Does Java need inline(value) types?

What project Valhalla can bring to Java from a performance perspective.

Sergey Kuksenko
Java Platform Group
Oracle
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Safe Harbor

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Who am I?

- Java/JVM Performance Engineer at Oracle, @since 2010
- Java/JVM Performance Engineer, @since 2005
- Java/JVM Engineer, @since 1996
What is Valhalla?
Valhalla Goals

• Provide denser memory layout (inline/value types)
• Specialized generics (including primitive, value types)
• Smooth library migration
• JVM cleanup (e.g. Nestmates a.k.a. JEP-181)
Object Identity is the root of all evil

Identity (object-oriented programming)

From Wikipedia, the free encyclopedia

An identity in object-oriented programming, object-oriented design and object-oriented analysis describes the property of objects that distinguishes them from other objects. This is closely related to the philosophical concept of identity.

In philosophy, identity, from Latin: identitas ("sameness"), is the relation each thing bears only to itself. The notion of identity gives rise to many philosophical problems, including the identity of indiscernibles, and questions about change and personal identity over time.
Identity gives

- Indirection
- Allocation in heap
- Nullability
- Mutability
- Reference equality (==)
- Locking
- Puzzlers, e.g.

```java
Integer.valueOf(42) == Integer.valueOf(42)
```

but

```java
Integer.valueOf(420) != Integer.valueOf(420)
```
Why JVM can’t eliminate it?
JVM can!
JVM can!
JVM can, but ... 

~300 articles per year!
Inline class

- is a class
- no identity
- immutable
- not nullable
- no synchronization
public  inline class Complex   {

    private  double re;
    private  double im;

    public Complex(double re, double im) {
        this.re = re;
        this.im = im;
    }
    public double re() { return re; }
    public double im() { return im; }
    ...
}
public inline class Complex {
  private double re;
  private double im;

  public Complex(double re, double im) {
    this.re = re;
    this.im = im;
  }
  public double re() { return re; }
  public double im() { return im; }
  ...
}
Identity of indiscernibles

‘==’ on inline classes

- deny at all
- always false
- substitutability check
  (recursive ‘==’ for each field)

OR

current choice
Inline class means inlineable

• JVM decides if:
  • allocate on heap
    OR
  • put on stack (locals, parameters, result)
  • inline into container class
  • inline into array (flattened array)
Inline class

- Inline types are subtypes of Object (interface)
- Inline arrays are covariant with Object[] (arrays of interface)
Boxing vs boxing

• V? - nullable twin of ‘V’
• means all values of V + ‘null’

compare to:

• Integer - nullable twin of ‘int’
• But Integer has full identity
Float like a butterfly, Sting like a bee
Code like a class, Work like an int
Local variable

```java
int count(Complex c) {
    Complex z = c;
    for (int i = 1; i < MAX_ITERATION; i++) {
        if (z.modulus() >= 4.0) return i;
        z = z.square().add(c);
    }
    return MAX_ITERATION;
}
```
Local variable

```java
int count(Complex c) {
    Complex z = c;
    for (int i = 1; i < MAX_ITERATION; i++) {
        if (z.modulus() >= 4.0) return i;
        z = z.square().add(c);
    }
    return MAX_ITERATION;
}
```

<table>
<thead>
<tr>
<th></th>
<th>average time(ns)</th>
</tr>
</thead>
<tbody>
<tr>
<td>Reference</td>
<td>485</td>
</tr>
<tr>
<td>primitive</td>
<td>350</td>
</tr>
<tr>
<td>Inline</td>
<td>350</td>
</tr>
</tbody>
</table>
Local variable

```java
int count(Complex c) {
    Complex z = c;
    for (int i = 1; i < MAX_ITERATION; i++) {
        if (z.modulus() >= 4.0) return i;
        z = z.square().add(c);
    }
    return MAX_ITERATION;
}
```

heap allocations(bytes/op)

<table>
<thead>
<tr>
<th>Reference</th>
<th>2120</th>
</tr>
</thead>
<tbody>
<tr>
<td>primitive</td>
<td>0</td>
</tr>
<tr>
<td>Inline</td>
<td>0</td>
</tr>
</tbody>
</table>
Reference vs Inline (Mandelbrot, 500x500)

Inline class:

- 4x less data loads
- 42x less L1 cache misses
- 5x less L3 cache misses
- 5x less dTLB misses
Scalability (Mandelbrot, 500x500)

- Reference
- Inline

16x
Method parameters/result

```java
static Value ackermann(Value x, Value y) {
    return x.isZero() ? y.inc() :
        (y.isZero() ? ackermann(x.dec(), new Value(1)) :
            ackermann(x.dec(), ackermann(x, y.dec())));
}
```
static Value ackermann(Value x, Value y) {
    return x.isZero() ? y.inc() :
    (y.isZero() ? ackermann(x.dec(), new Value(1)) :
     ackermann(x.dec(), ackermann(x, y.dec())));
}
Array access

ref

a
b
c

VS

a
b
c
Array Random Access

- Reference
- Primitive
- Inline

access time (ns) vs size of array (K elements)
Array Sequential Access

- Reference
- primitive
- Inline

access time (ns)

size of array (K elements)
Collateral Damage in legacy world
The following section is intended to outline Valhalla current status and development. It is intended for information purposes only, and may not be incorporated into any contract (or slowdown blaming). Any adverted performance regression maybe a subject to removal.
Inline in heap ("boxing")

- Object o = my_inline_value;
- Interface i = my_inline_value;
- Value? nullable_value = my_inline_value;
- JVM decided
Reference comparison

Glorious pre Valhalla past

just compare it

Brighter post Valhalla future

if <both refs are inline>
  if <classes are same>
    then
      check substitutability
    else
      false
  else
    just compare it
If object is inline class

ptr

Mark Word

Klass ptr

My Object

Class Desc
If object is inline class

Use Mark Word

ptr

Mark Word
Klass ptr

My Object

Class Desc
Mark Word

// 64 bits:
// --------
// unused:25 hash:31 -->| unused:1 age:4 biased_lock:1 lock:2 (normal object)
// JavaThread*:54 epoch:2 unused:1 age:4 biased_lock:1 lock:2 (biased object)
// "1" :54 epoch:2 unused:1 age:4 biased_lock:1 lock:2 (biased always locked object)
// PromotedObject*:61 --------------------->| promo_bits:3 ----->| (CMS promoted object)
// size:64 ------------------------------>| (CMS free block)

// unused:25 hash:31 -->| cms_free:1 age:4 biased_lock:1 lock:2 (COOPs && normal object)
// JavaThread*:54 epoch:2 cms_free:1 age:4 biased_lock:1 lock:2 (COOPs && biased object)
// narrowOop:32 unused:24 cms_free:1 unused:4 promo_bits:3 ----->| (COOPs && CMS promoted object)
// unused:21 size:35 -->| cms_free:1 unused:7 ------------------>| (COOPs && CMS free block)
Mark Word

// 64 bits:
// -----
// unused:25 hash:31 -->| unused:1 age:4 biased_lock:1 lock:2 (normal object)
// JavaThread*:54 epoch:2 unused:1 age:4 biased_lock:1 lock:2 (biased object)
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// narrowOop:32 unused:24 cms_free:1 unused:4 promo_bits:3 ----->| (COOPs && CMS promoted object)
// unused:21 size:35 -->| cms_free:1 unused:7 ------------------>| (COOPs && CMS free block)

// [ <unused> | larval |1| epoch | age | 1 | 01] permanently locked
Reference comparison a.k.a. ‘acmp’

```java
...
if (o1 == o2) {
    ... = x + 1;
} else {
    ... = x - 1;
}
...
```
acmp  -XX:--EnableValhalla

cmp %rdx,%rcx

inc %eax  dec %eax
acmp  -XX:+EnableValhalla

cmp %rdx,%rcx

test %rcx,%rcx

mov $0x405,%r10
and (%rcx),%r10
cmp $0x405,%r10

test %rdx,%rdx

mov 0x8(%rdx),%r11
mov 0x8(%rcx),%r10
cmp %r11,%r10

dec %eax

#invokedynamic ...

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acmp  -XX:+EnableValhalla

cmp  %rdx,%rcx

inc  %eax

dec  %eax

mov  0x8(%rdx),%r11
mov  0x8(%rcx),%r10
cmp  %r11,%r10

mov  $0x405,%r10
and  (%rcx),%r10
cmp  $0x405,%r10

test  %rdx,%rdx

mov  $0x405,%r10
and  (%rcx),%r10
cmp  $0x405,%r10

test  %rcx,%rcx

#invokedynamic  ...

Reference World

Value World

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acmp

• Complex code
• Additional loads
• invokedynamic prevents loop unrolling
acmp performance

- Baseline, isolated
- Baseline, unrolled loop
- Valhalla

Time, ns

% of equal elements
synchronized(obj)

Glorious pre Valhalla past

do all synch stuff

Brighter post Valhalla future

if <ref is inline class>
then
  throw exception
else
  do all synch stuff
synchronized(obj)

do all synch stuff

Glorious pre Valhalla past

if <ref is inline class>
then
  throw exception
else
  do all synch stuff

Brighter post Valhalla future

< 1% difference
Arrays (Object[]) 

Object[] may be:

- Good old reference array
- Array of inline classes in heap
  - references, but not nullable
- Flattened array of inline classes
Arrays (Object[])
Arrays (Object[])
Arrays (Object[]) 

• Any access to Klass ptr required clearing:
  • and $0x1fffffff,%reg

• HotSpot is good enough at eliminating it
  • knowing that it isn’t an array
Load from Object[]

Glorious pre Valhalla past

element size is the same:
just load it

Brighter post Valhalla future

if <array is flattened>
    find element size
    load it
    do boxing if needed
else
    just load it
Store to Object[]

Glorious pre Valhalla past

do ArrayStoreCheck
store if ok

Brighter post Valhalla future

do ArrayStoreCheck
if <array is flattened>
  find element size
  do unboxing if needed
store
else
store
Object[] access

• Targeting benchmarks: −2% . . . − 10%
• Solution: aggressive loop hoisting and loop duplication (in progress)
Inline vs inline

Integer[] \texttt{i1} = \texttt{new Integer[1000]};
Integer[] \texttt{i2} = \texttt{new Integer[1000]};

@Setup
public void setup() {
    for (int \texttt{i} = \texttt{0}; \texttt{i} < \texttt{1000}; \texttt{i}++)
        \texttt{i1[i]} = \texttt{i2[i]} = \texttt{i};
    \texttt{i2[999]} = \texttt{394857623};
}

@Benchmark
public boolean arrayEquals() {
    return Arrays.equals(\texttt{i1}, \texttt{i2});
}
Inline vs inline

Integer[] i1 = new Integer[1000];
Integer[] i2 = new Integer[1000];

@Setup
public void setup() {
    for (int i = 0; i < 1000; i++)
        i1[i] = i2[i] = i;
    i2[999] = 394857623;
}

@Benchmark
public boolean arrayEquals() {
    return Arrays.equals(i1, i2);
}
Methods inline tree

-XX:-EnableValhalla
...
- @ java.util.Arrays::equals [...] (inlined: inline (hot))
  - @ java.util.Objects::equals [...] (inlined: inline (hot))

-XX:+EnableValhalla
...
- @ java.util.Arrays::equals [...] (inline failed: already compiled into a big method)
Methods inline tree

-XX:-EnableValhalla
...
- @ java.util.Arrays::equals [..., bytes=57, insts=352]
  (inlined: inline (hot))
- @ java.util.Objects::equals [..., bytes=23, insts=128]
  (inlined: inline (hot))

-XX:+EnableValhalla
...
- @ java.util.Arrays::equals [..., bytes=57, insts=1760]
  (inline failed: already compiled into a big method)
Current status

• Checked ~30 big benchmarks:
  • No regressions more than 2%

• Checked ~1600 microbenchmarks:
  • 1200 – ±0% at the first run
  • 300 – fixed
  • 100 – less 5% (in progress)
  • 1 – 14% regression (in progress)
Brighter post Valhalla future
## Arithmetic types

- **Complex matrix multiplication (100x100)**

<table>
<thead>
<tr>
<th>Method</th>
<th>Time</th>
</tr>
</thead>
<tbody>
<tr>
<td>ref Complex</td>
<td>12.6 ms</td>
</tr>
<tr>
<td>inline Complex</td>
<td>2.7 ms</td>
</tr>
<tr>
<td>inline Complex + cache friendly algorithm</td>
<td>2.1 ms</td>
</tr>
</tbody>
</table>
java.util.Optional

public class HashMap<K, V> ... {
    ...
    /**
     * Returns the value to which the specified key is mapped,
     * or null if this map contains no mapping for the key.
     * ...
     */
    public V get(Object key)
public class HashMap<K, V> ... {
    ...
    /**
     * Returns the value to which the specified key is mapped,
     * or null if this map contains no mapping for the key.
     * ...
     */
    public V get(Object key)
    Not a good idea
java.util.Optional

```java
class HashMap<K, V> {
    public Optional<V> get(Object key) {
        // Returns an Optional describing the value to which the specified
        // key is mapped, or an empty Optional if this map contains no
        // mapping for the key.
    }
}
```

What if?
java.util.Optional

- 1000000 gets from 1000000 map

<table>
<thead>
<tr>
<th>Method</th>
<th>Time (ms)</th>
</tr>
</thead>
<tbody>
<tr>
<td>ref Optional</td>
<td>112.7</td>
</tr>
<tr>
<td>inline Optional</td>
<td>71.5</td>
</tr>
</tbody>
</table>
Example: map with complex key

- Map from $<K_1, \ldots, K_N> \rightarrow <V>$

  Two ways to implement:

1. Map<$\text{CompositeKey}<K_1, \ldots, K_N>, V>$

2. Map<$K_1$, Map $\ldots$, Map<$K_N$, V>...$>
Example: Map from \(<\text{Integer, Integer}> \rightarrow \text{<Integer}>\)

- 1000000 gets from 1000000 map

<table>
<thead>
<tr>
<th>Time (ms)</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>ref CompositeKey</td>
<td>146</td>
</tr>
<tr>
<td>Map of Map ...</td>
<td>122</td>
</tr>
<tr>
<td>inline CompositeKey</td>
<td>85</td>
</tr>
</tbody>
</table>
HashMap inside

VS
HashMap experiments

- Classic HashMap.get(): ~75 ns
- Experimental HashMap.get(): ~60 ns

? Faster only for large maps
? ‘put’ is slower
... to be continued
Iteration

HashMap<Integer, Integer> map; // map.size() == 1000000

@Benchmark
public int sumValuesInMap() {
    int s = 0;
    for (Integer i : map.values()) {
        s += i;
    }
    return s;
}
Iteration

HashMap<Integer, Integer> map; // map.size() == 1000000

@Benchmark
public int sumValuesInMap() {
    int s = 0;
    for (Integer i : map.values()) {
        s += i;
    }
    return s;
}

Lucky case 33 ms
Unlucky case 55 ms
Iteration

HashMap<Integer, Integer> map; // map.size() == 1000000

@Benchmark
public int sumValuesInMap() {
    int s = 0;
    for (Iterator<Integer> iterator = map.values().iterator();
         iterator.hasNext(); ) {
        s += iterator.next();
    }
    return s;
}

<table>
<thead>
<tr>
<th>scalarized</th>
<th>Lucky case</th>
<th>33 ms</th>
</tr>
</thead>
<tbody>
<tr>
<td>on heap</td>
<td>Unlucky case</td>
<td>55 ms</td>
</tr>
</tbody>
</table>
Inside HashMap

abstract class HashIterator {
    ...
    Node<K,V> next;       // next entry to return
    Node<K,V> current;    // current entry

    final Node<K,V> nextNode() {
        ...
        if ((next = (current = e).next) == null && (t = table) != null) {
            do {}
            while (index < t.length && (next = t[index++]) == null);
        }
        return e;
    }
    ...
}
Inside HashMap

**abstract class** HashIterator {

...  
Node<K, V> next;  // next entry to return  
Node<K, V> current;  // current entry  

final Node<K, V> nextNode() {

...  
if ((next = (current = e).next) == null && (t = table) != null) {
  do {}  
  while (index < t.length && (next = t[index++]) == null);  
}
return e;
}
...  

Write to reference field
GC write barriers
inline Cursor

```java
public interface Cursor<V> {
    boolean hasElement();
    V get();
    Cursor<V> next();
}
```
inline Cursor

@Benchmark
dpublic int sumValuesInMap() {
    int s = 0;
    for (Cursor<Integer> cursor = map.values().cursor();
         cursor.hasElement();
         cursor = cursor.next()) {
        s += cursor.get();
    }
    return s;
}
inline Cursor

@Benchmark
dpublic int sumValuesInMap() {
    int s = 0;
    for (Cursor<Integer> cursor = map.values().cursor();
         cursor.hasElement();
         cursor = cursor.next()) {
        s += cursor.get();
    }
    return s;
}
Move/Copy data

• Reference – easy, only reference is moved
• Inline – all data should be moved
  – Sort:
    Reference – default TimSort from JDK
    Inline – reimplemented TimSort
    indexed Inline – sort ‘int’ indices first, then copy
size == 400, fit into L1 cache

- Reference
- Inline
- indexed Inline

Consistent performance with varying class size.
size == 4000, fit into L2 cache

```
time, microseconds
```

- Reference
- Inline
- indexed Inline

```
class size (bytes)
```

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size==40000, fit into L3 cache
size == 400000, slightly > L3 cache
size == 4000000, much more L3 cache
Move/Copy data

• Dense location is better than moving less
Inline classes

• Dense, HW-friendly memory layout
• More control on heap allocations:
  • less GC pressure
  • less GC barriers

Better performance!
Links

• Wiki:
  https://wiki.openjdk.java.net/display/valhalla/Main

• Mailing lists:
  http://mail.openjdk.java.net/mailman/listinfo/valhalla-dev
  http://mail.openjdk.java.net/mailman/listinfo/valhalla-spec-observers

• Repository:
  http://hg.openjdk.java.net/valhalla