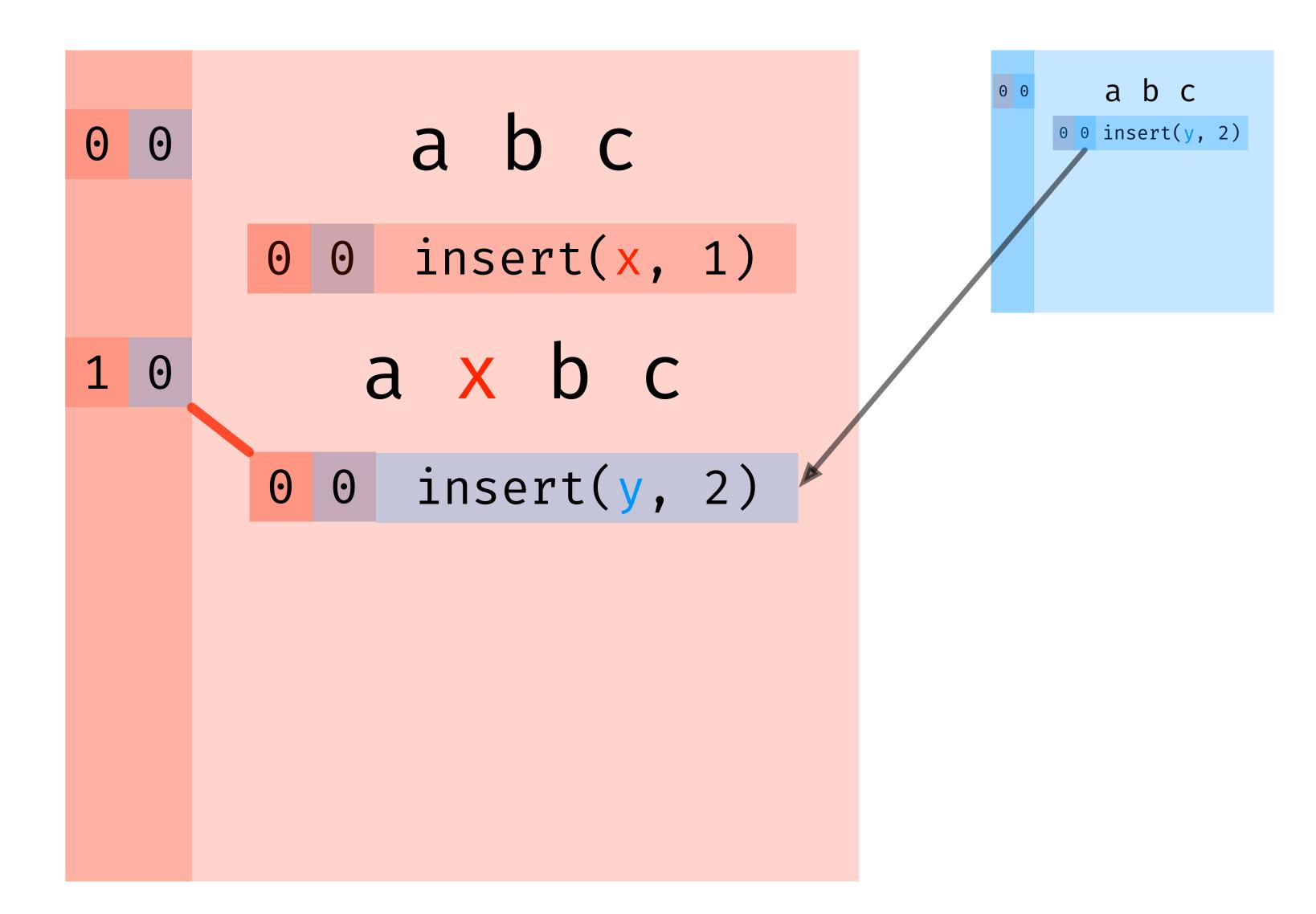


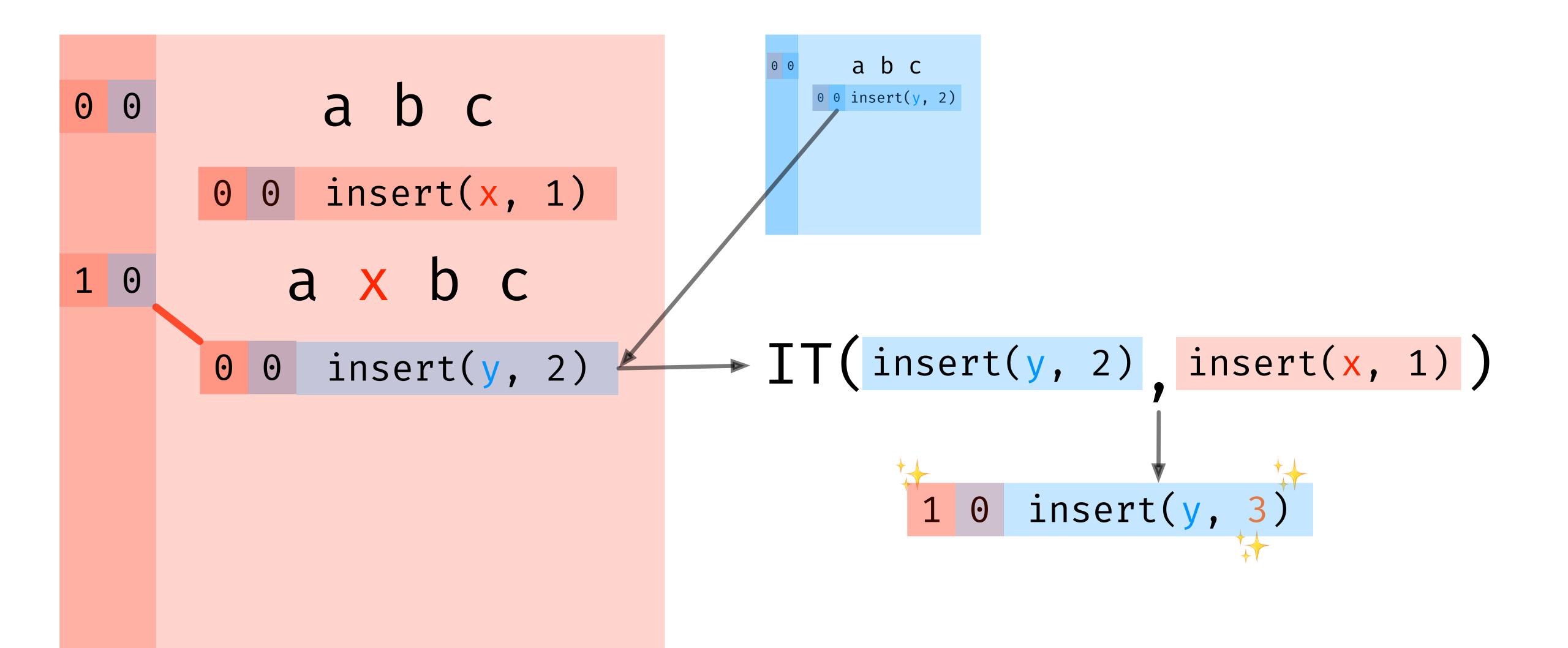


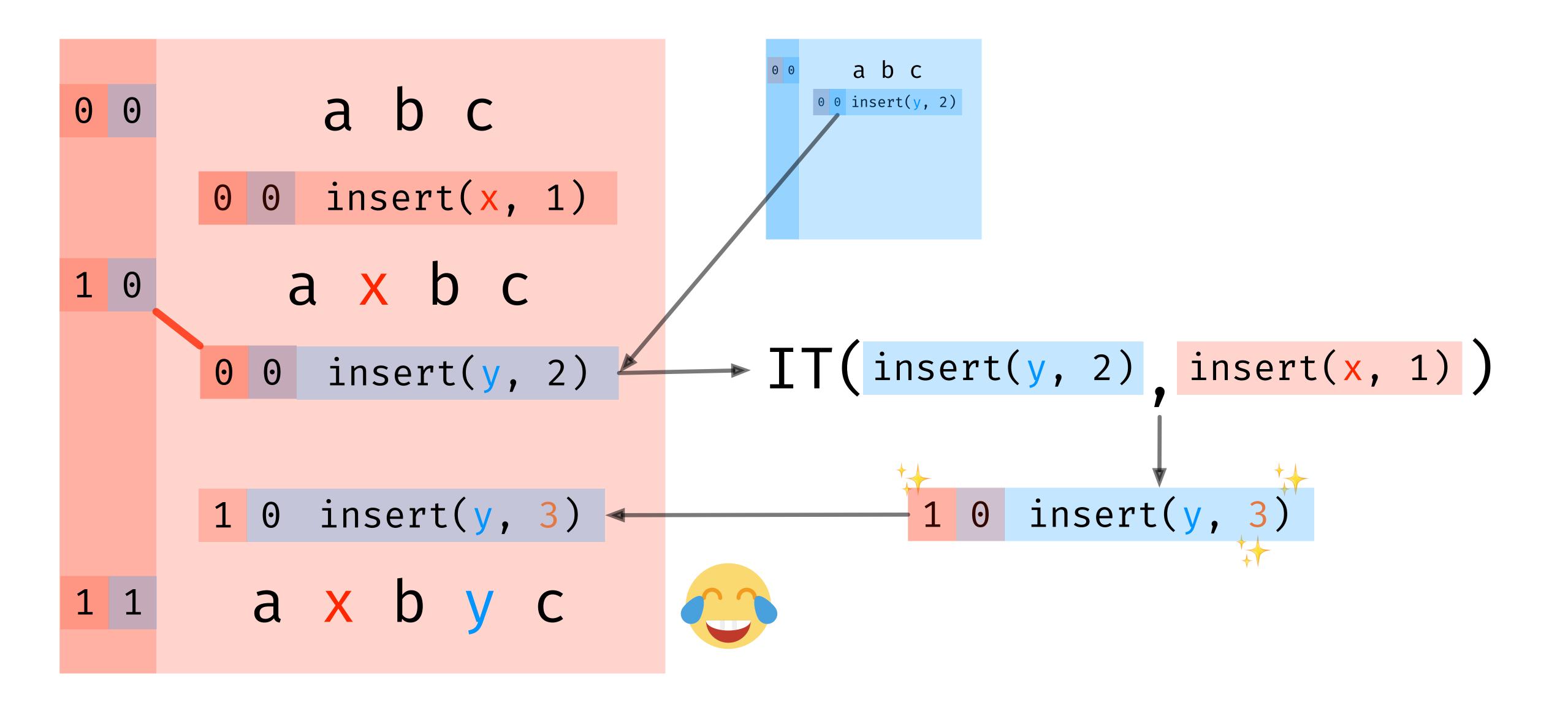
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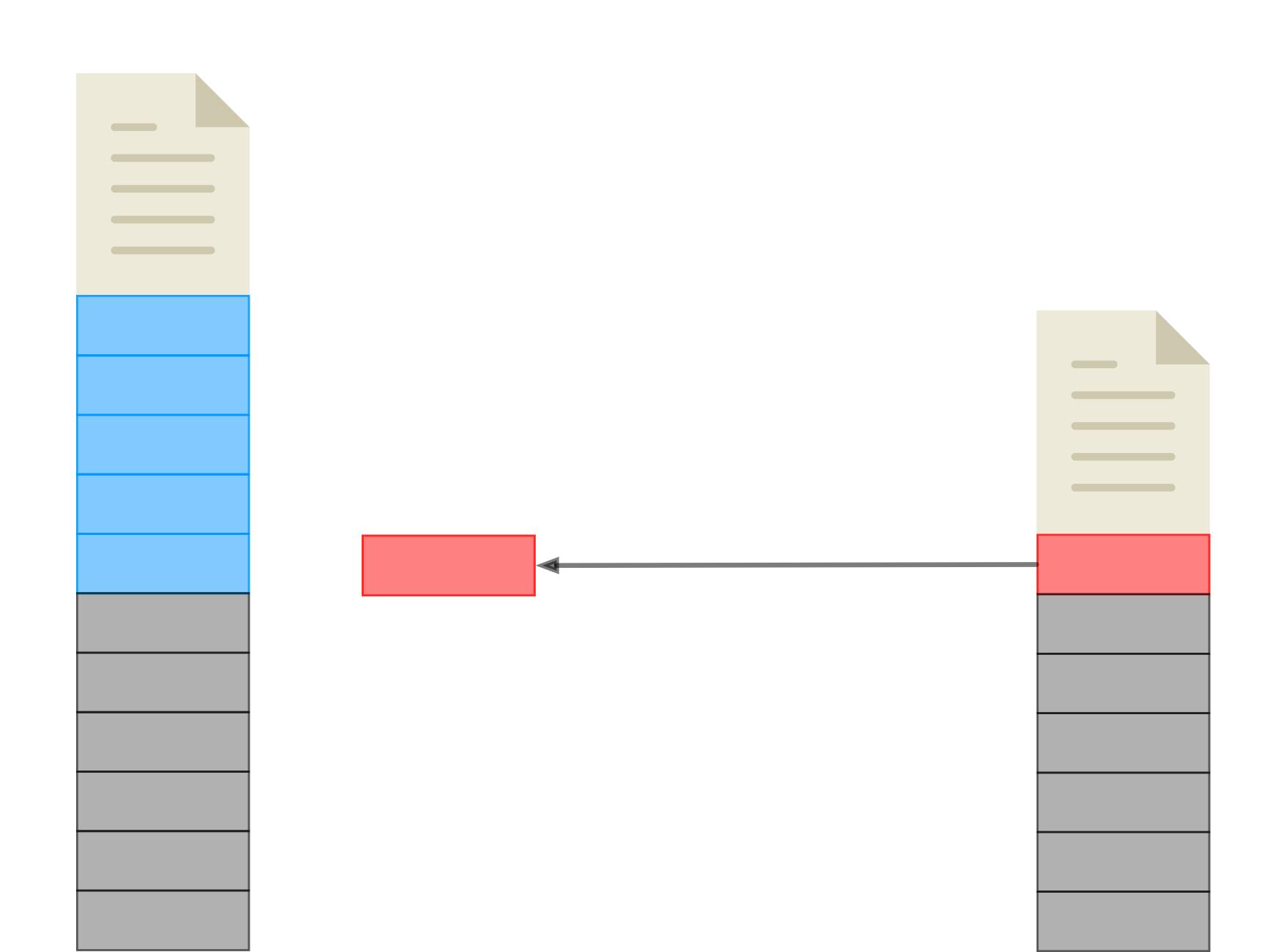


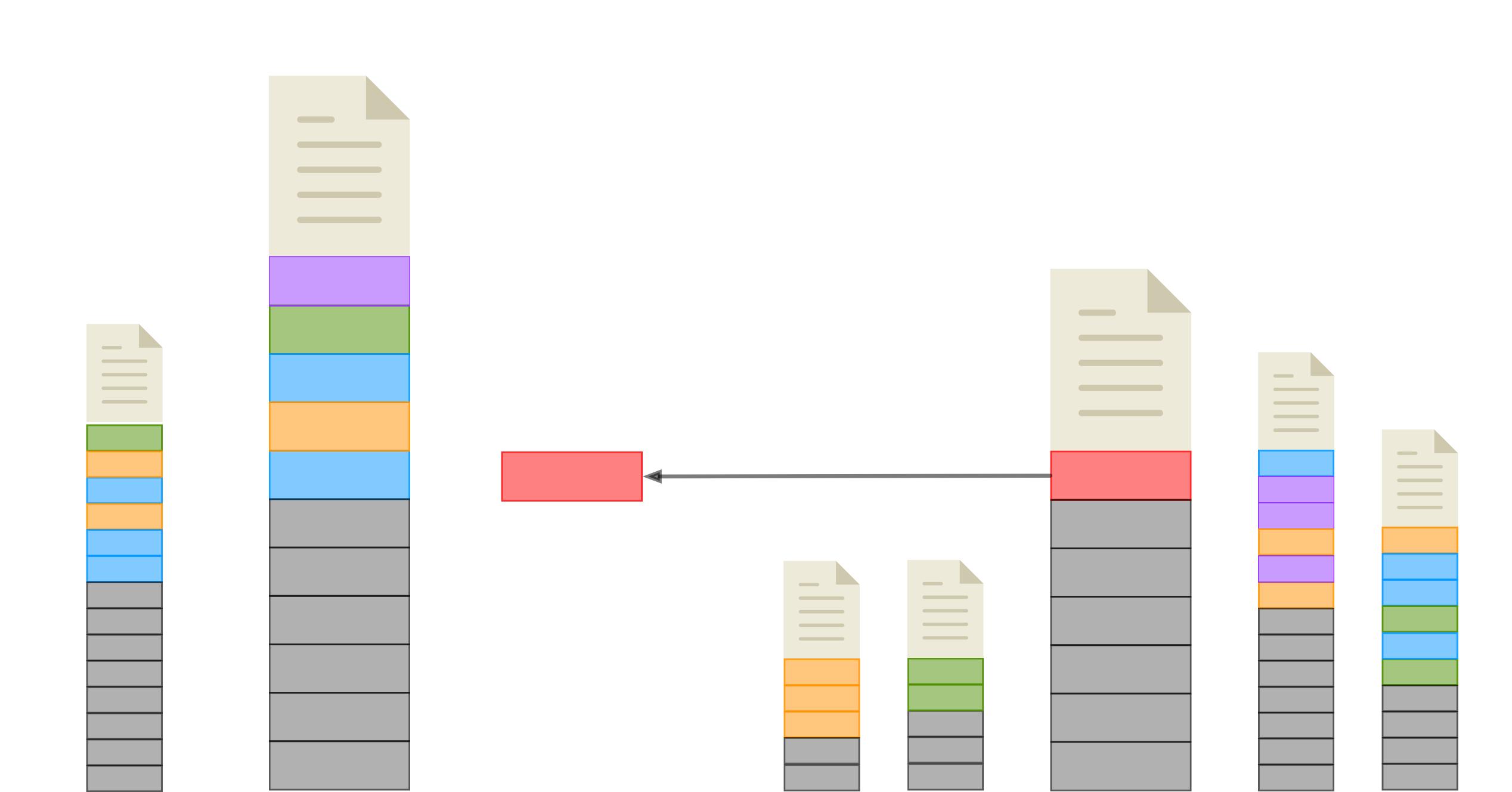












Concurrent Undo Oper Environments Usi-

Transparent Adaptation of Single-User Real-Time

Operation Context and Context-based Operational Transformation

David Sun Computer Science Division, EECS University of California Berkeley, CA

davidsun@cs.berkeley.edu

ABSTRACT

Operational Transformation

'T) is a technique for consisundo, and is being applied aborative applications. The is crucial in determining its l new problems, as well as The theory of causality has OT systems, but it is inadectness requirements. Past atches to work around this y intricate and complicated esigned, implemented, and T algorithms, we reflected

Chengzheng Sun School of Computer Engineering Nanyang Technological University Singapore CZSun@ntu.edu.sg

1. INTRODUCTION

Operational Transformation (OT) was originally invented for consistency maintenance in plain-text group editors [4]. In over 15 years, OT has evolved to support an increasing number of applications, including group undo [15, 19, 18, 21], group-awareness [28], operation notification and compression [20], spreadsheet and table-centric applications [14, 27], HTML/XML and tree-structured document editing [3, 7], word processing and slide creation [29, 25, 24], transparent and heterogenous application-sharing [1, 10, 24], and mobile replicated computing and database systems [6, 16].

To effectively and efficiently support existing and new applications, we must continue to improve the capability and

and new problems. The ation for OT is crucial in pinning of all existing OT y [9, 17, 4, 22]: causally ed in their causal order; nsformed before their exausality is inadequate to r correct transformatic

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> City University of Hong Kong Kowloon, Hong Ko

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Tombstone Transformation Functions for Ensuring Editing Systems Consistency in Col

Gérald Oster Institute for Information Systems ETH Zurich Email: osterg@inf.ethz.ch

Abstract-In collaborative editing, consistency ma of the copies of shared data is a critical issue. In decade, Operational Transformation (OT) approach re a suitable mechanism for maintaining consistency. Unfo none of the published propositions relying on this app able to satisfy the mandatory correctness properties TP_2 defined in the Ressel's framework. This paper this correctness issue by proposing a new way to mo state by retaining tombstones when elements are rei instantiation of the proposed model for a linear data and the related transformation functions are provide

I. Introduction

Collaborative editing systems allow users to edi document from multiple sites across Internet. De the work context, users can work synchronously chronously. Synchronous collaboration is also calle editing since when a user performs some modif the document, these modifications are instantly s users who can see them without any delay. In t in asynchronous collaboration, users may not same time. They work in isolation: they can dec publish their modifications and when to integrate n performed by other users.

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Achieving Convergence, Causality Preservation, and Intention Preservation in Real-Time Cooperative Editing Systems

CHENGZHENG SUN Griffith University **XIAOHUA JIA** City University of Hong Kong YANCHUN ZHANG The University of Southern Queensland YUN YANG Deakin University and DAVID CHEN Griffith University

Real-time cooperative editing systems allow multiple users to view and edit the same text/graphic/image/multimedia document at the same time from multiple sites connected by communication networks. Consistency maintenance is one of the most significant challenges in designing and implementing real-time cooperative editing systems. In this article, a consistency model, with properties of convergence, causality preservation, and intention preservation, is proposed as a framework for consistency maintenance in real-time cooperative editing systems. Moreover, an integrated set of schemes and algorithms, which support the proposed consistency model, are devised and discussed in detail. In particular, we have contributed (1) a novel generic operation transformation control algorithm for achieving

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ENGZHENG SUN

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fith University

Du Li Nokia Research Center Palo Alto, CA, USA lidu008@gmail.com ience

School of Computer Science Fudan University Shanghai, China ninggu@fudan.edu.cn

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1. INTRODUCTION

Undo is a key feature in interactive applications. Many familiar single-user applications, including text editors, word nnnar single-user applications, including text editors, word processors, design tools and even Web browsers, allow the user to undo operations in a chronological order. Undo is user to undo operations in a chronological order. Ondo is often used for user-level error recovery, e.g., to fix typos. It can also encourage users to explore unfamiliar capabilities in an application provided that the effects of erroneous

In collaborative applications, undo is significantly more operations can always be undone. that conadorative applications, undo is significantly more challenging [9, 2, 14]. When the users are distributed and work in parallel, local and remote operations could be inwork in paramet, local and remote operations could be in-terleaved arbitrarily due to concurrency. A user request to terleaved arbitrarily due to concurrency. A user request to undo could mean to undo a local operation or a remote operation. To undo a remote operation is even trickier because there can be multiple operations by several remote users and there can be multiple operations by several remote users and it is important to unambiguously specify which operation by which user. Otherwise, the undo action could be interby which user. Otherwise, the undo action could be interpreted arbitrarily at different sites, leading to unpredictable consequences in the system. Therefore, in a distributed enconsequences in the system. Therefore, in a distributed environment, it becomes necessary to support selective undo

Department of Computer Science

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an important mechanism for error recovery and exploration of alternatives in interactive and

llaborative applications, an undo facility should have the capability of undoing any operation at

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Tombstone Transformation Functions for Ensuring Editing Systems

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"Unfortunately, implementing OT sucks. There's a million algorithms with different tradeoffs, mostly trapped in academic papers. The algorithms are really

hard and time consuming to implement correctly. Wave

took two years to write and if we rewrote it today, it

would take almost as long to write a second time."

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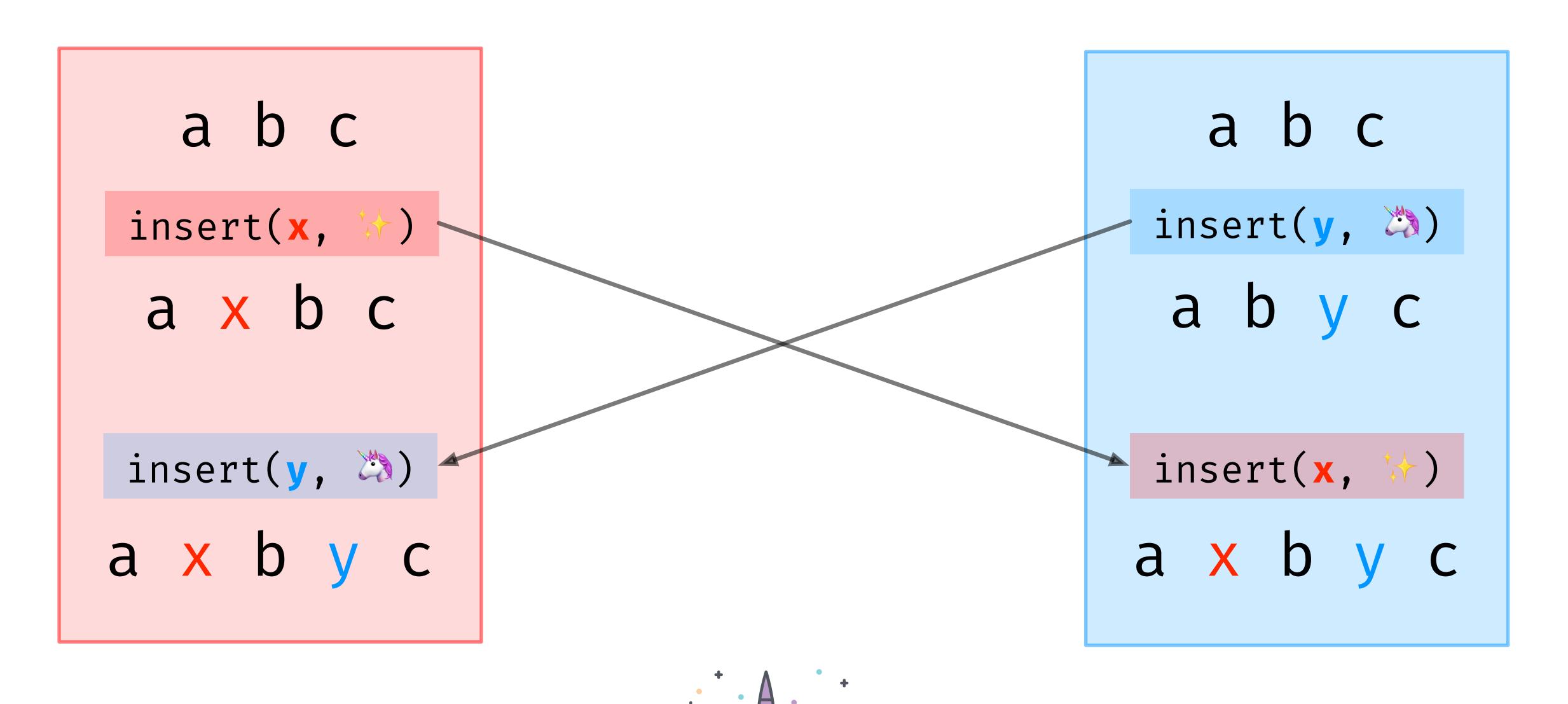
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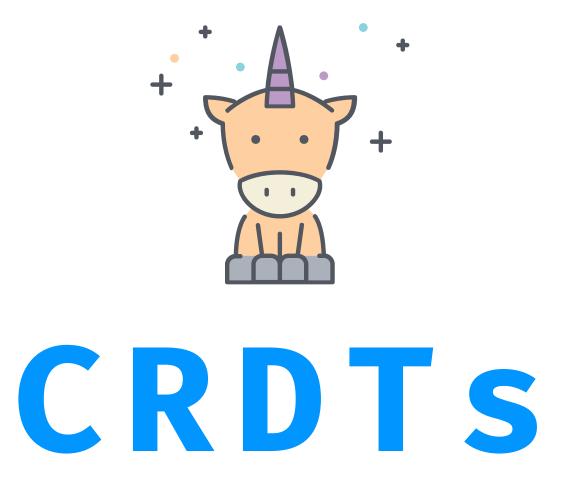
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Conflict-Free Replicated Data Types

Data Consistency for P2P Collaborative Editing

Gérald Oster
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Pascal Urso, Pascal Molli, Abdessamad Imine Université Henri Poincaré, Nancy 1 LORIA {urso,molli,imine}@loria.fr

ABSTRACT

Peer-to-peer (P2P) networks are very efficient for distributing content. We want to use this potential to allow not only distribution but collaborative editing of this content. Existing collaborative editing systems are centralised or depend on the number of sites. Such systems cannot scale when deployed on P2P networks. In this paper, we propose a new model for building a collaborative editing system. This model is fully decentralised and does not depend on the number of sites.

Categories and Subject Descriptors

C.2.4 [Computer-Communication Networks]: Distributed Systems—Distributed applications; H.5.3 [Information Interfaces and Presentation]: Group and Organization Interfaces—Collaborative computing, theory and models

General Terms

Algorithms, Design, Human Factors

Keywords

CSCW, Collaborative editing, Optimistic replication, Concurrency control

1. INTRODUCTION

Currently, peer-to-peer systems demonstrated how they can ensure scalable content distribution. In their survey [4], Androutsellis-Theotokis et al. wrote:

"Peer-to-peer content distribution systems rely on the replication of content on more than one node for improving the availability of content, enhancing performance and resisting censorship attempts."

We want to reuse these characteristics not only for content distribution but also for content editing. Currently, P2P

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CSCW'06, November 4–8, 2006, Banff, Alberta, Canada. Copyright 2006 ACM 1-59593-249-6/06/0011 ...\$5.00.

networks mainly distribute immutable contents, we want to distribute updates on this content and manage collaborative editing on it. We are convinced that if we can deploy a group editor framework on a P2P network, we open the way for P2P content editing. It means that all existing collaborative editing applications such as CVS and Wiki can be redeployed on P2P networks and take advantage of availability improvements, performance enhancements and censorship resistance of P2P networks. For instance, Wikipedia [2] is currently a collaborative encyclopædia that has collected more than 4,700,000 articles in more than 200 languages. Wikipedia has more than 50 million of page requests per day. 200,000 changes are made every day [1]. However, Wikipedia needs a costly infrastructure to handle the load. Hundreds of thousands of dollars are spent every year to fund the infrastructure. A P2P massive collaborative editing system would allow to distribute the service, tolerate failures, improve performances, resist to censorship and share the cost of the underlying infrastructure.

Collaborative editing systems such as CVS or Wikis are currently centralised and cannot be adapted to peer-to-peer networks. Collaborative systems based on the operational transformation approach [7, 23] can be decentralised.

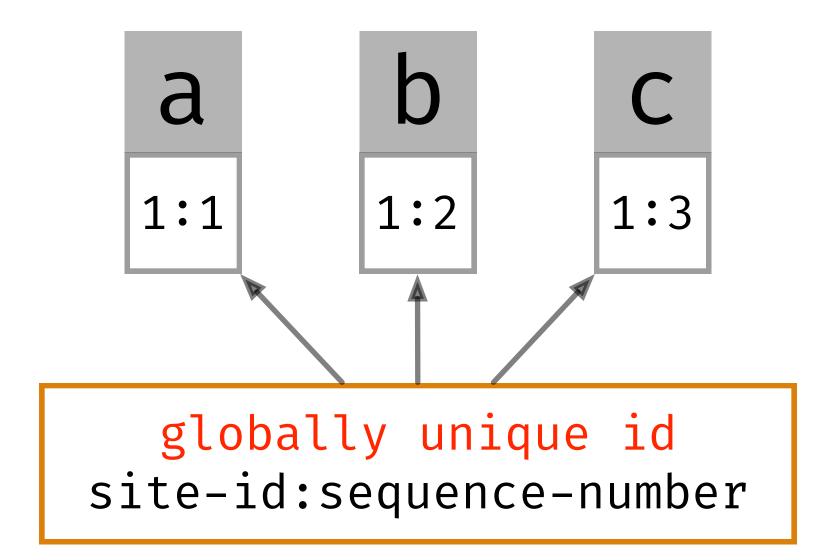
However, existing algorithms such as GOTO [23], ABT[16] and SOCT2 [21] rely on vector clocks to detect concurrent operations. OT approach supposes that each operation is immediately executed locally, stored in a local log and then broadcast to other sites in order to be re-executed and stored in their logs. A vector clocks is associated to each operation. A vector clocks [17] is an array of logical clocks, one clock per site. It is used to detect the happened-before relationship and therefore the concurrency between operation. It causes no problem if the number of sites is fixed and low but if the number of sites grows, the size of the vector clocks is unbounded. Thus, messages exchanged between sites will grow as well as the size of local operation. Also, the time efficiency of operation on vectors clocks will decline as vectors clocks grow. Clearly, vectors clocks prevent these algorithms to scale and represent a serious bottleneck for their deployment on P2P networks.

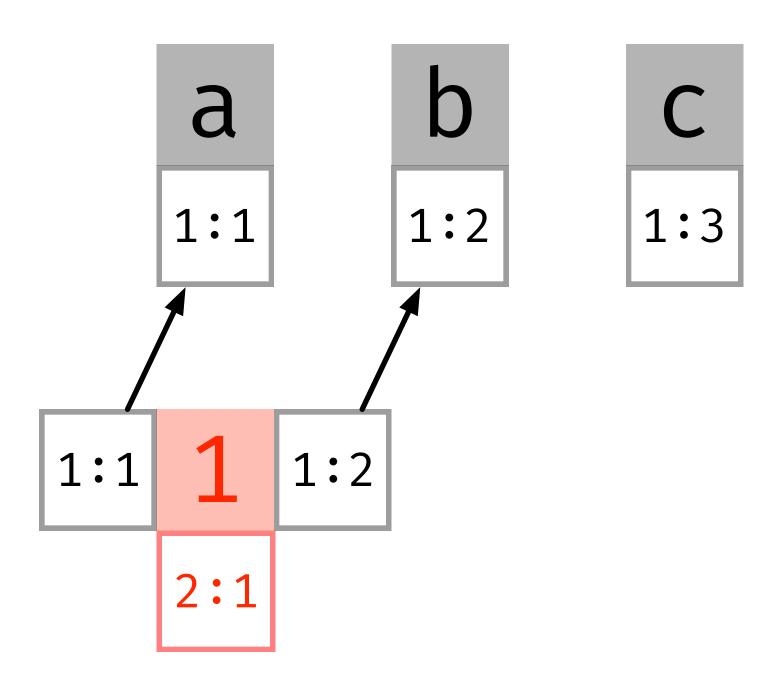
In this paper we propose a new model called WOOT for building group editors that is suitable for dynamic P2P systems. Compared to existing decentralised group editor models, the number of sites involved in group editing is not a variable

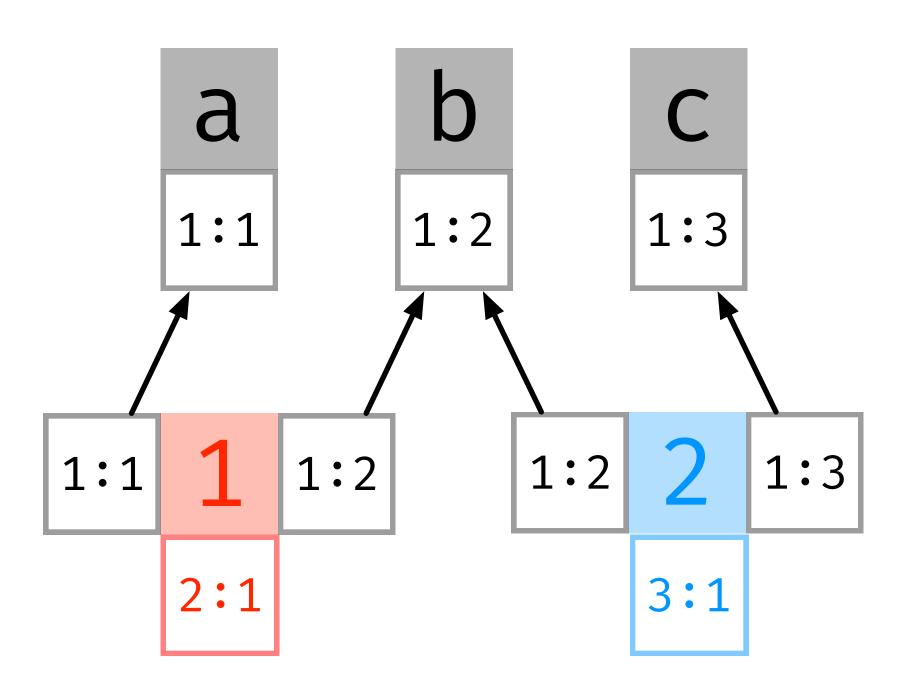
The remainder of this paper is organised as follows: Section 2 describes the WOOT approach and details the consistency model. Section 3 presents a formal definition of

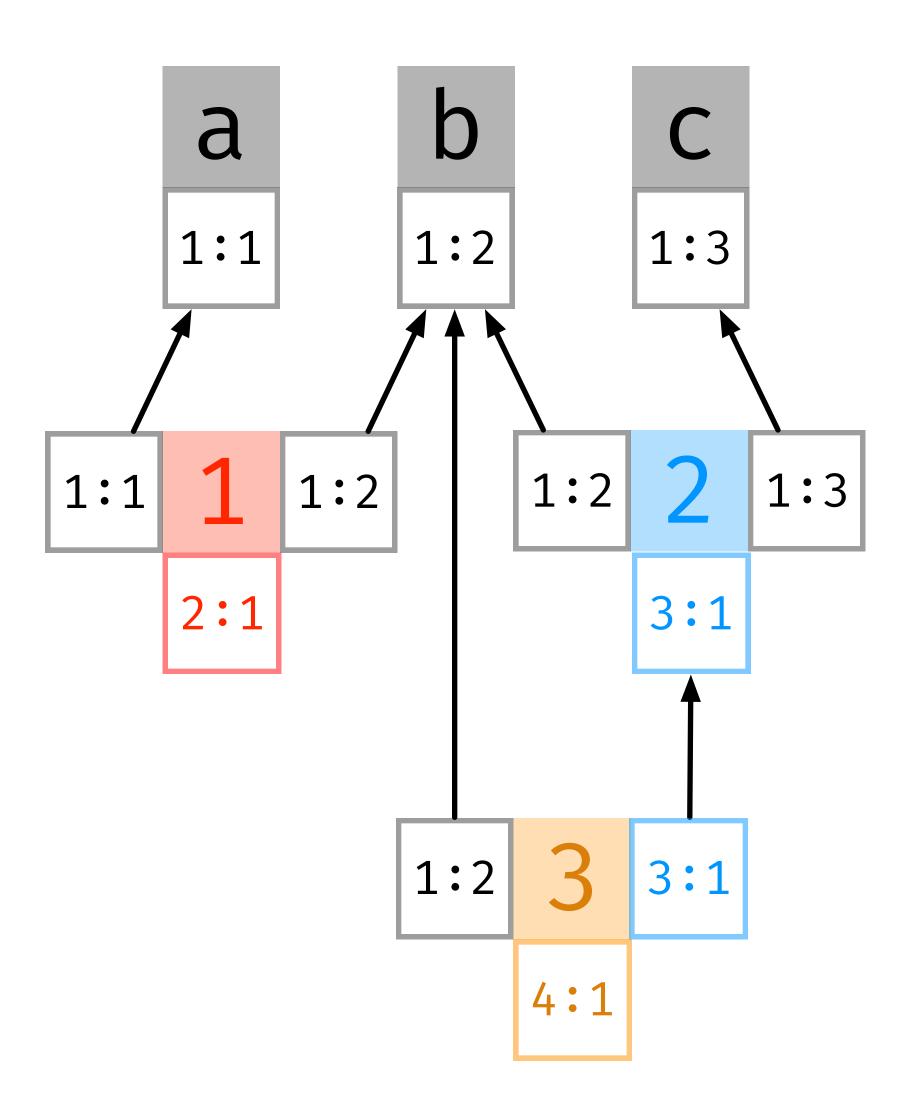
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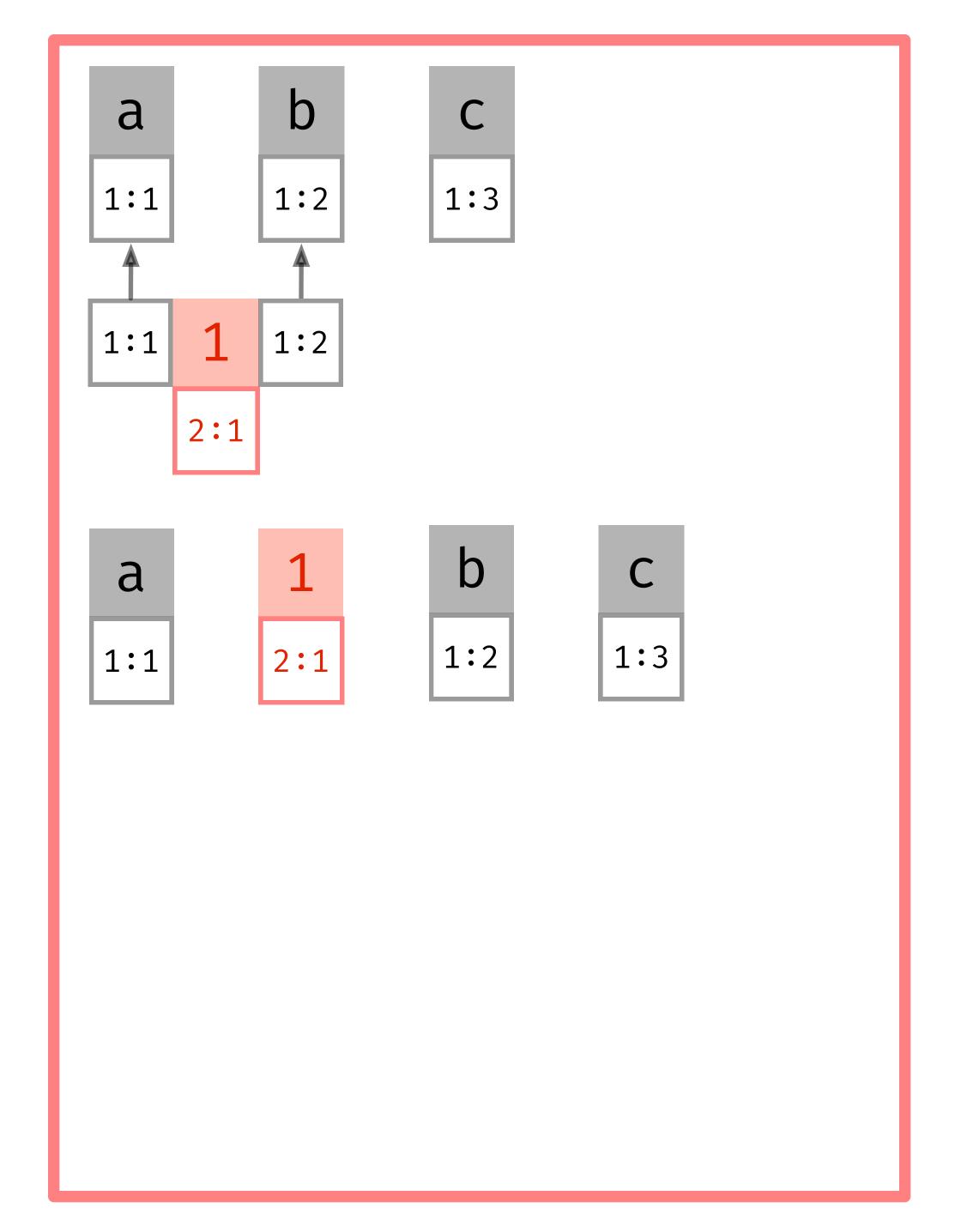


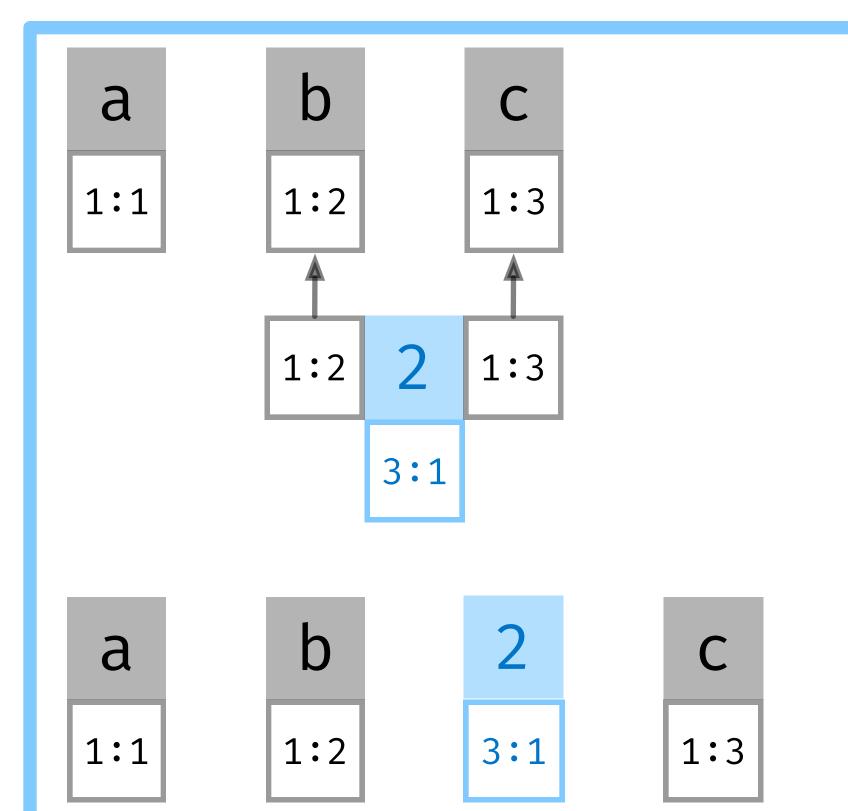


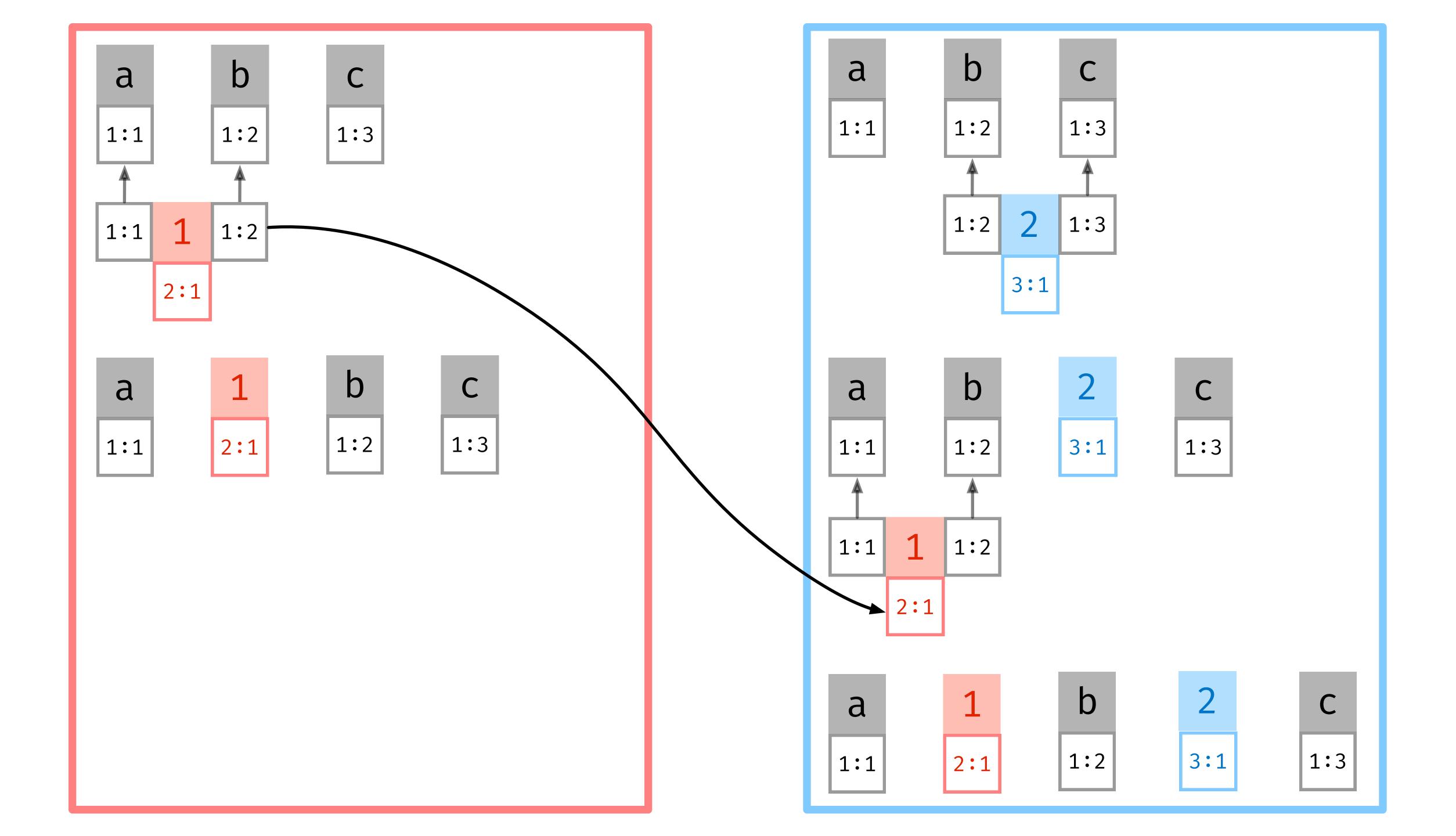


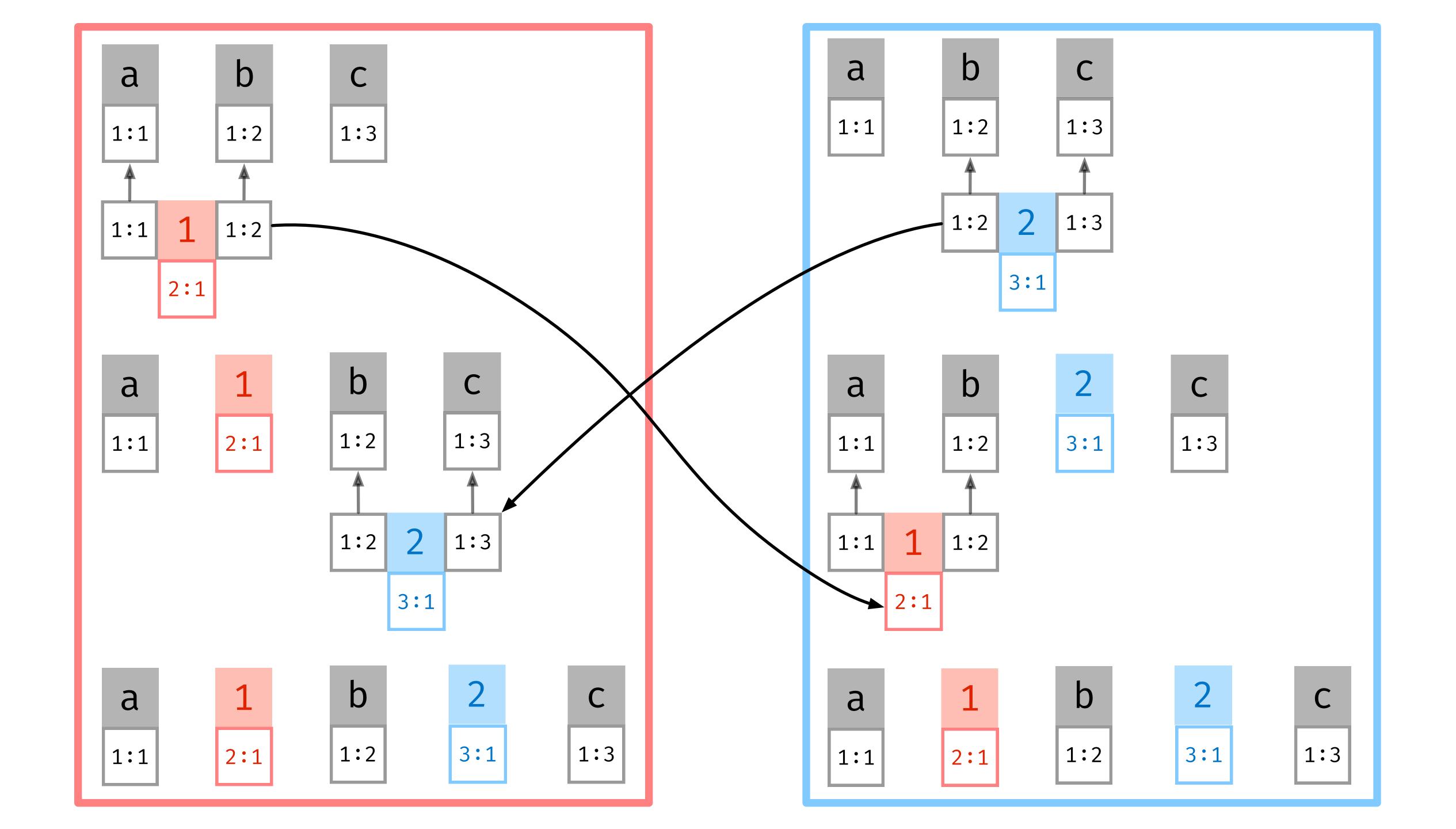
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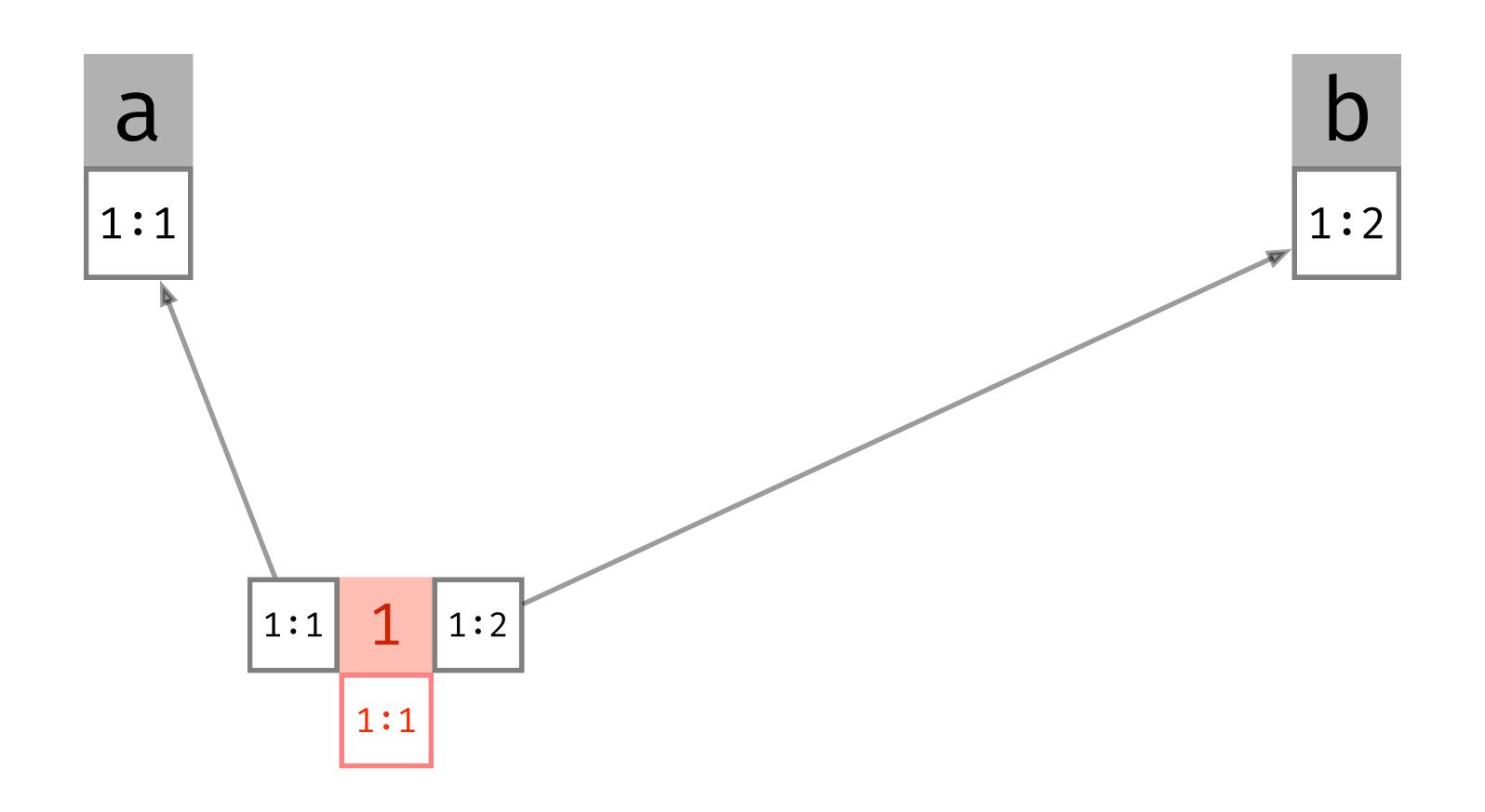


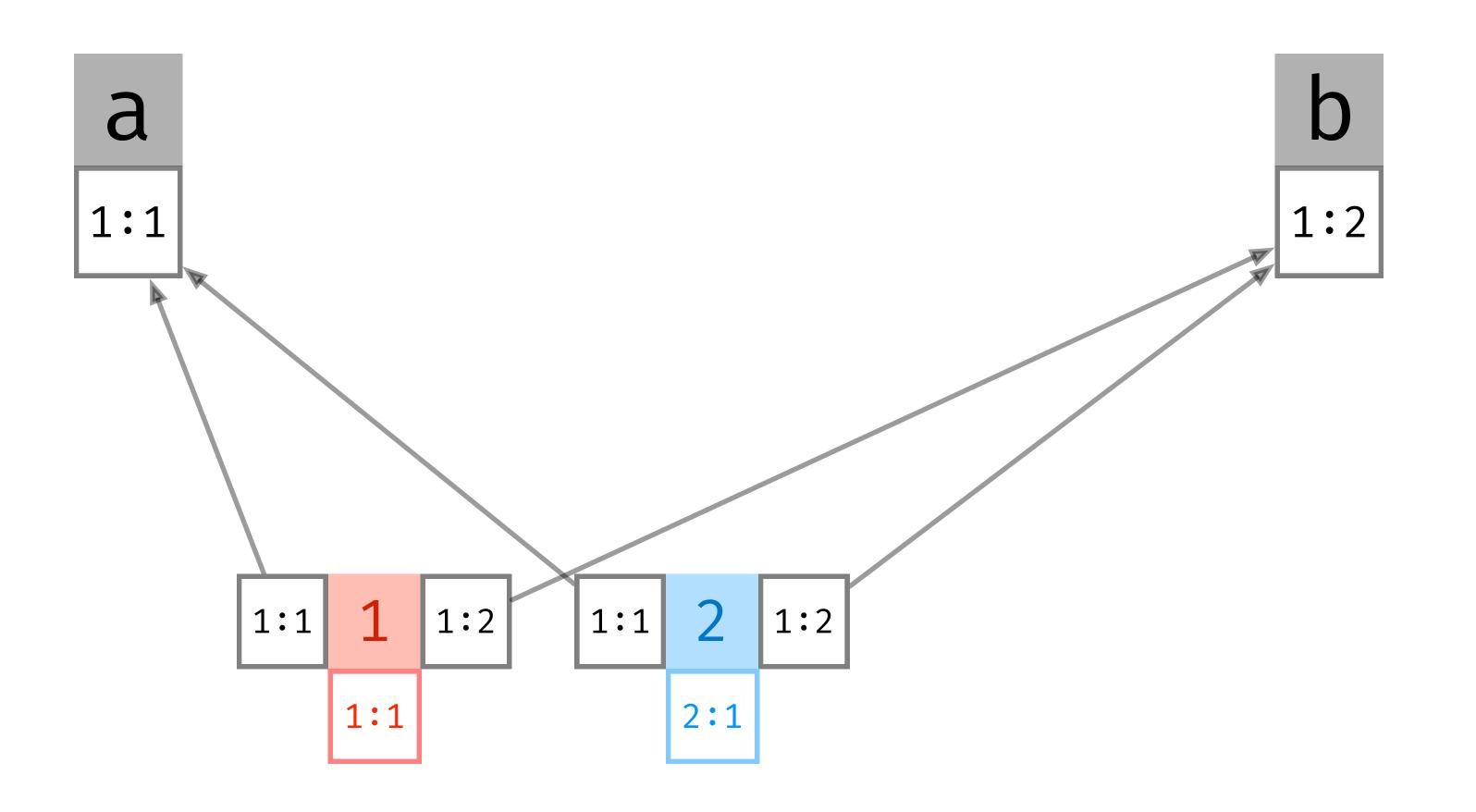
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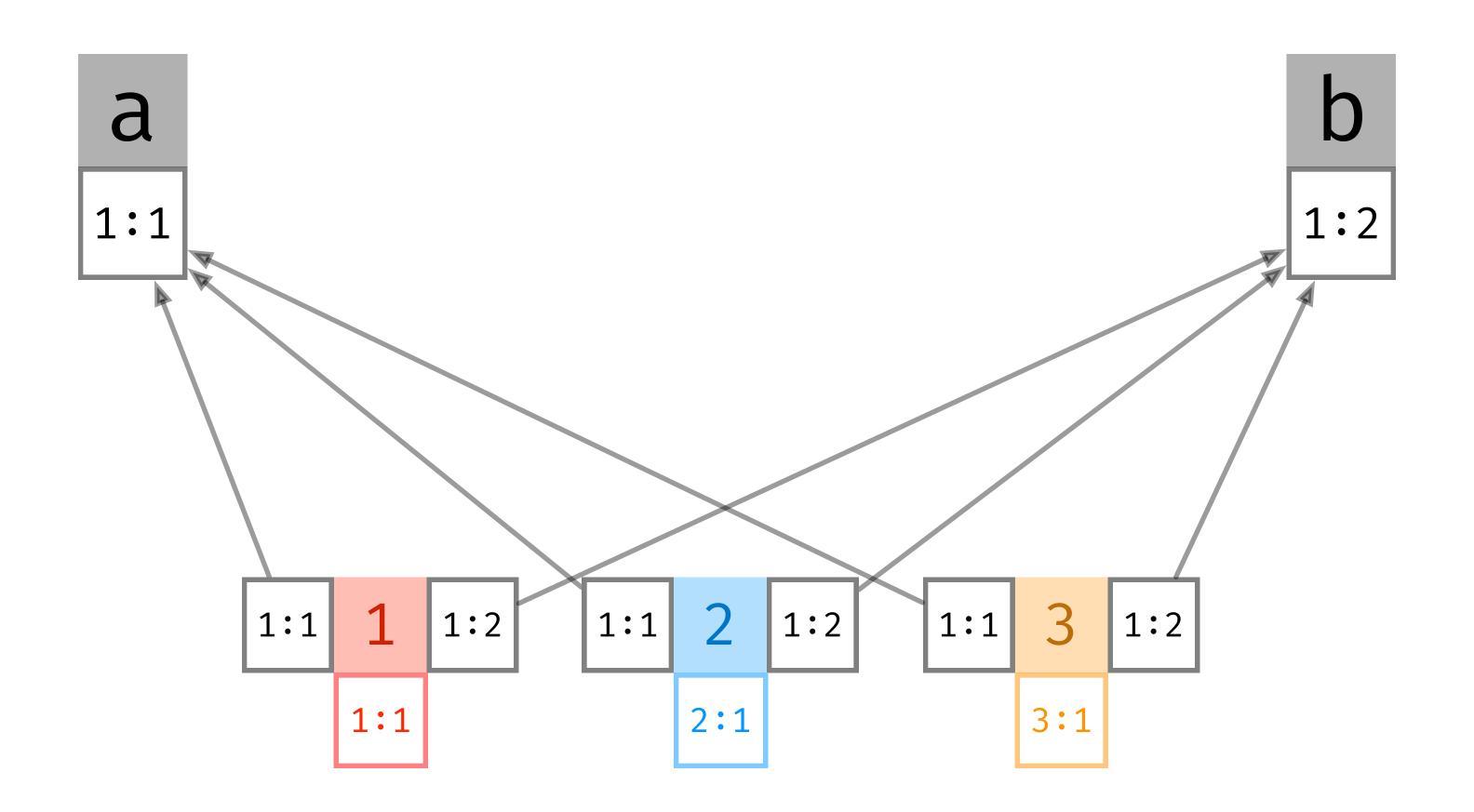
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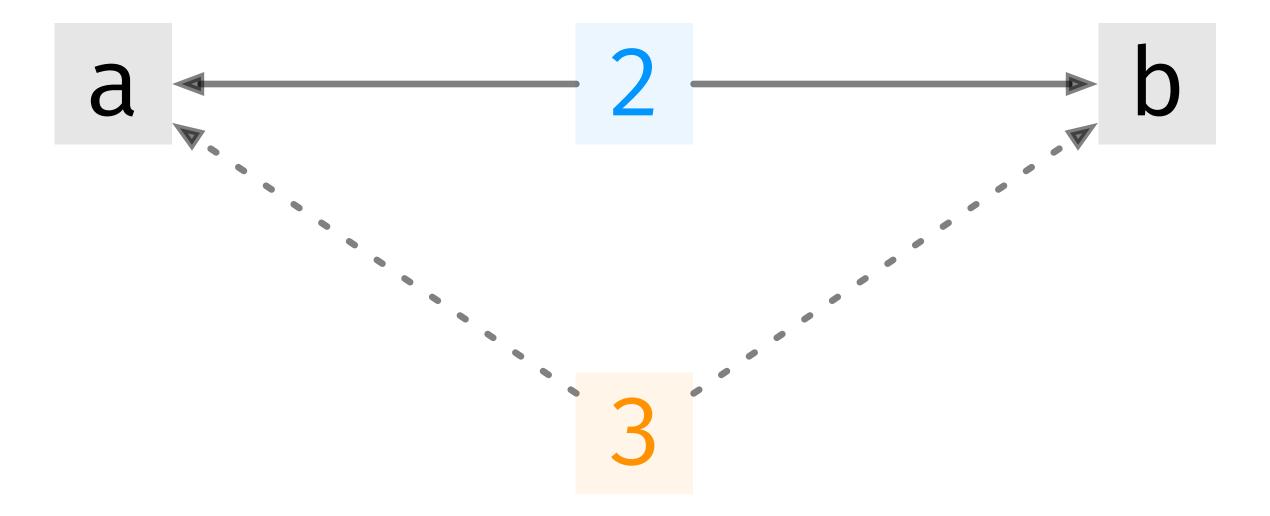
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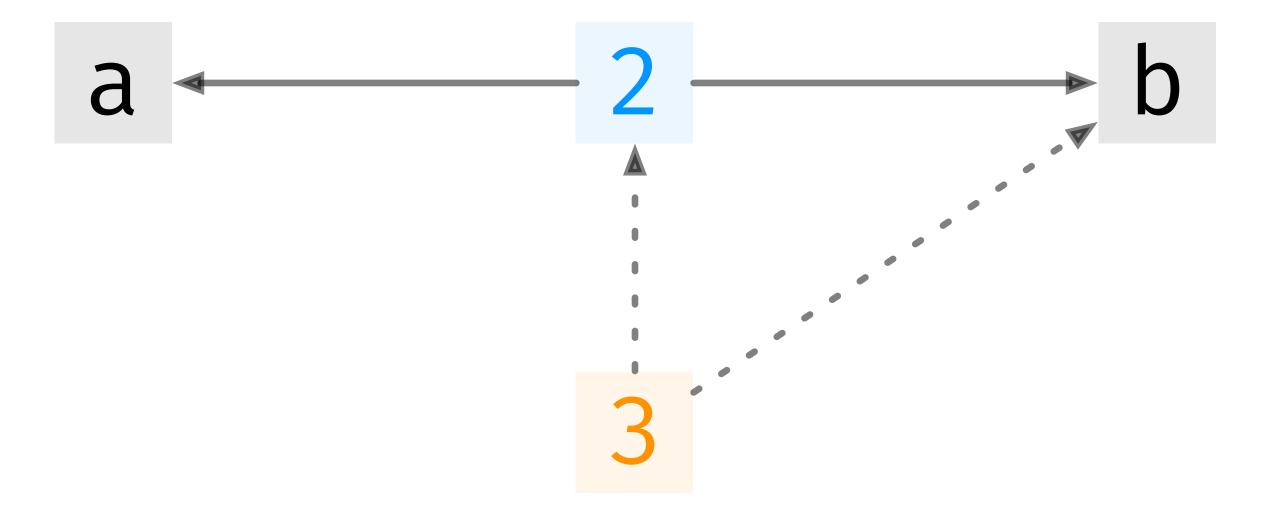


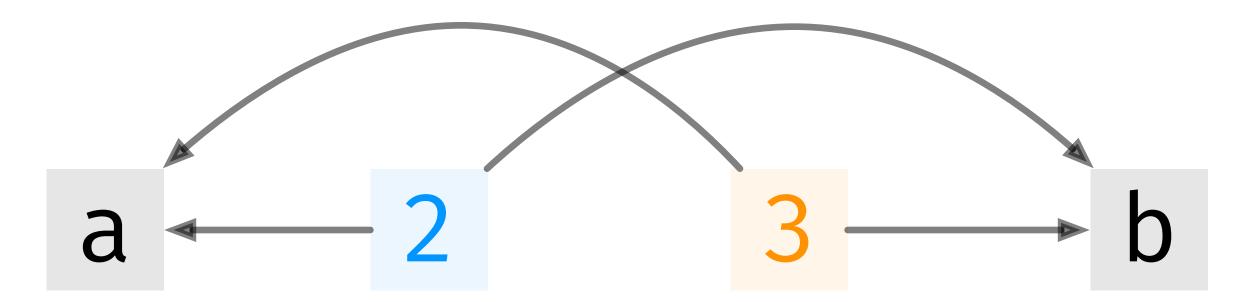


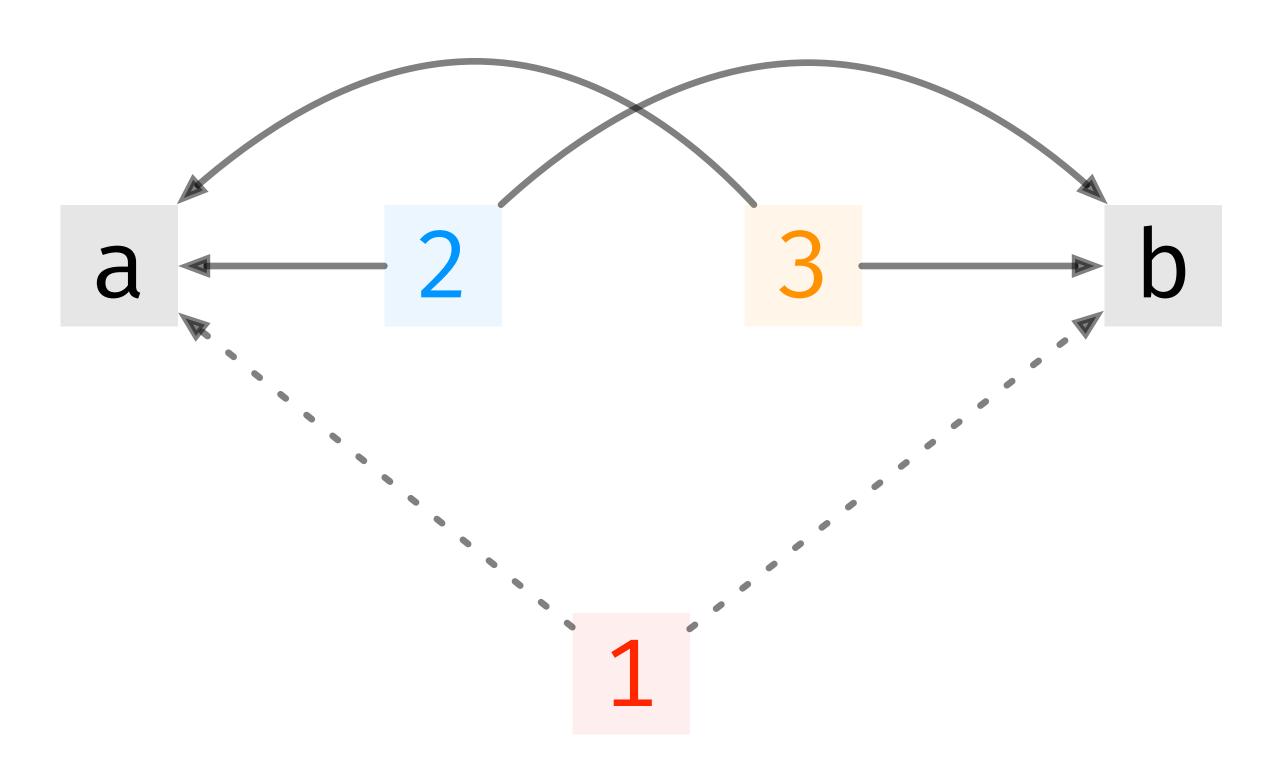


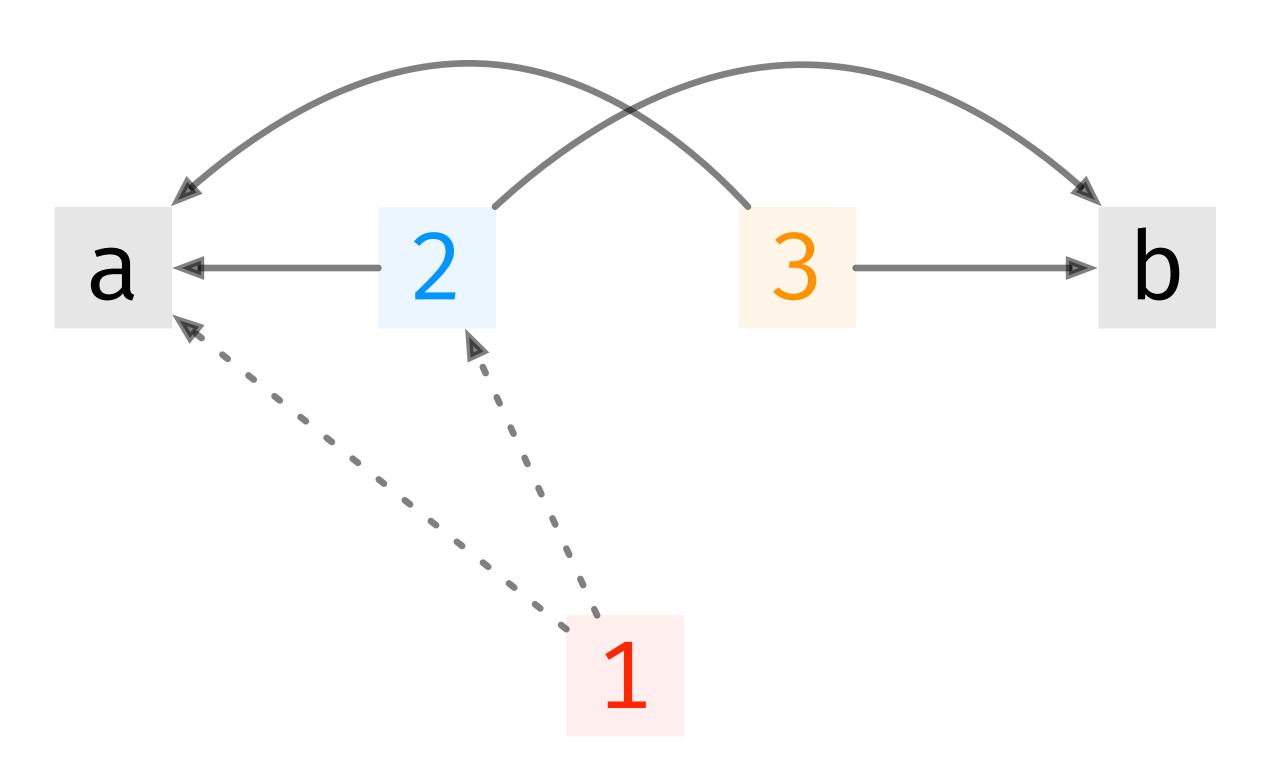


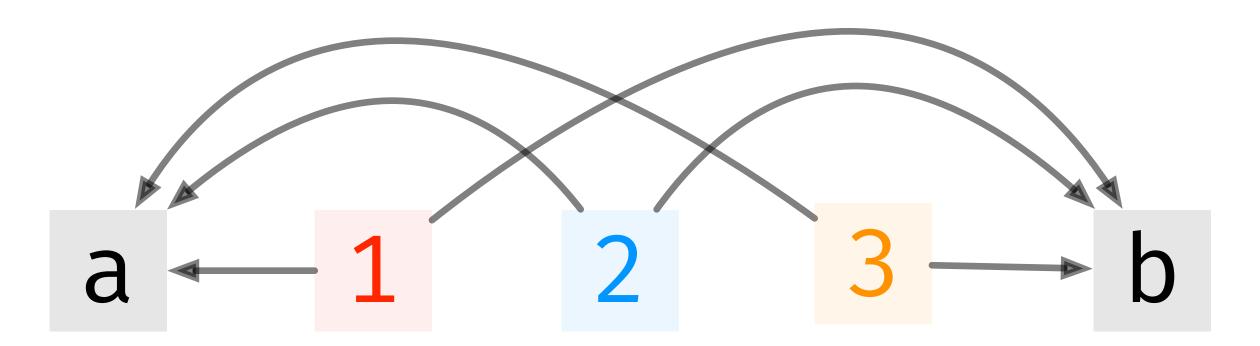












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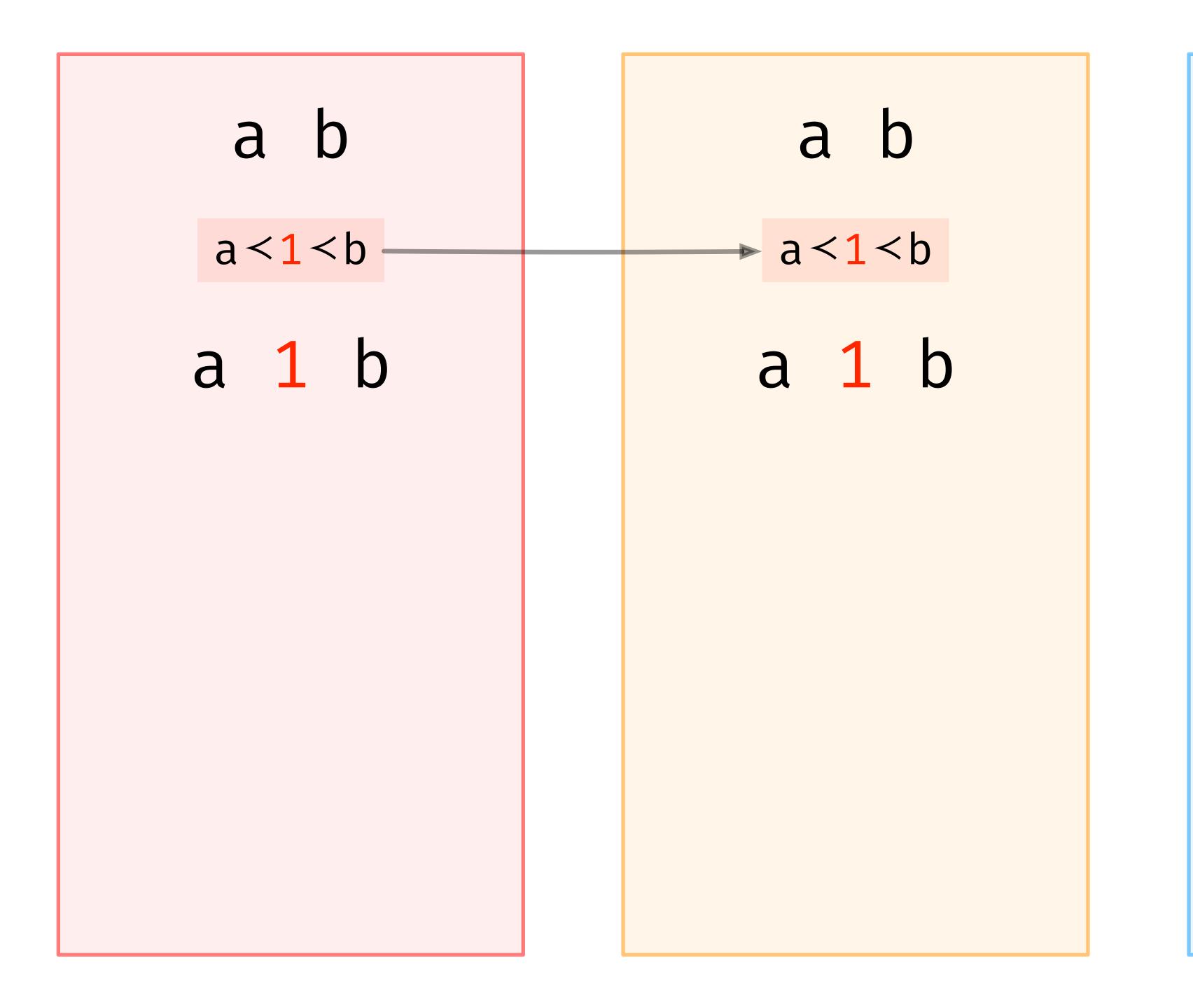
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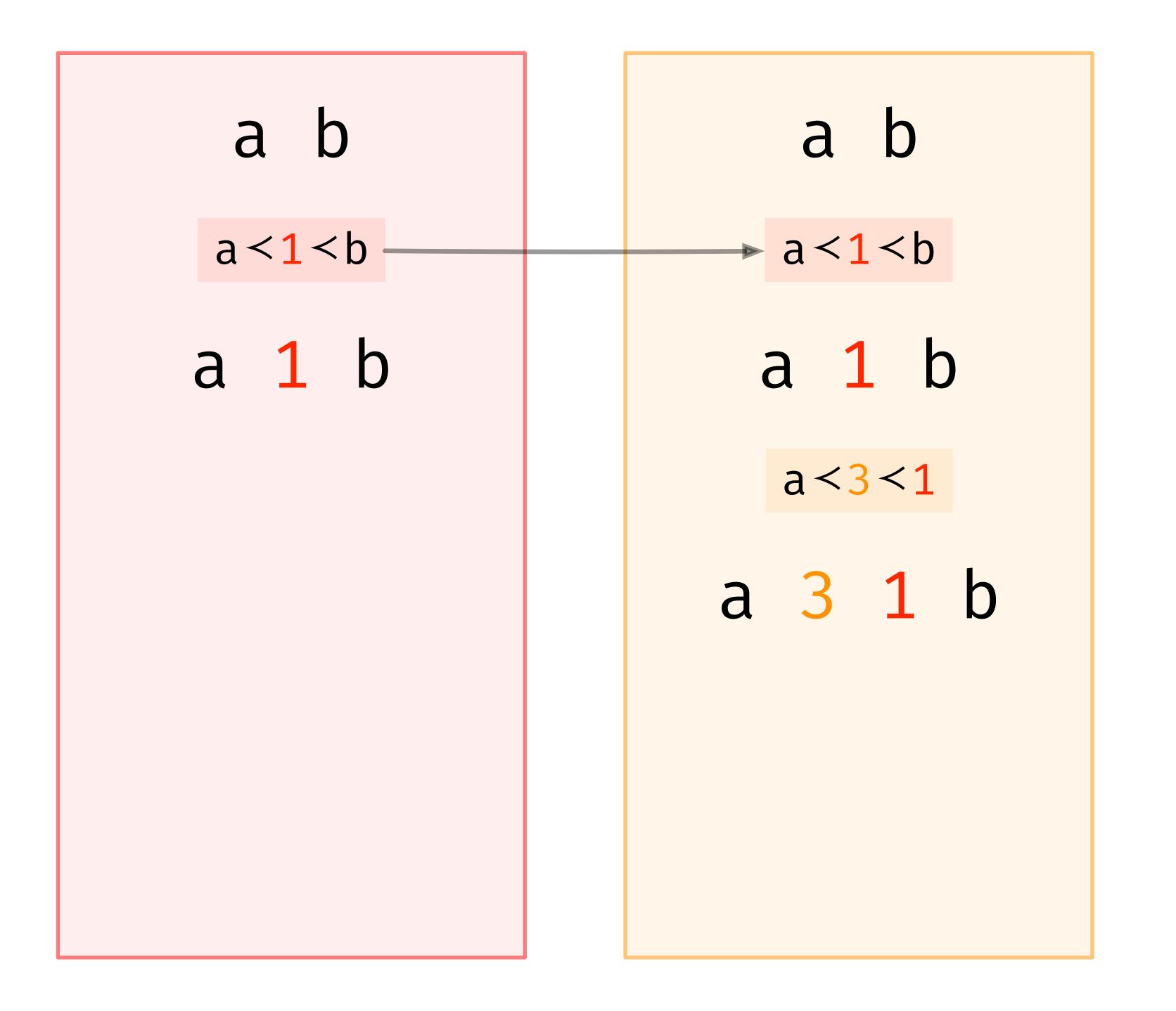
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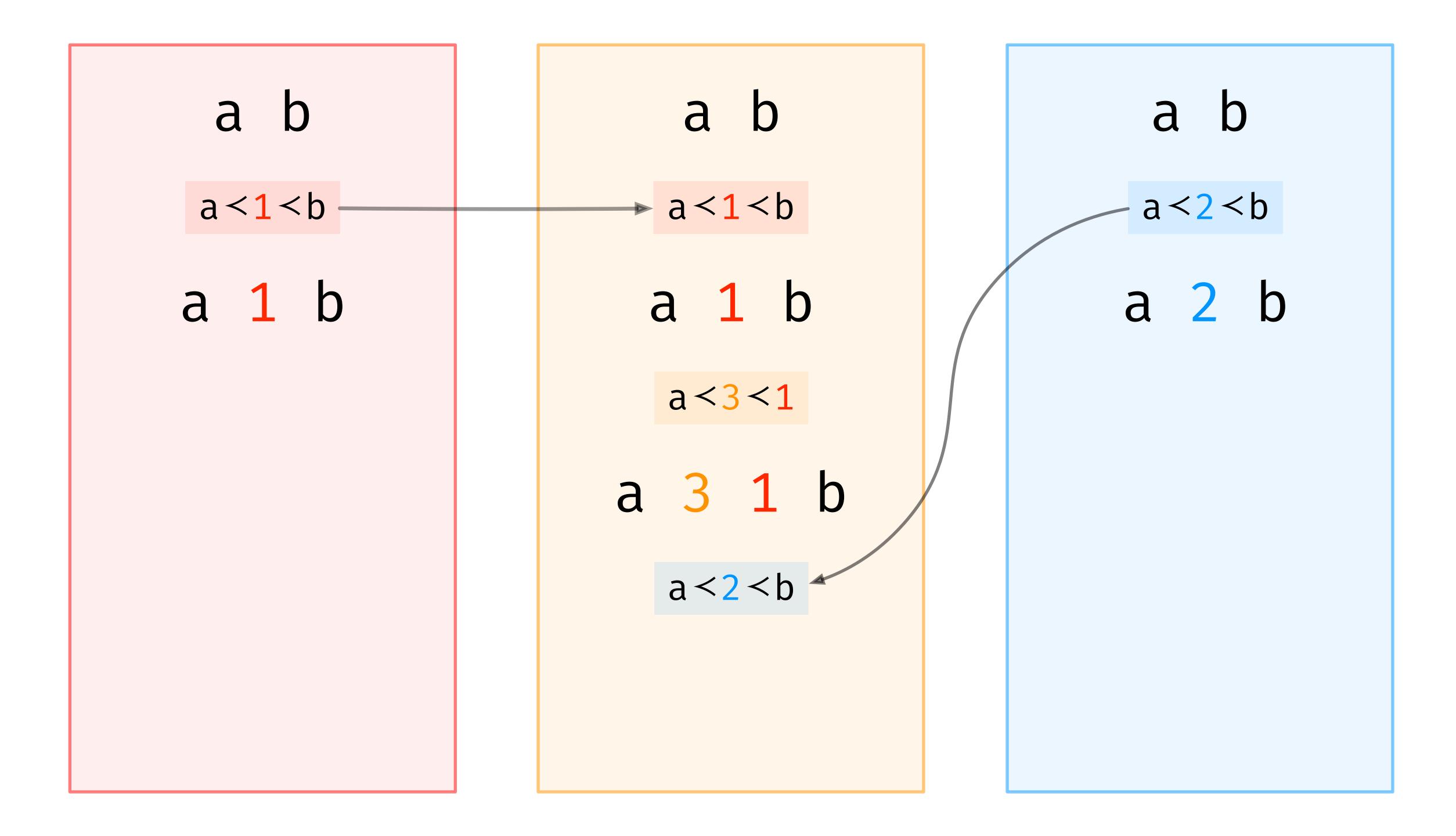
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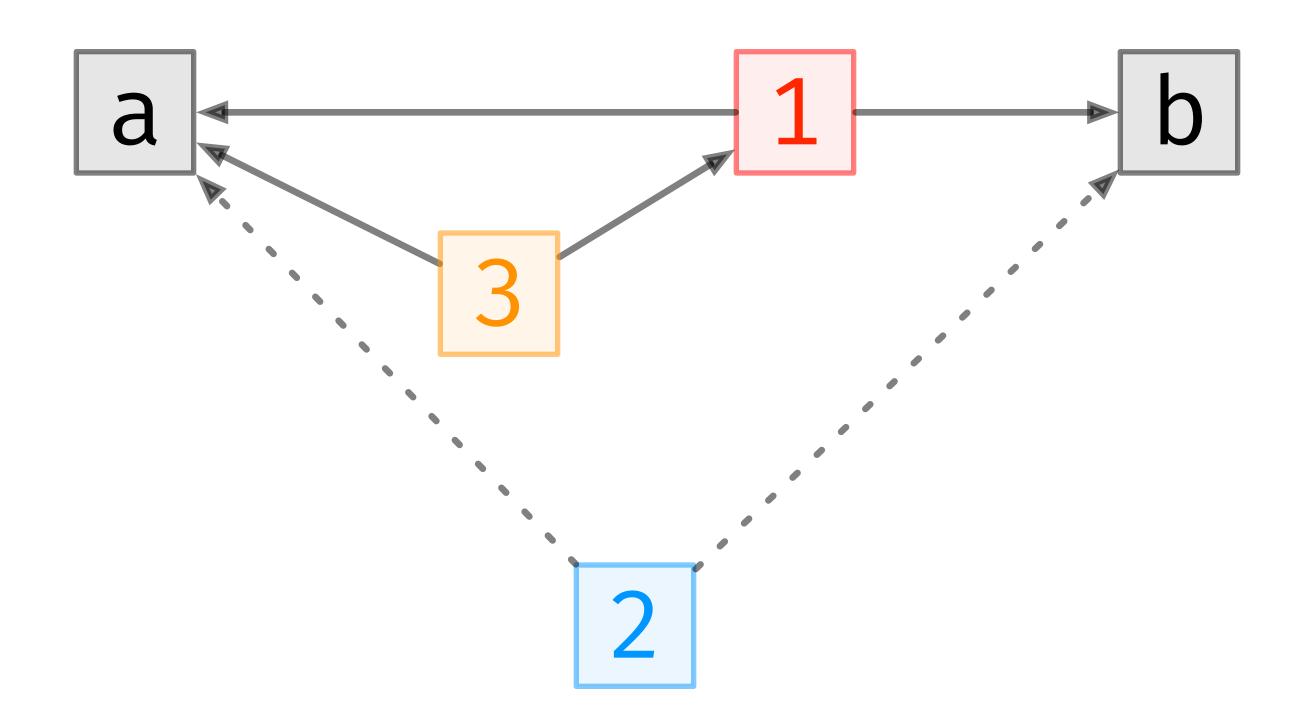


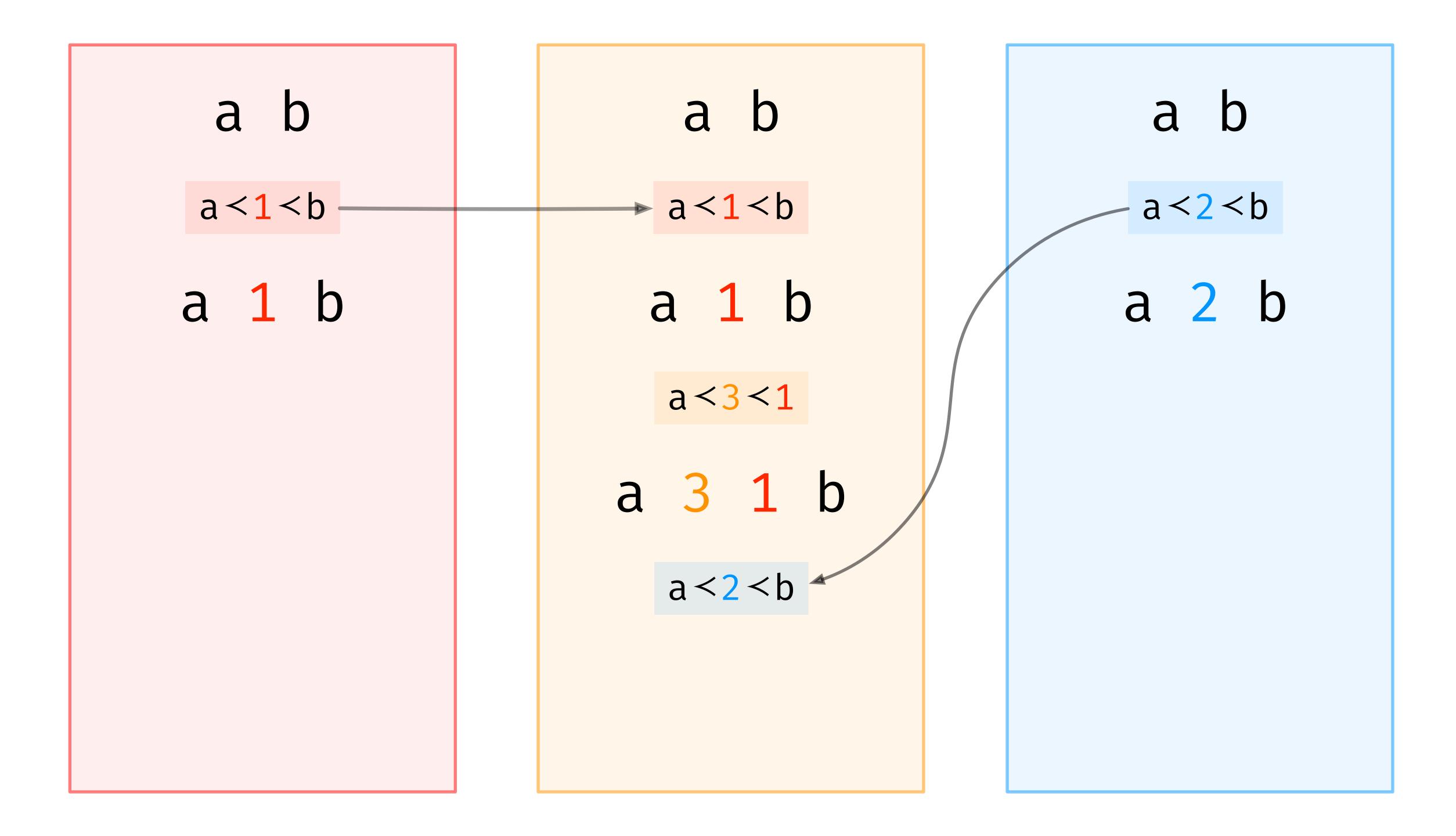
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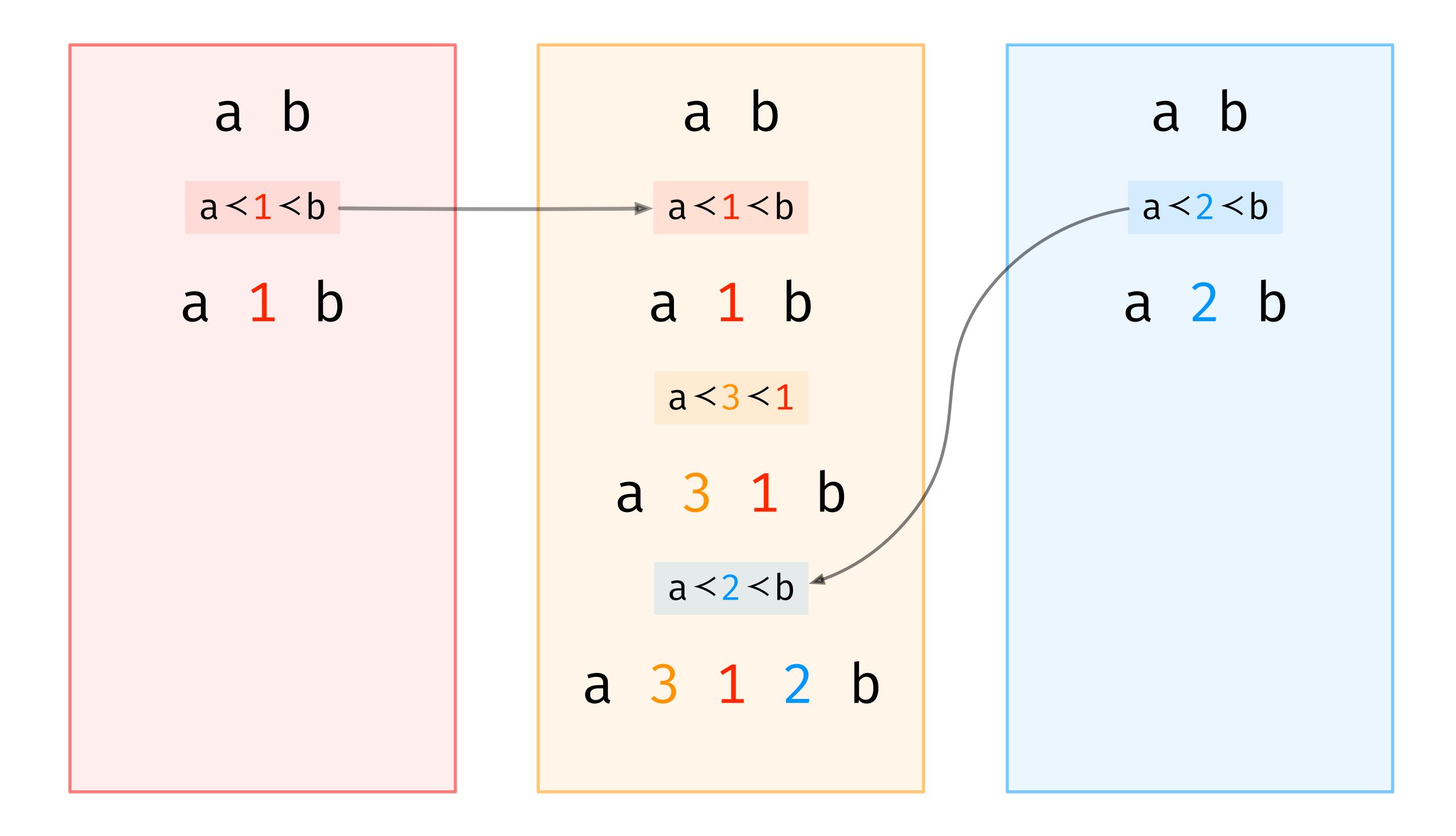


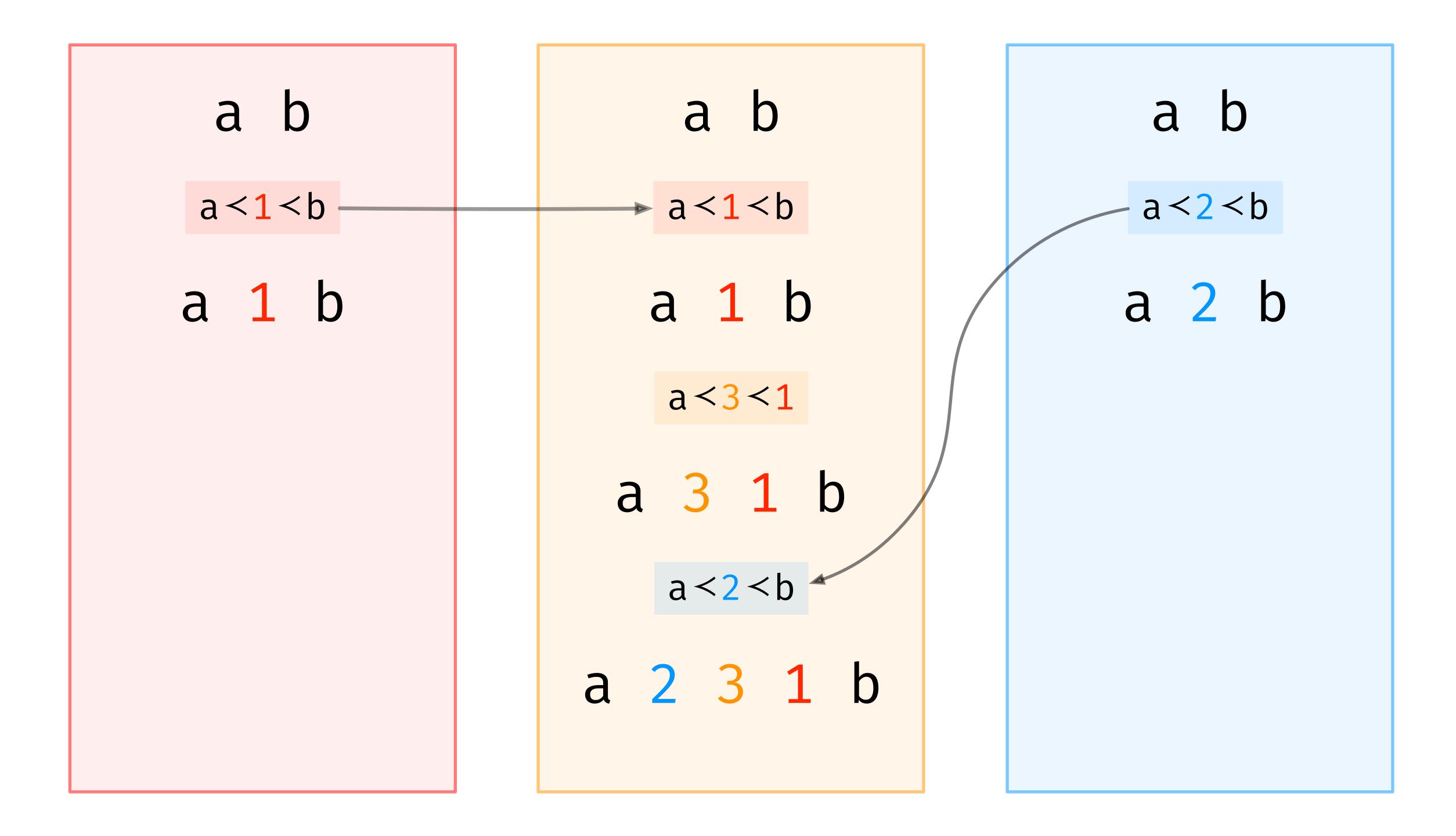
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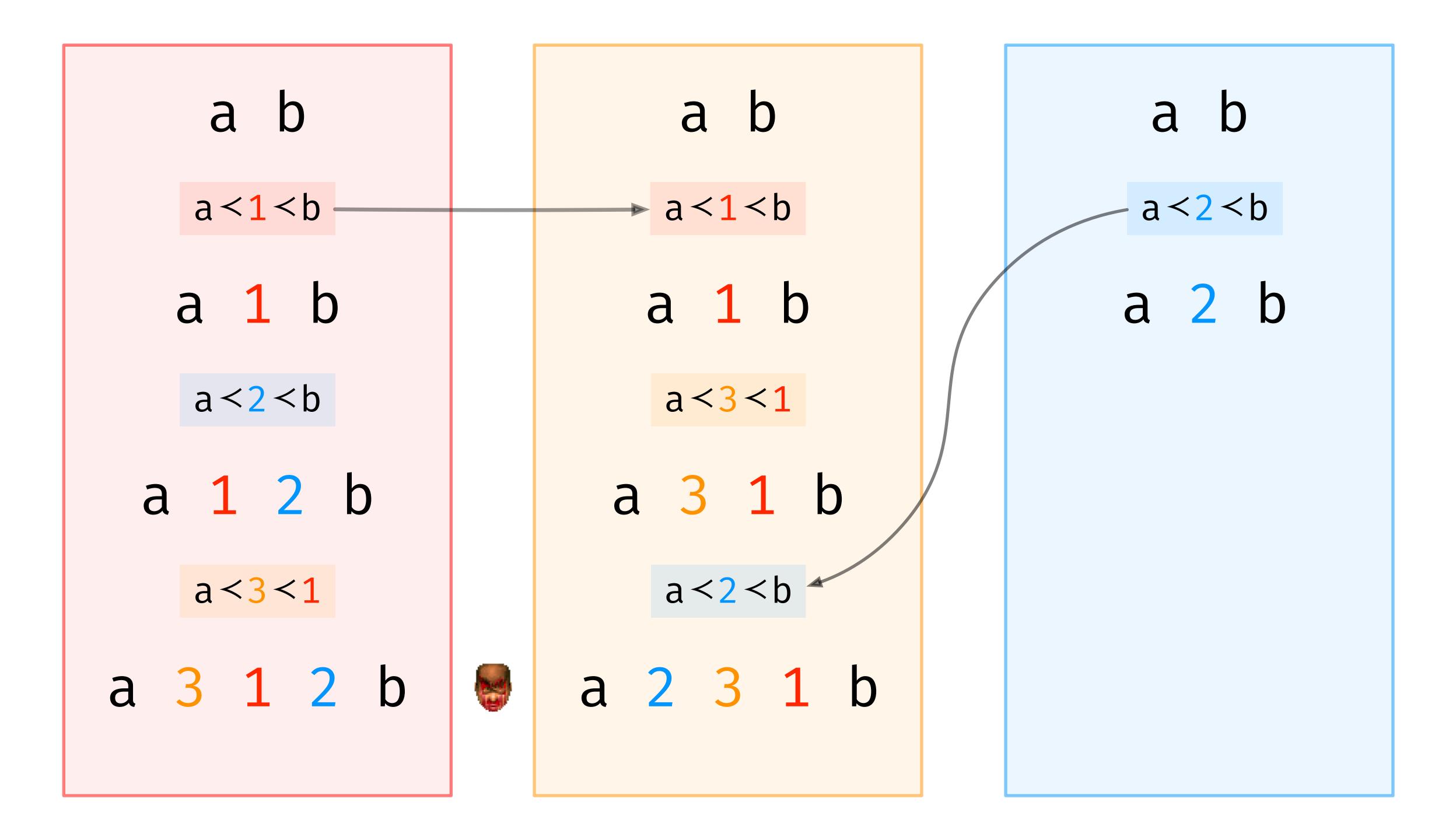


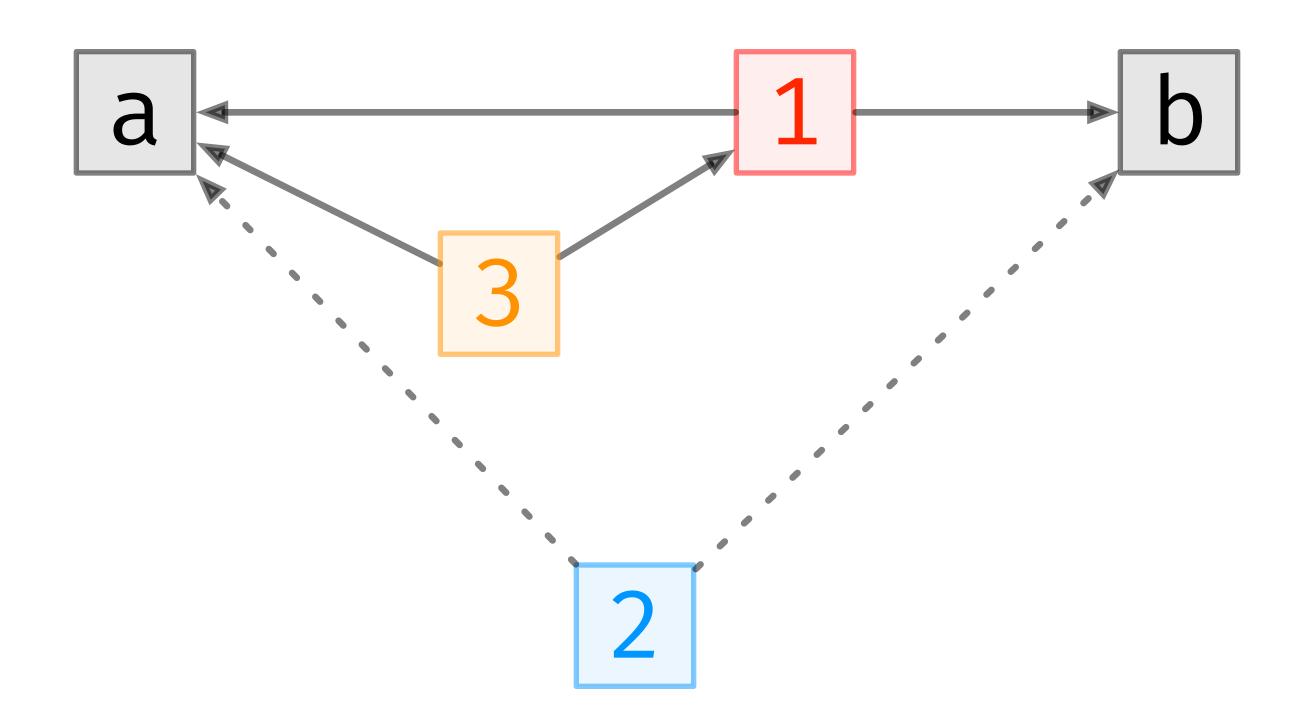


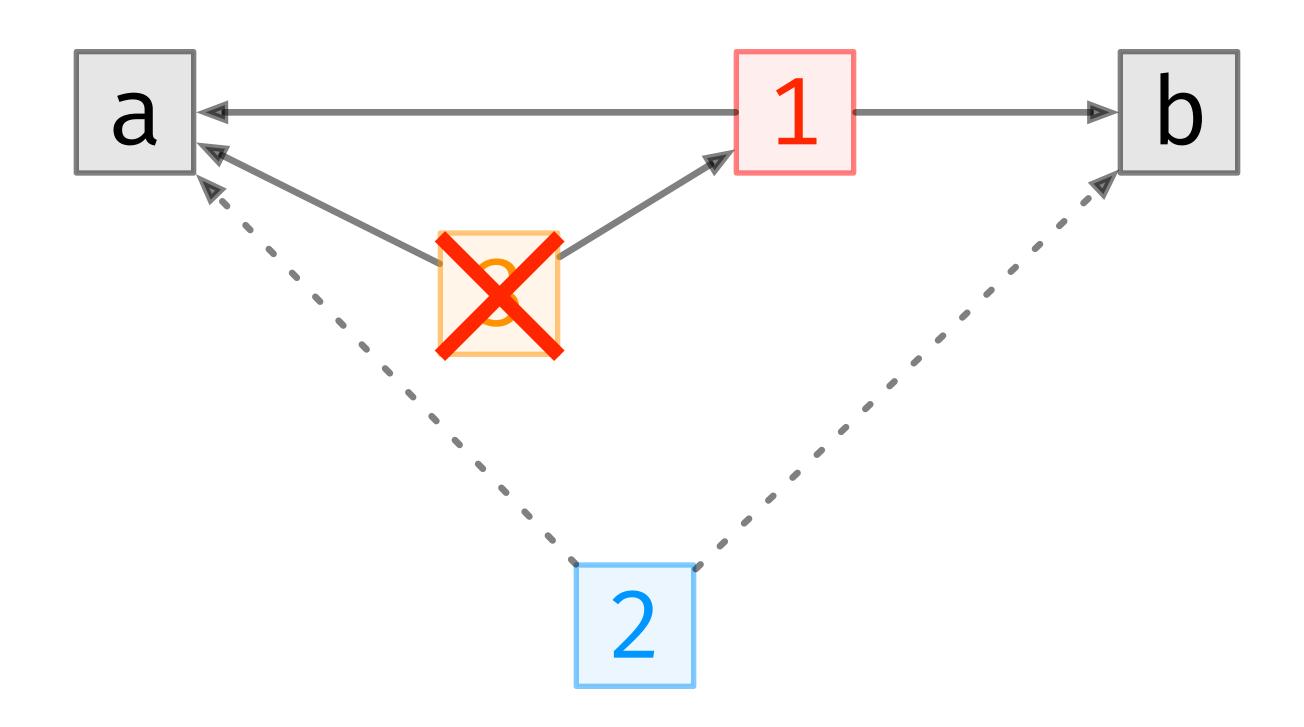


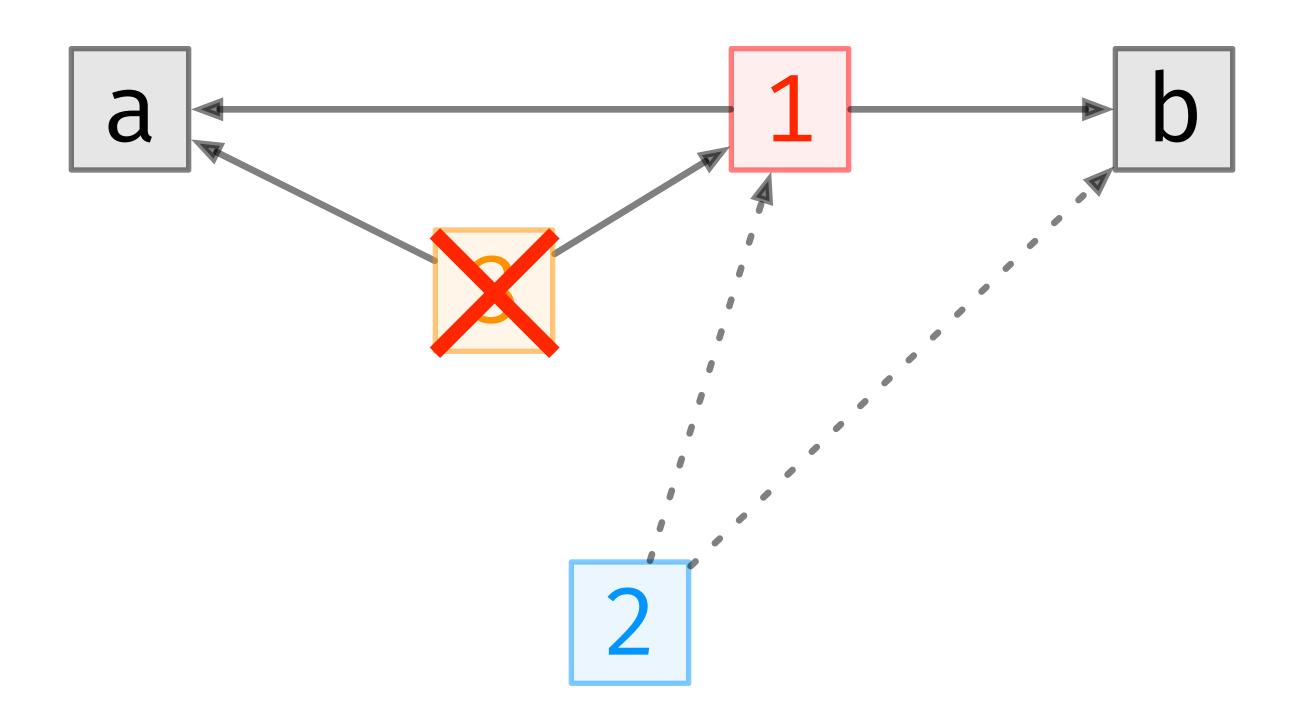
a b a b a b a < 2 < ba < 1 < b \rightarrow a < 1 < ba 2 b a 1 b a 1 b a < 2 < ba < 3 < 1a 1 2 b a 3 1 b a < 2 < b

a b a b a b a < 2 < ba < 1 < b \rightarrow a < 1 < ba 2 b a 1 b a 1 b a < 3 < 1a < 2 < ba 1 2 b a 3 1 b a < 2 < ba < 3 < 1









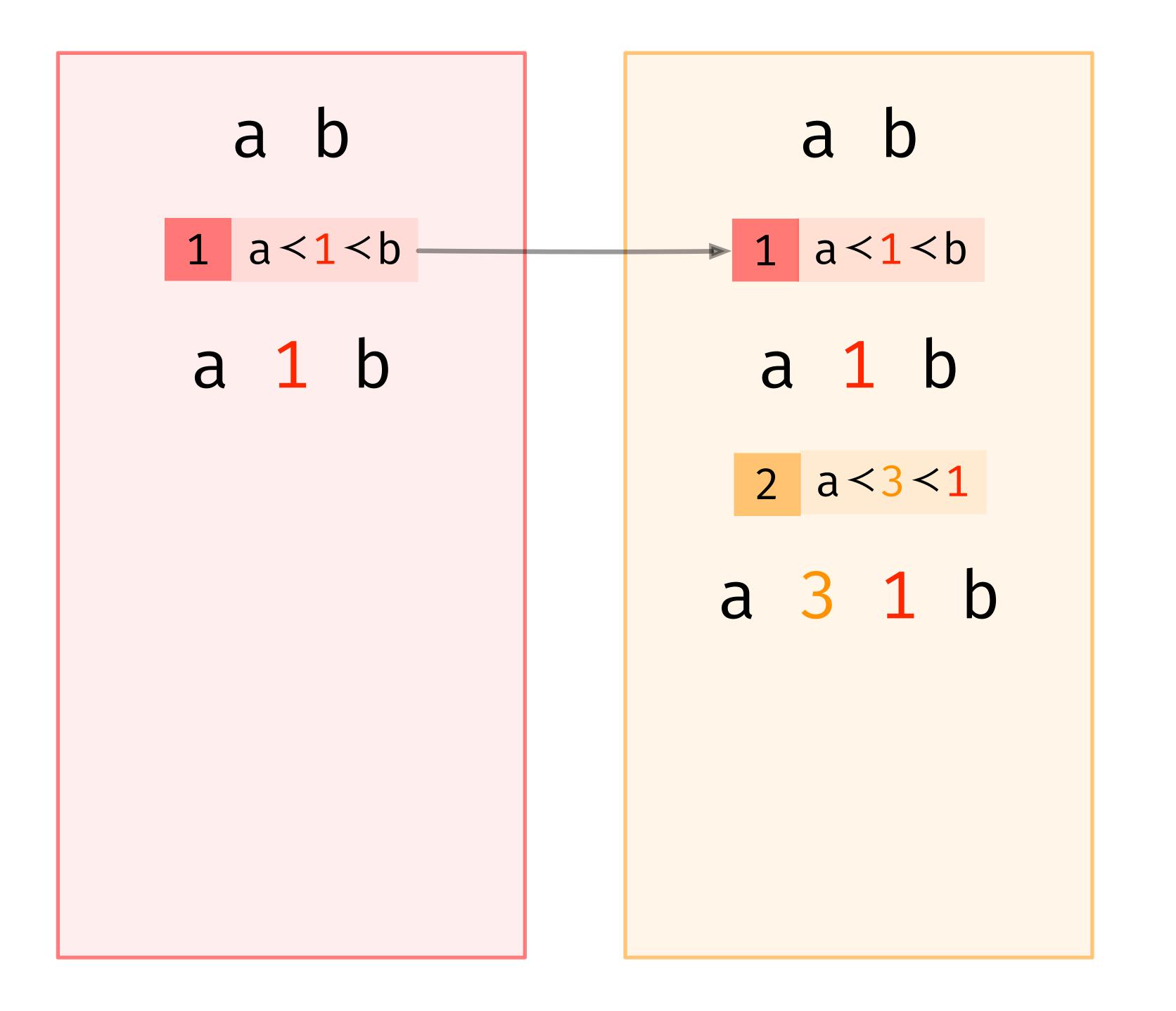
1 a<1<b

a 1 b

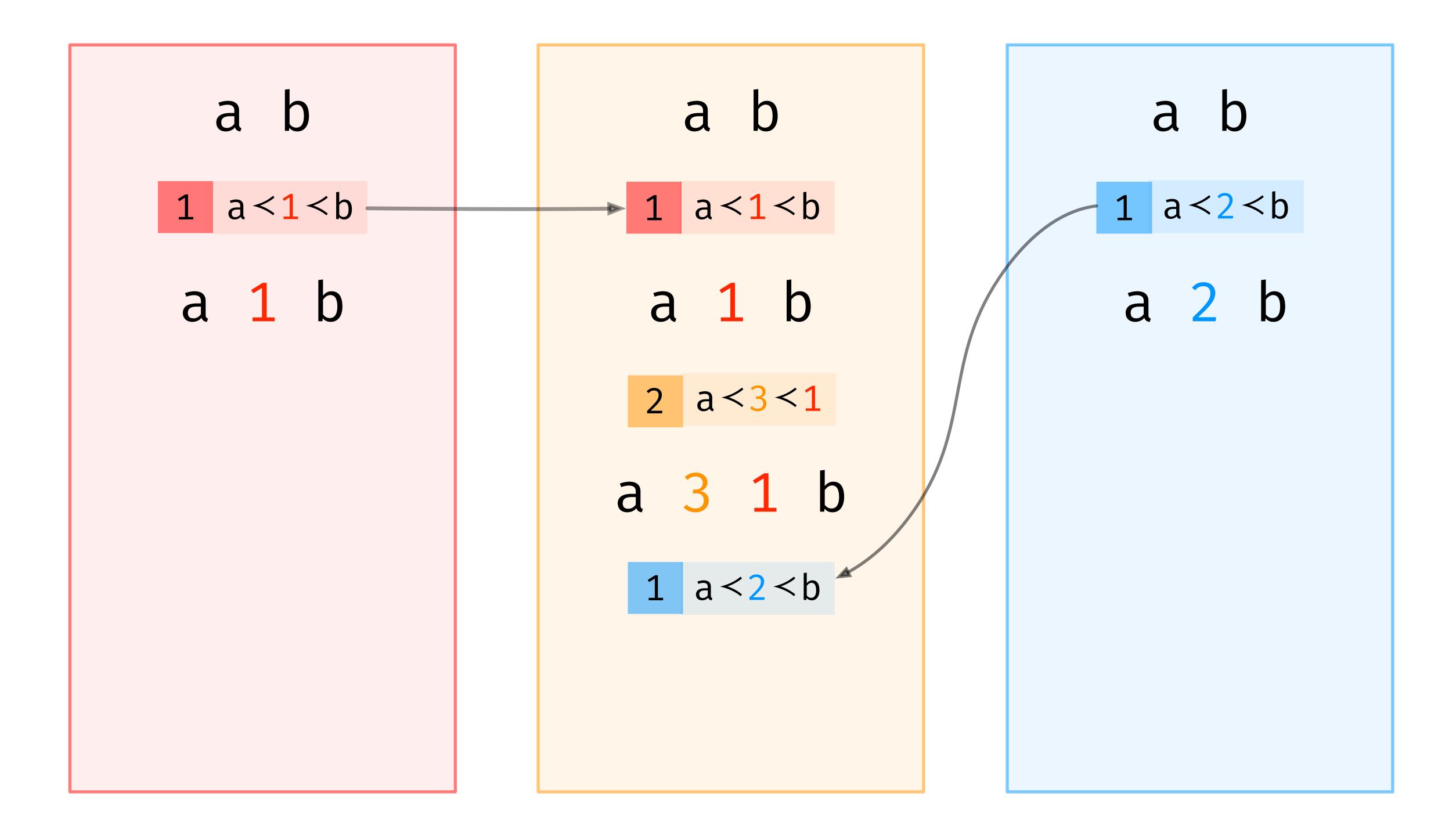
a b

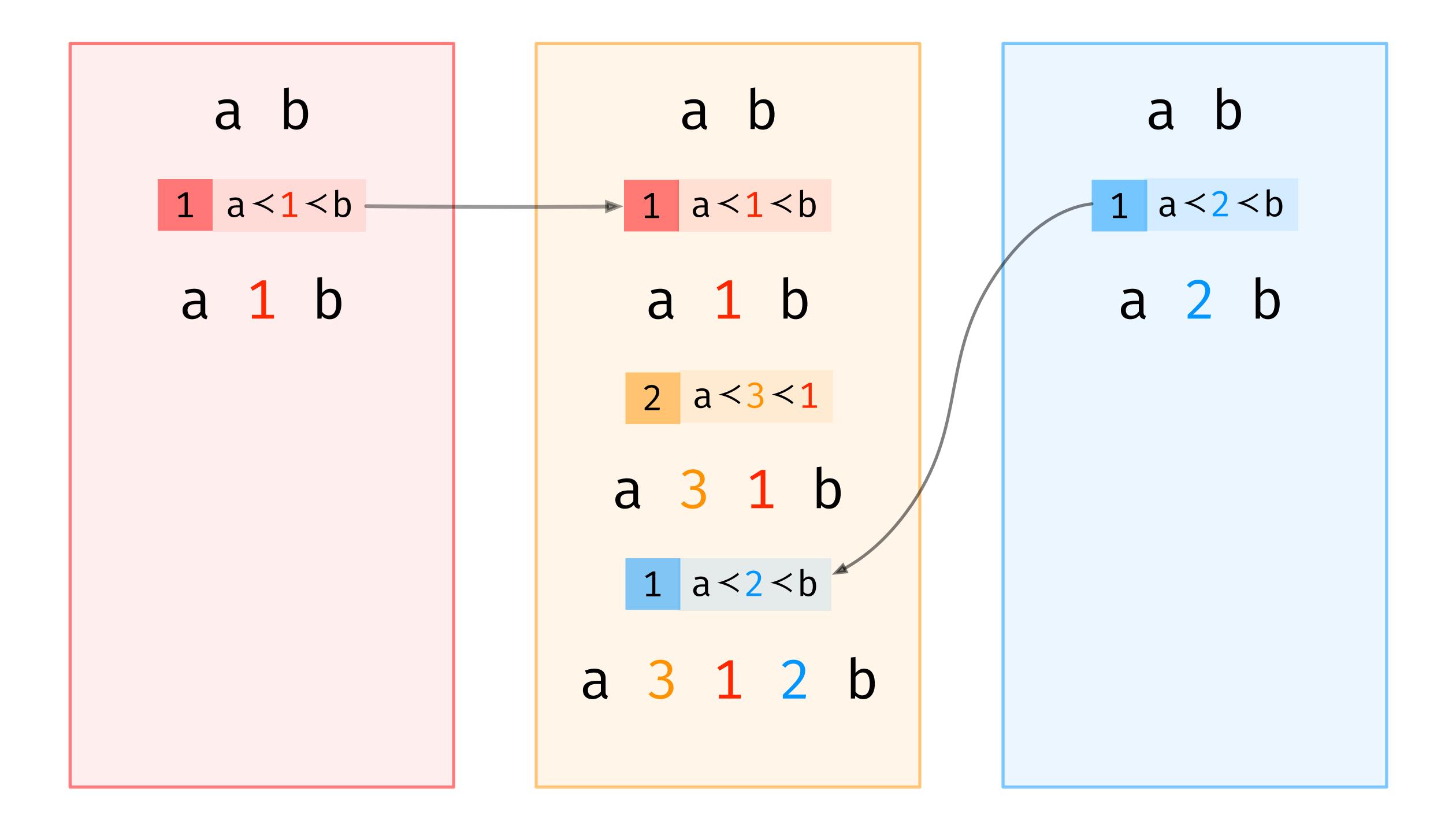
a b

1 a < 2 < b



a ba < 2 < ba 2 b





a b a b a b 1 a < 2 < b1 a<1<b-1 a<1<b a 2 b a 1 b a 1 b 1 a < 2 < b 2 a < 3 < 1 a 3 1 b a 1 2 b 1 a < 2 < b

a b a b 1 a<1<b-1 a < 1 < b a 1 b a 1 b 1 a < 2 < b 2 a < 3 < 1 a 1 2 b a 3 1 b 1 a<2<b 2 a < 3 < 1

a b1 a < 2 < ba 2 b

a b a b 1 a<1<b-1 a<1<b a 1 b a 1 b 1 a < 2 < b 2 a < 3 < 1 a 1 2 b a 3 1 b 1 a < 2 < b 2 a < 3 < 1

a ba < 2 < ba 2 b

Supporting String-Wise Operations and Selective Undo for Peer-to-Peer Group Editing

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ABSTRACT

Real-time group editing has been envisioned as an effective manner of collaboration. For years, operational transformation (OT) has been the standard concurrency control mechanism for real-time group editing, due to its potential for high responsiveness to local editing operations. OT algorithms are generally non-trivial to be error-free and are computation intensive. Recently, commutative replicated data types (CRDT) have appeared as an alternative to OT. The state-of-the-art OT and CRDT work still lacks the basic functionality found in single-user text editors. In particular, there is no published work that supports both string-wise operations and selective undo. This paper presents an approach that combines and extends OT and CRDT strengths. It is fully decentralized and supports string-wise editing operations and selective undo. Our performance study shows that it provides sufficient responsiveness to the end-users.

Categories and Subject Descriptors

C.2.4 [Computer-Communication Networks]: Distributed Systems—Distributed applications; H.5.3 [Information Interfaces and Presentation]: Group and Organization Interfaces—Collaborative computing

General Terms

Algorithms, Performance

Keywords

Real-time collaborative editor, commutative replicated data type, operation transformation.

1. INTRODUCTION

A real-time group editor allows multiple users to simultaneously edit the same document from different places. Fully decentralized, or peer-to-peer, collaboration has generally the advantage of availability, scalability and resistance to censorship and surveillance, over collaboration via a central server.

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Copyright 2014 ACM 978-1-4503-3043-5/14/11 ...\$15.00. http://dx.doi.org/10.1145/2660398.2660401. Operational transformation (OT) has been established as a concurrency control mechanism for real-time group editing due to its potential for high responsiveness to local operations [4, 5, 6, 7, 8, 10, 14, 15, 16, 17, 18, 19, 21]. Local operations are executed immediately at local peers and later transformed and integrated at remote peers. OT algorithms are sophisticated. Counterexamples of several published OT algorithms have been reported. Moreover, they have time complexity in the length of operation history of the document been edited, which potentially grows endlessly.

Recently, a new class of mechanisms called commutative replicated data types (CRDT) have been proposed [2, 11, 12, 13, 22, 23]. Concurrent operations of a CRDT are mutually commutative, so that a document is eventually kept consistent at all peers.

A real-time group editor should support at least the most basic functionality found in a single-user text editor. At its minimum, it should support insertion and deletion of single characters and strings of characters, as well as the undo and redo of the insertions and deletions. String-wise operations are important as they are the basis for other useful operations like copy-paste, select-delete and find-replace. Surprisingly, there is currently no published work that supports both string-wise operations and their undo.

Our work supports both string-wise operations and their undo by combining and extending existing OT and CRDT approaches.

The rest of this paper is organized as follows. Section 2 presents background and related work. Section 3 gives an overview of the approach. Section 4 presents the view-model architecture and the data structure of the model. Section 5 describes operations and updates in view and model. Section 6 describes how model and view are synchronized. Sections 7 and 8 describe how local and remote operations are integrated into the model. Section 9 shows the correctness of the approach. Section 10 presents performance results. Section 11 discusses some open issues. Section 12 concludes.

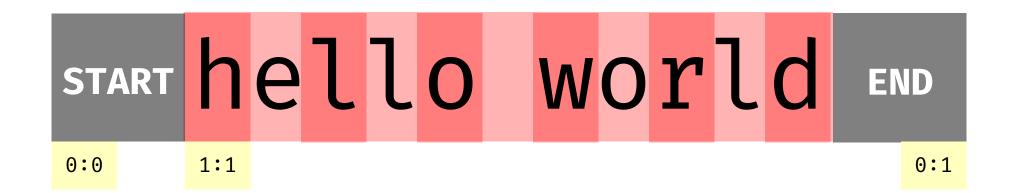
2. BACKGROUND

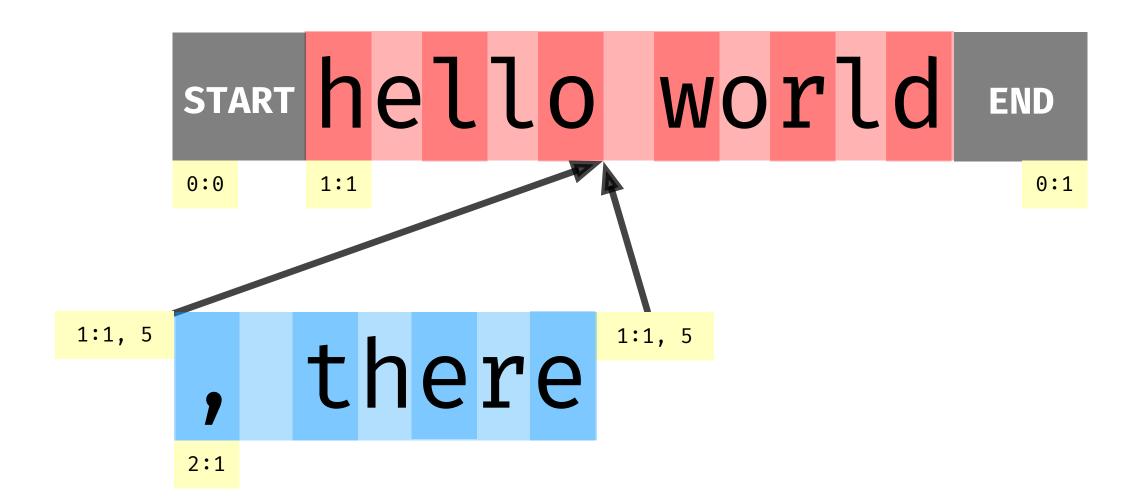
OT was first introduced in [4]. The basic idea is as the following. A shared document is replicated at different peers. An editing operation is first executed at a local peer and then propagated to remote peers. Suppose two peers start with "012". Peer 1 inserts "a" between "0" and "1" with ins(1, "a") and Peer 2 deletes "2" with del(2). The states after local executions at the two peers are "0a12" and "01". Now if the two peers execute the remote operations as is, the states at these peers become "0a2" and "0a1", which are inconsistent. With OT, the remote operations are transformed to include the executed concurrent operations, into ins(1, "a") and del(3) respectively. The two peers are in consistent state "0a1" after executing the transformed operations.

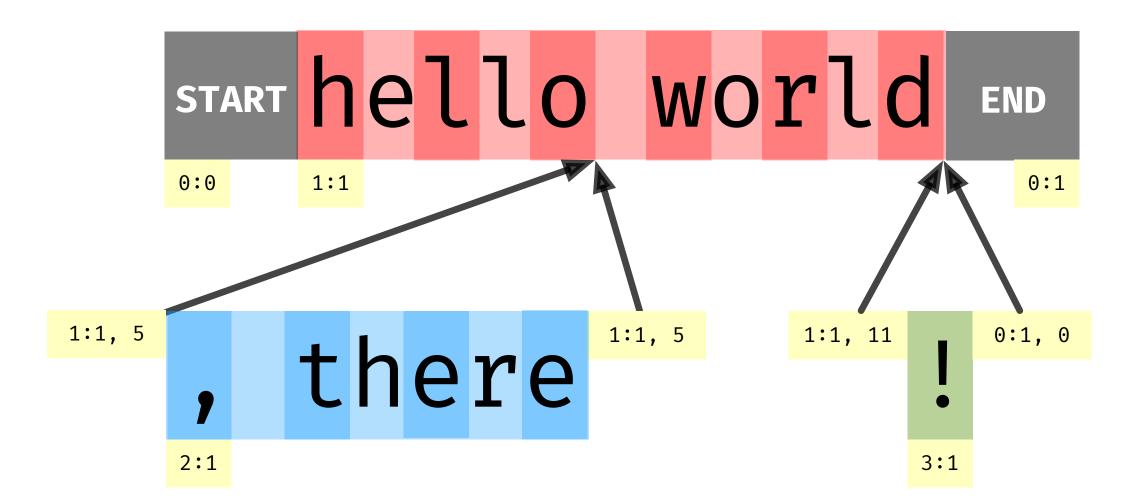
There are some challenges with this basic approach. First, a remote operation can only be transformed to include a concurrent

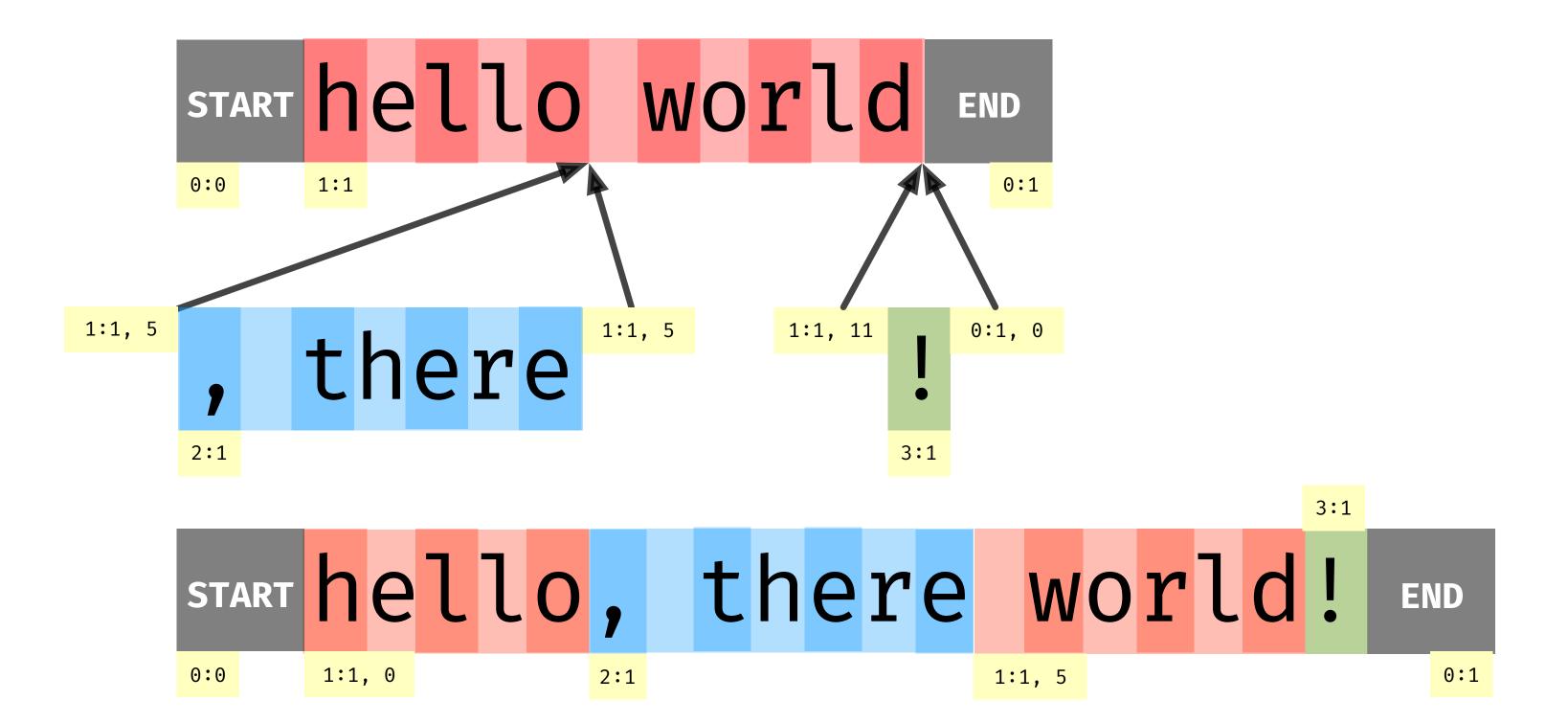
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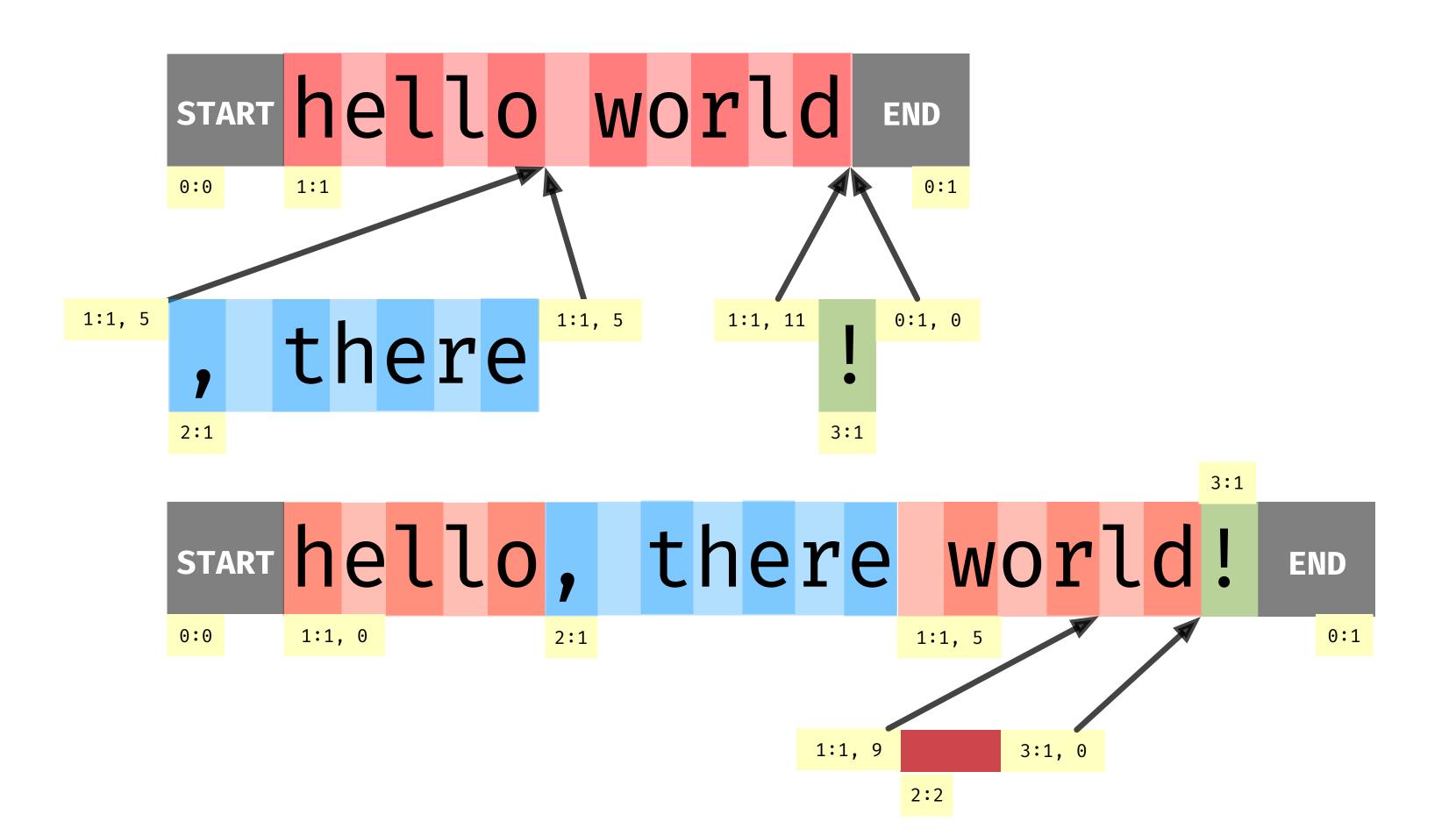
String-Wise WOOT 2014

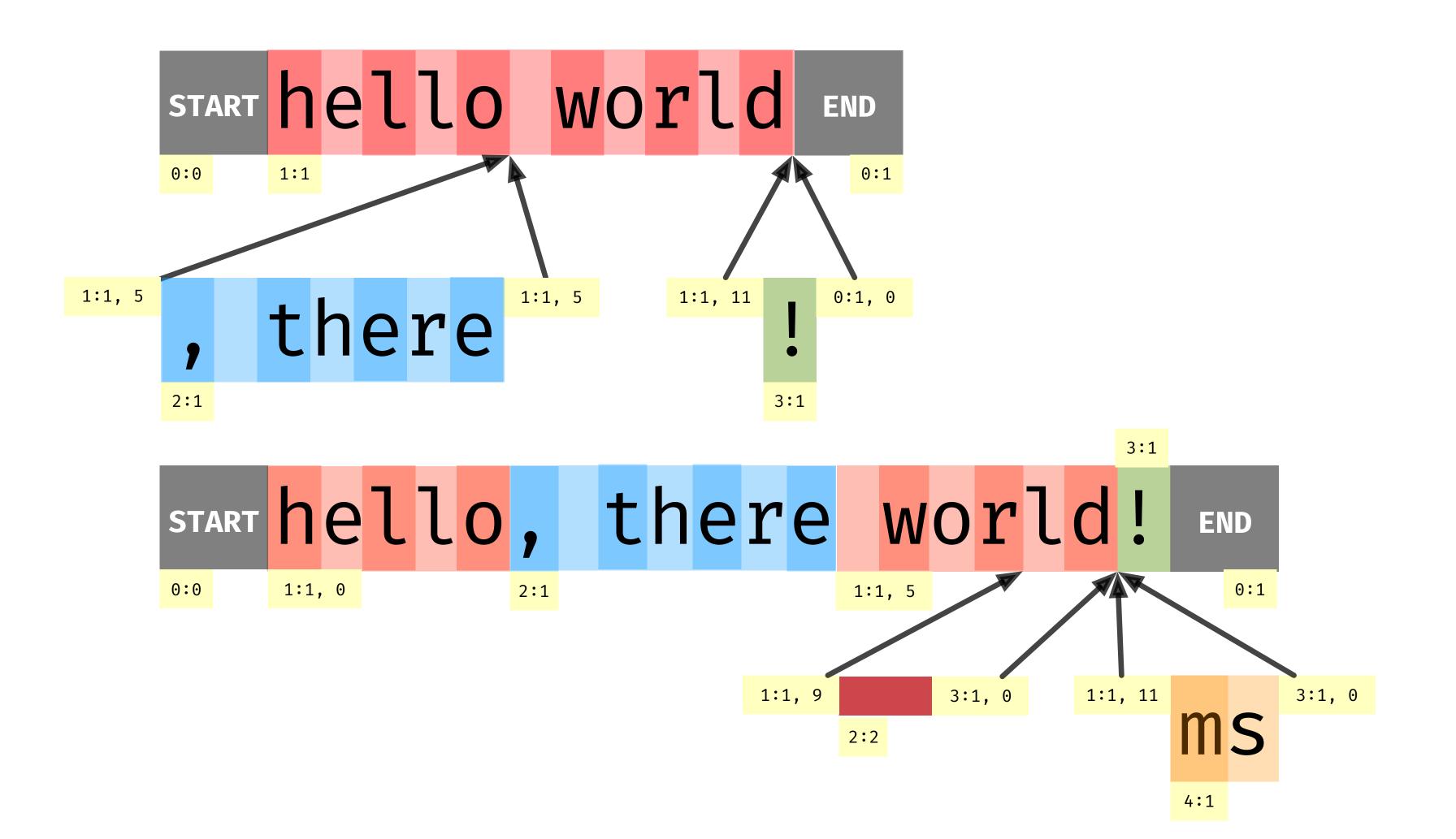


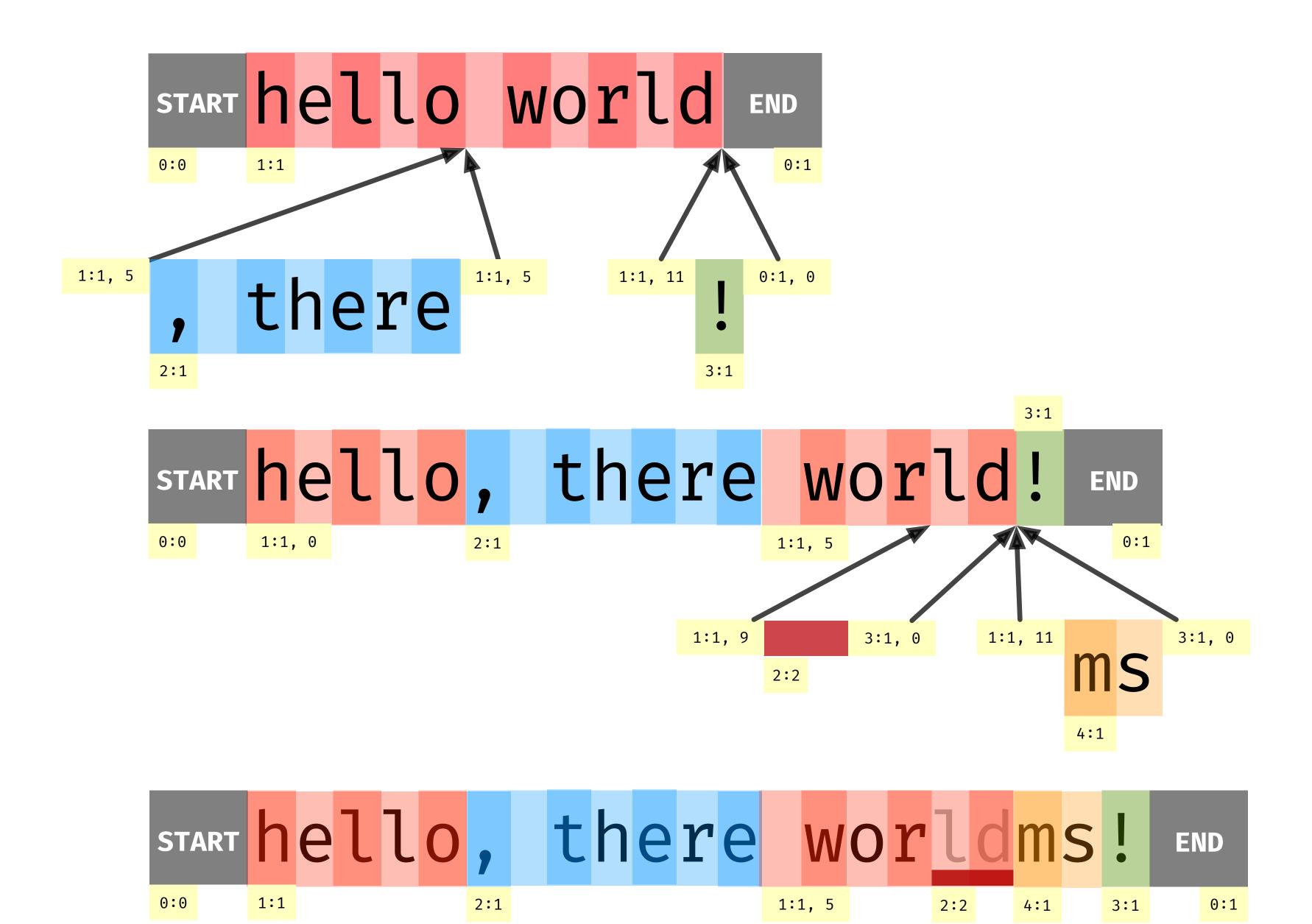






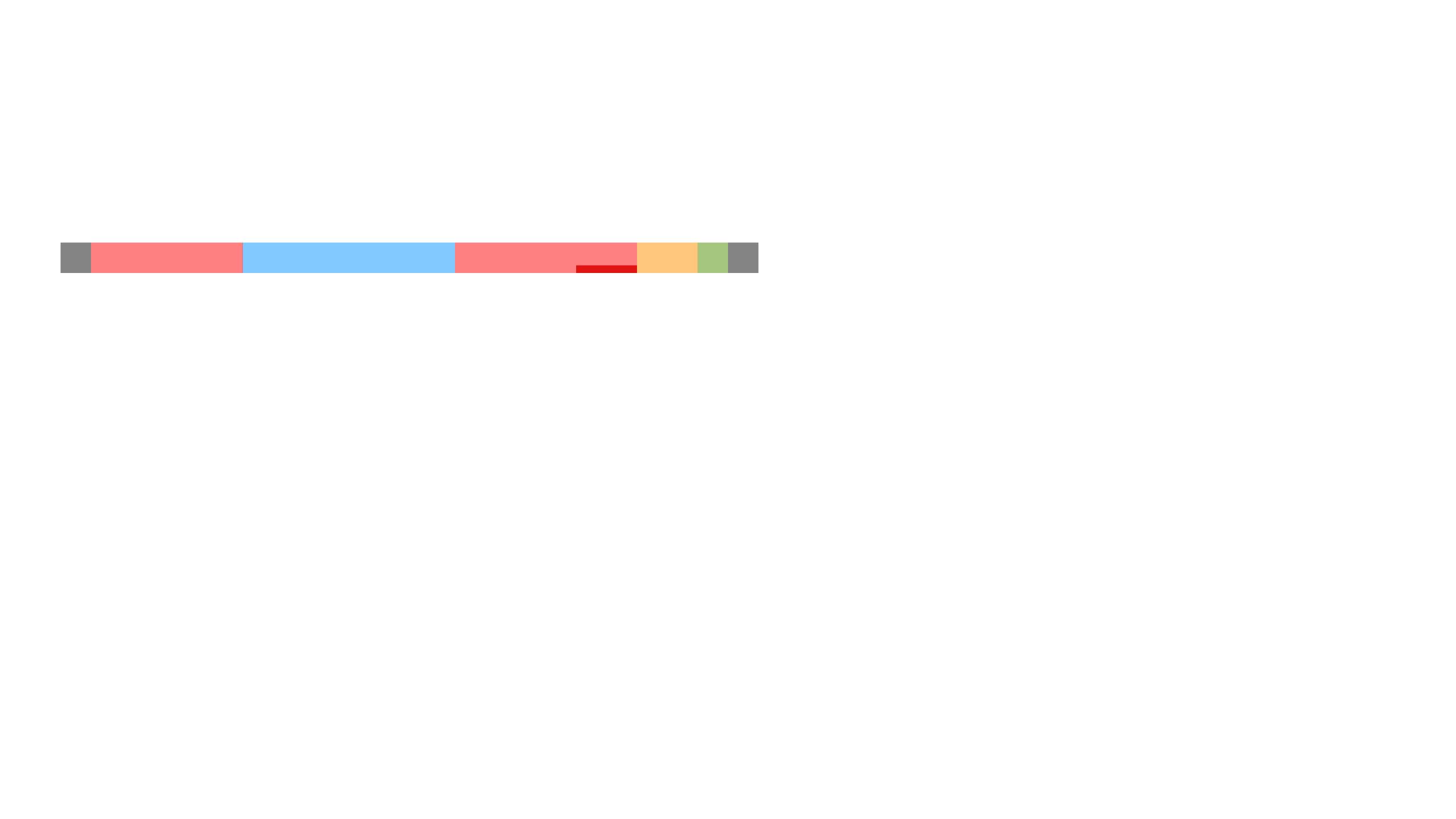


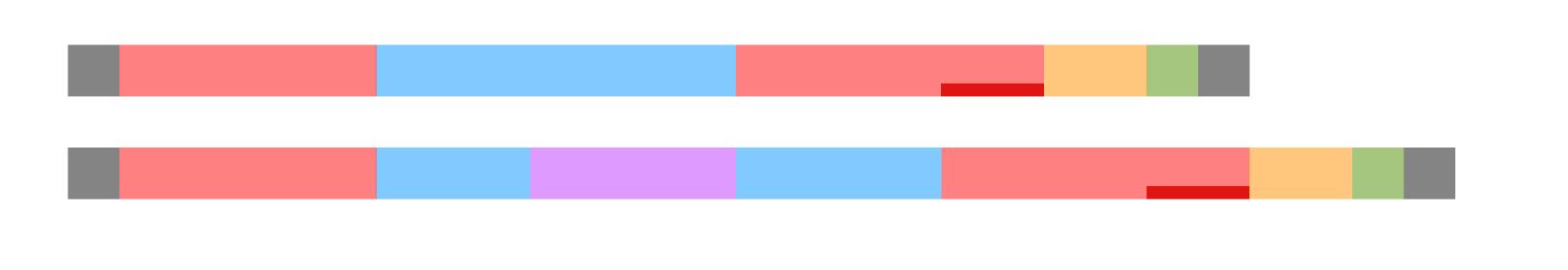


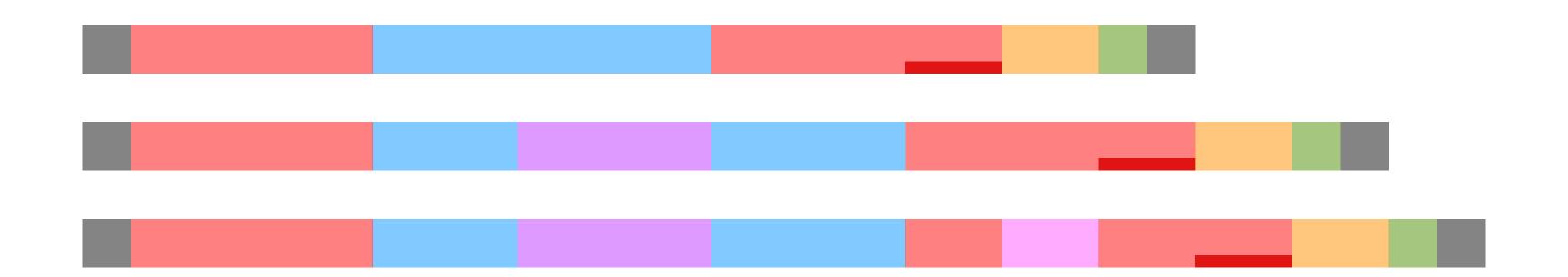


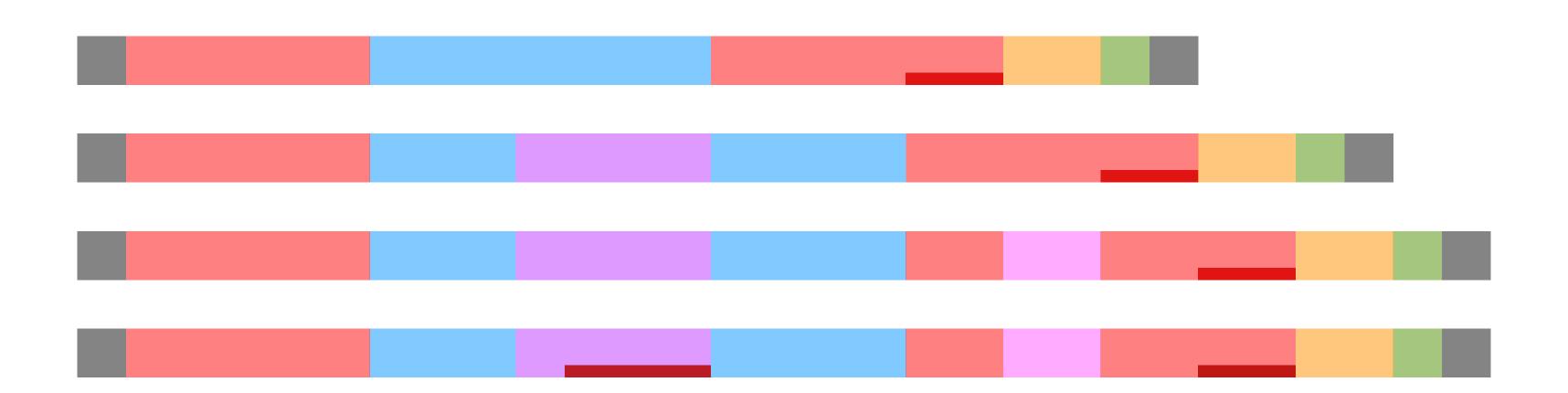
start hello, there worldms! END

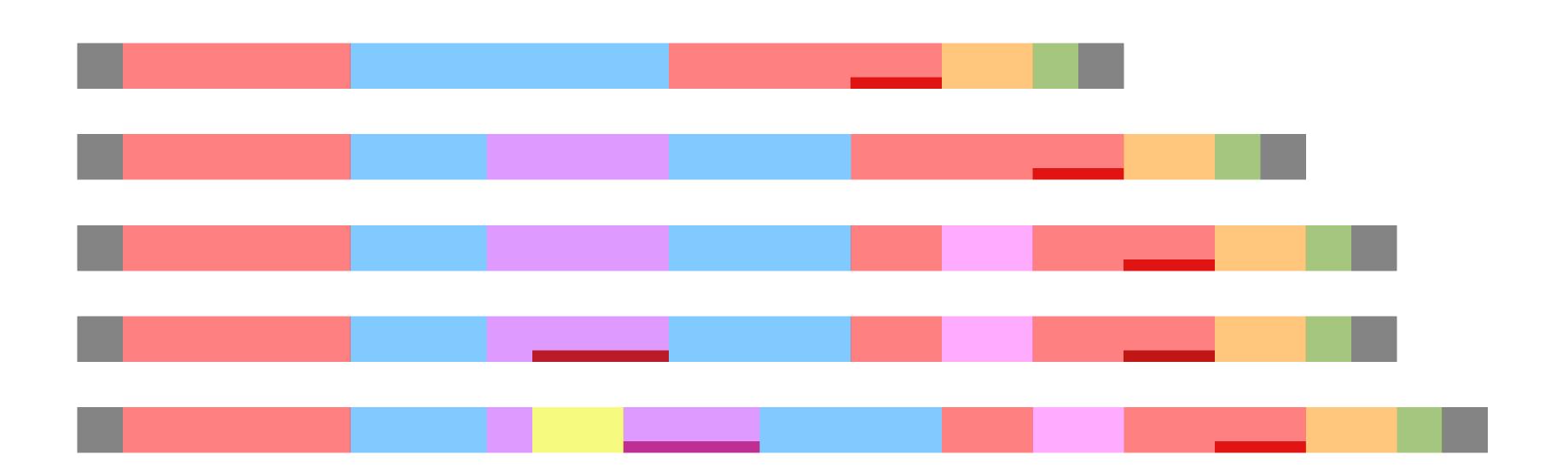
start hello, there worldms! END

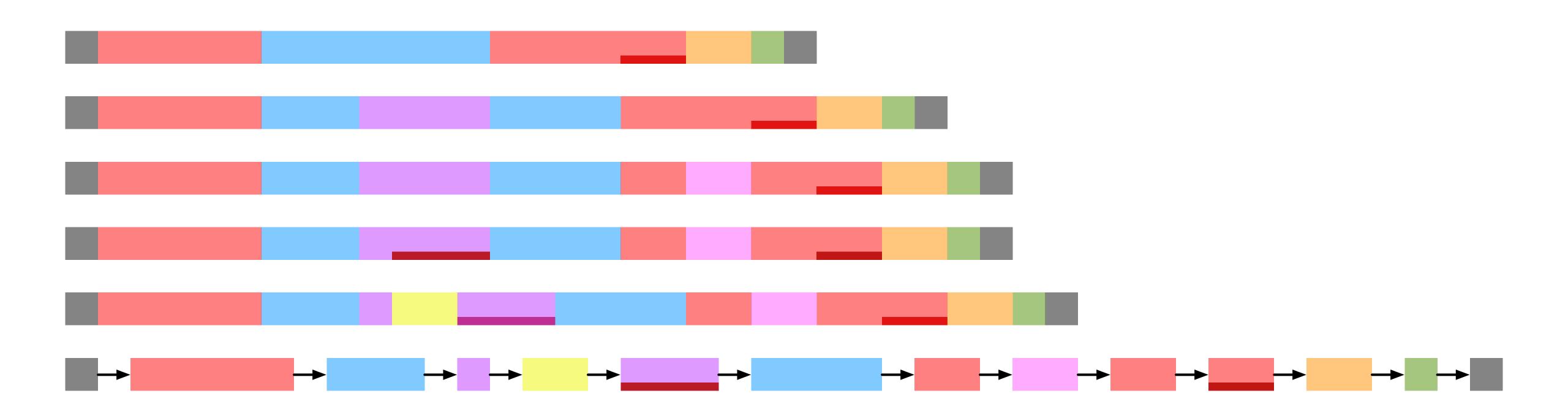


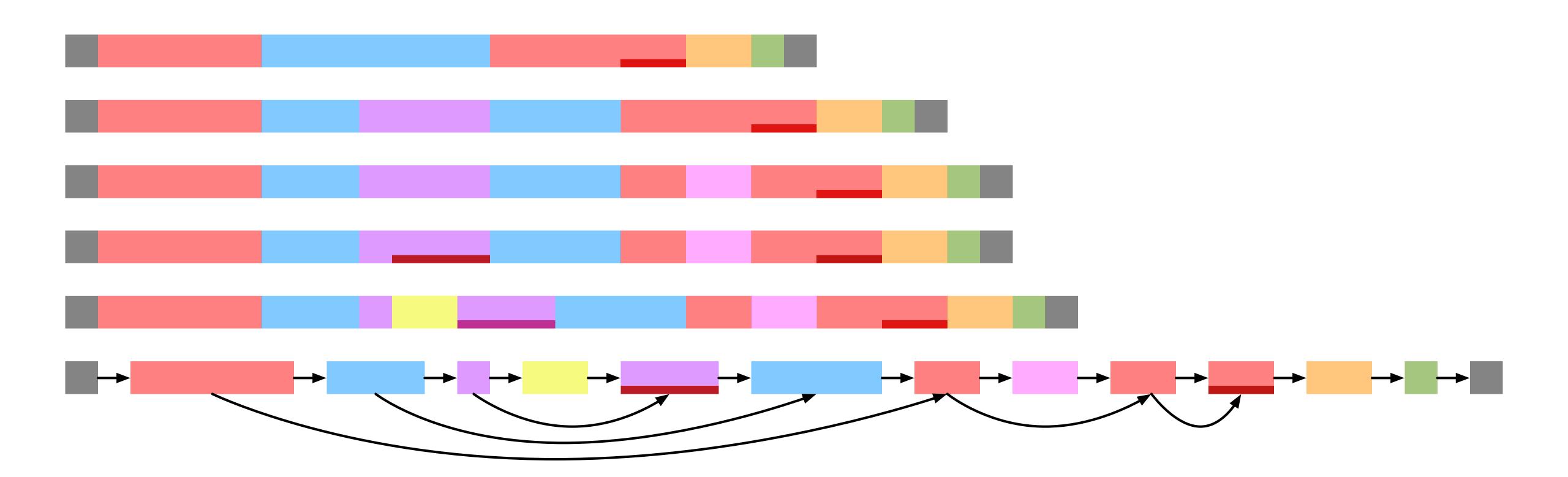


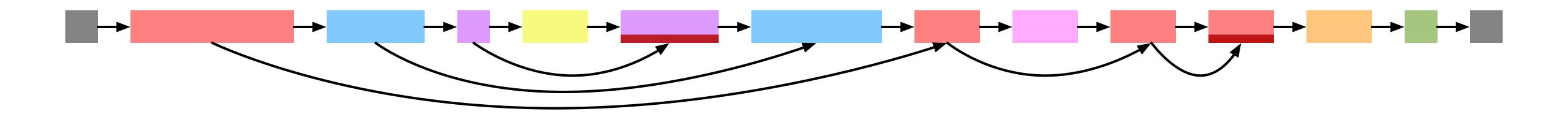


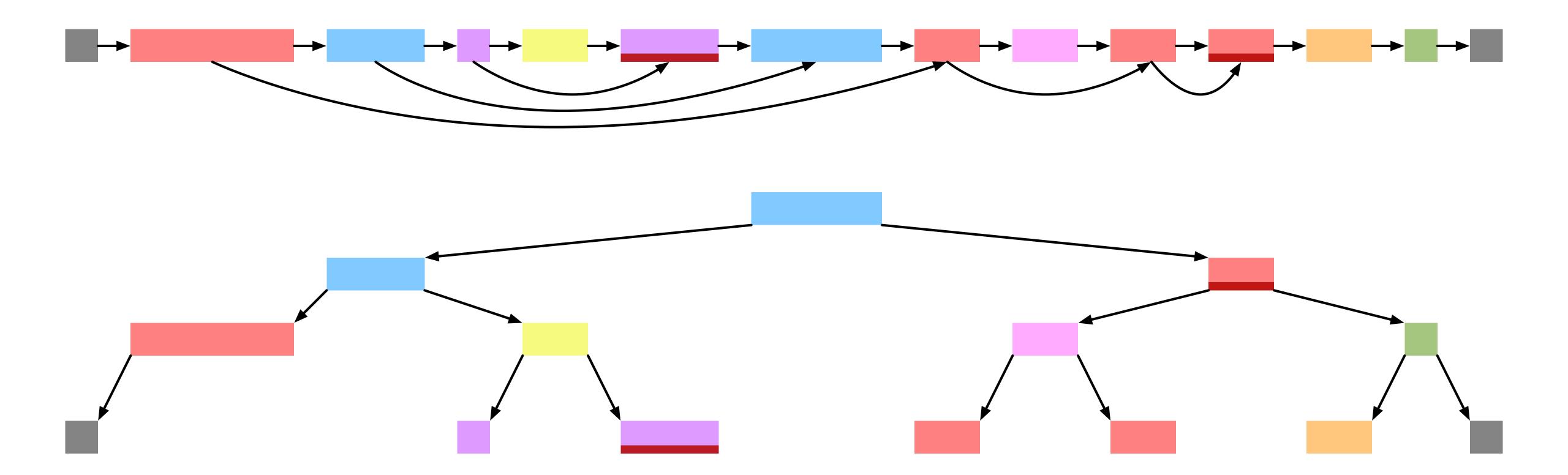


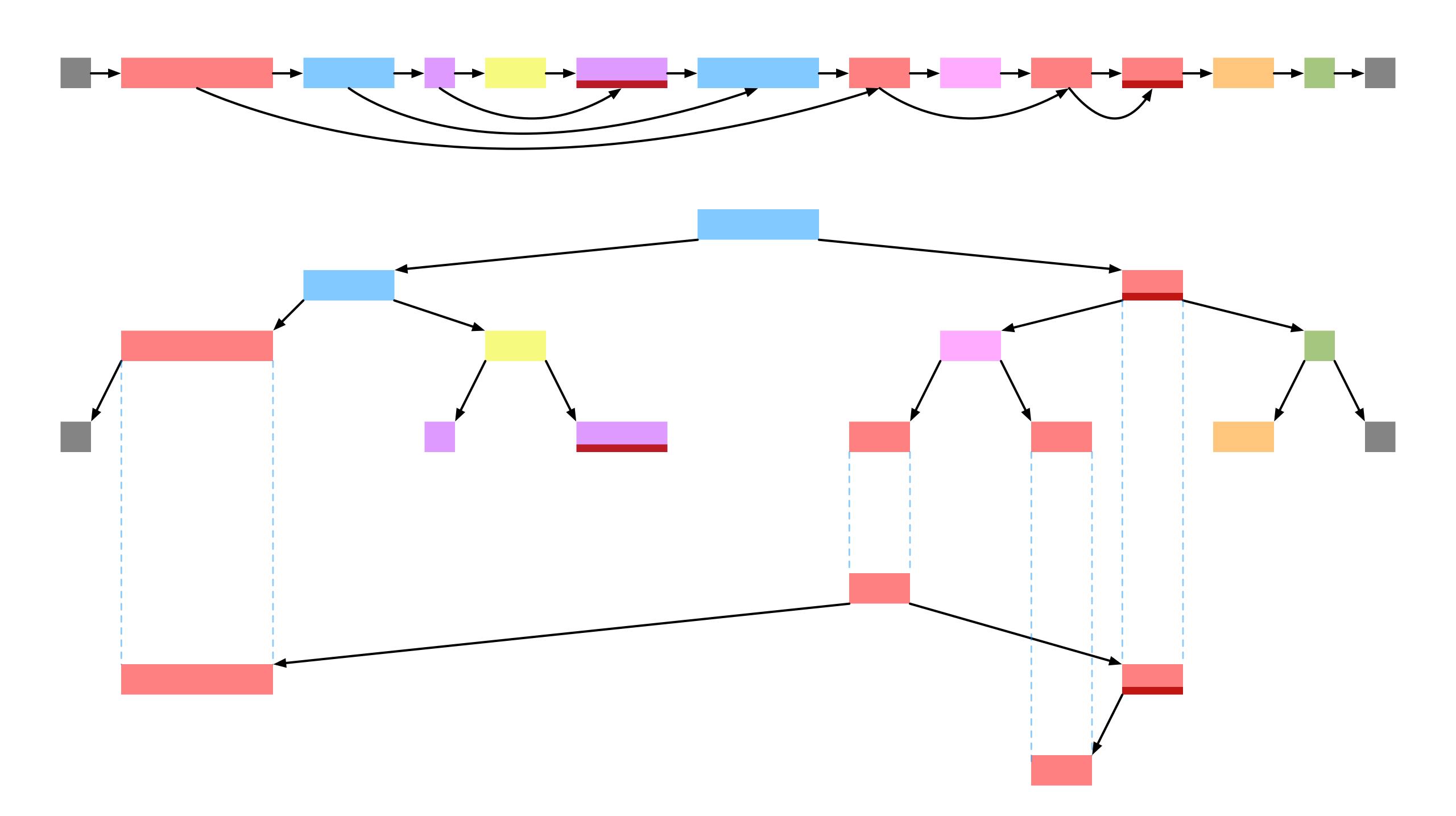


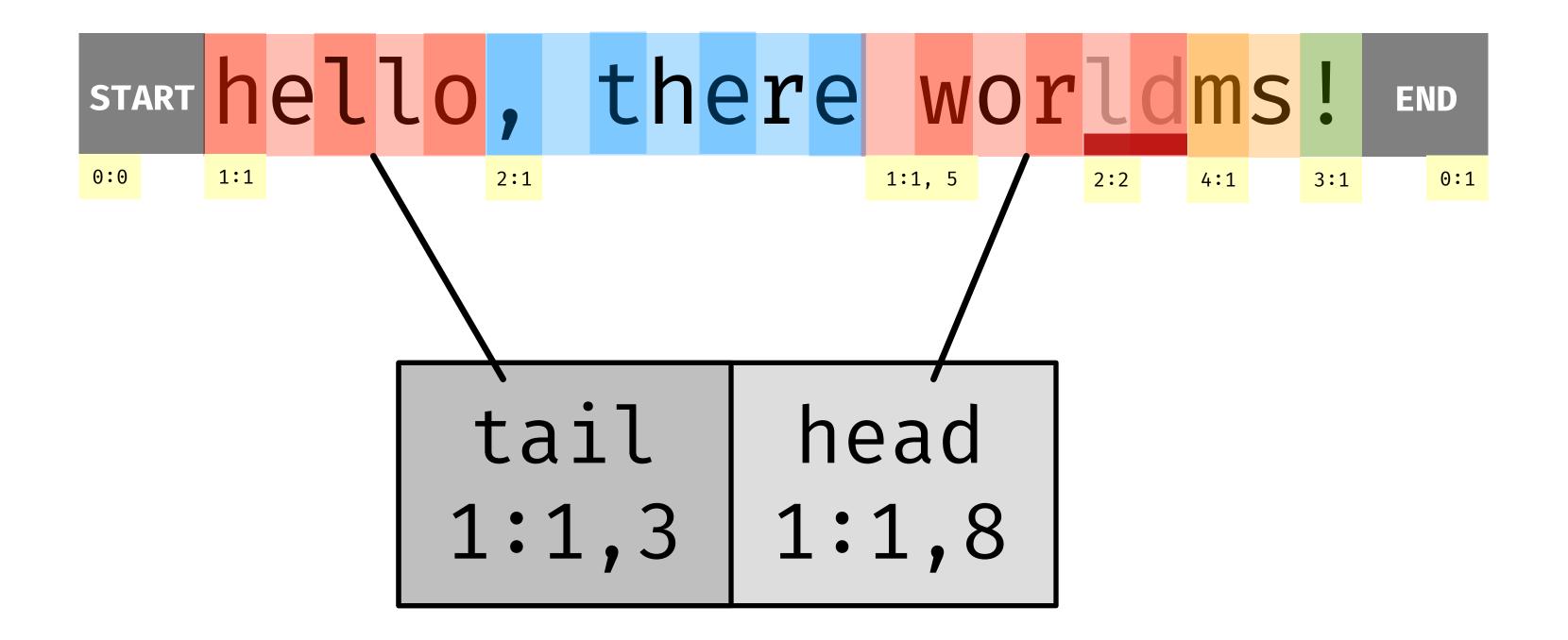














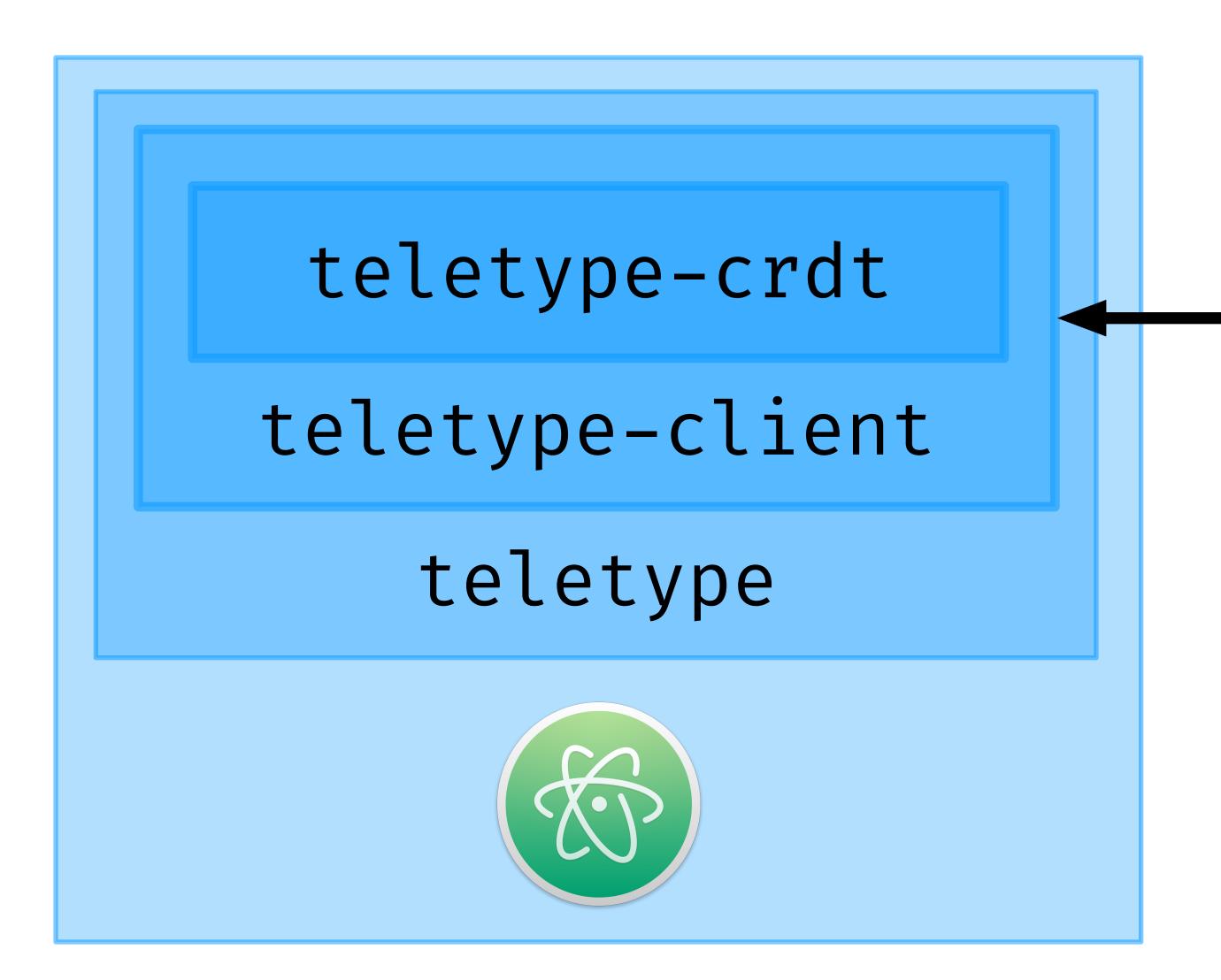
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teletype-crdt

teletype-client

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teletype-server



teletype-server



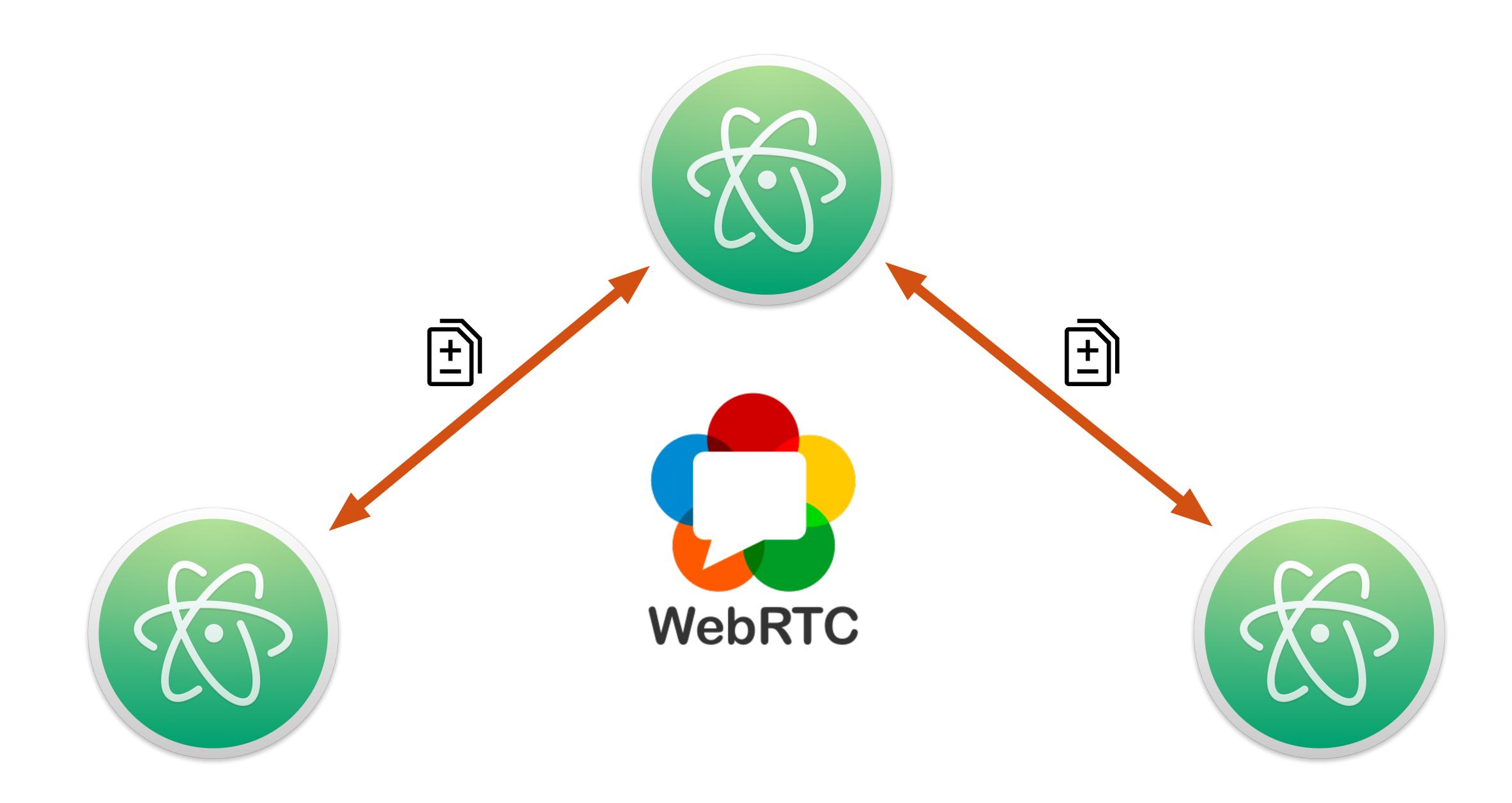




((•)) A







SPRAY: an Adaptive Random Peer Sampling Protocol

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ABSTRACT

The introduction of WebRTC has opened a new playground for large-scale distributed applications consisting of large numbers of directly-communicating web browsers. In this context, gossip-based peer-sampling protocols appear as a particularly promising tool thanks to their inherent ability to build overlay networks that can cope with network dynamics. However, the dynamic nature of browser-to-browser communication combined with the connection establishment procedures that characterize WebRTC make current peersampling solutions inefficient or simply unreliable. In this paper, we address the limitations of current peer-sampling approaches by introducing Spray, a novel peer-sampling protocol designed to avoid the constraints introduced by WebRTC. Unlike most recent peer-sampling approaches, Spray has the ability to adapt its operation to networks that can grow or shrink very rapidly. Moreover, by using only neighborto-neighbor interactions, it limits the impact of the threeway connection establishment process that characterizes WebRTC. Our experiments demonstrate the ability of Spray to adapt to dynamic networks and highlight its efficiency improvements with respect to existing protocols.

Categories and Subject Descriptors

C.2.1 [Computer Systems Organization]: Computer-Communication Networks—Network Architecture and Design

General Terms

Network, algorithm, simulation

Keywords

Large scale distributed applications, random peer sampling, browser-to-browser communication, WebRTC

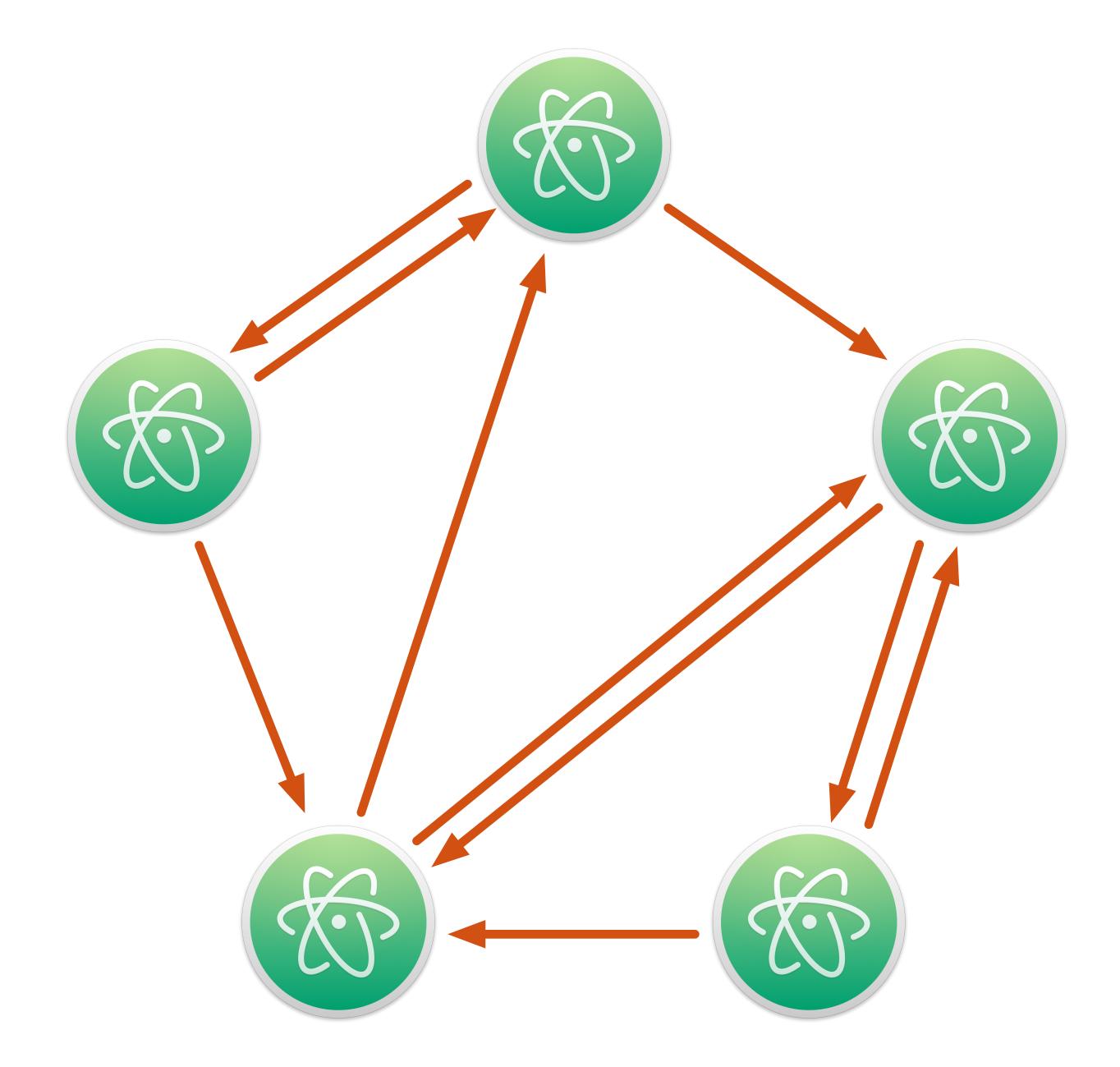
1. INTRODUCTION

Peer sampling constitutes a fundamental mechanism for many large-scale distributed applications both on the cloud [4] and in a peer-to-peer setting. Services such as dissemination [6, 19], aggregation [12] and network management [13,21] have been based on peer sampling and the recent introduction of WebRTC¹ opens the opportunity to deploy such applications on browsers that can run on laptops, desktops and mobile devices. In this context, WebRTC drastically simplifies deployment even within complex network systems that utilize firewalls, proxies and Net Address Translation (NAT).

Unfortunately, WebRTC has several constraints that make existing peer sampling services inefficient or unreliable. Browsers can run on small devices in mobile networks. Hence, keeping the number of connections as low as possible is a major requirement. However, peer sampling services such as CYCLON [20] do not adapt the number of connections to the real number of participants. For instance, a user must maintain 10 connections with other remote browsers when only 6 are enough. On the other hand, the peer sampling service SCAMP [10] is adaptive but uses random dissemination paths to establish connections which is much more costly and likely to fail in the WebRTC context.

In this paper, we introduce Spray, a random peer sampling protocol inspired by both Scamp [10] and Cyclon [20]. Compared to the state of art, (i) Spray dynamically adapts the neighborhood of each peer. Thus, the number of connections grows logarithmically compared to the size of the network. (ii) Spray only uses neighbor-to-neighbor interactions to establish connections. Thus, the connections are established in constant time. (iii) Spray quickly converges to a topology exposing properties similar to those of a random graph. Thus, the network becomes robust to massive failures, efficiently disseminates information etc. (iv) In the experimental setup, we show the adaptiveness of Spray and highlight its efficiency improvement compared to Cyclon and Scamp, at the cost of little overhead.

The rest of this paper is organized as follows: Section 2 reviews the related work. Section 3 details the SPRAY protocol. Section 4 shows the properties of SPRAY and compares them to state-of-the-art random peer sampling approaches. We conclude and discuss about the perspective in Section 5.



¹http://www.webrtc.org/



